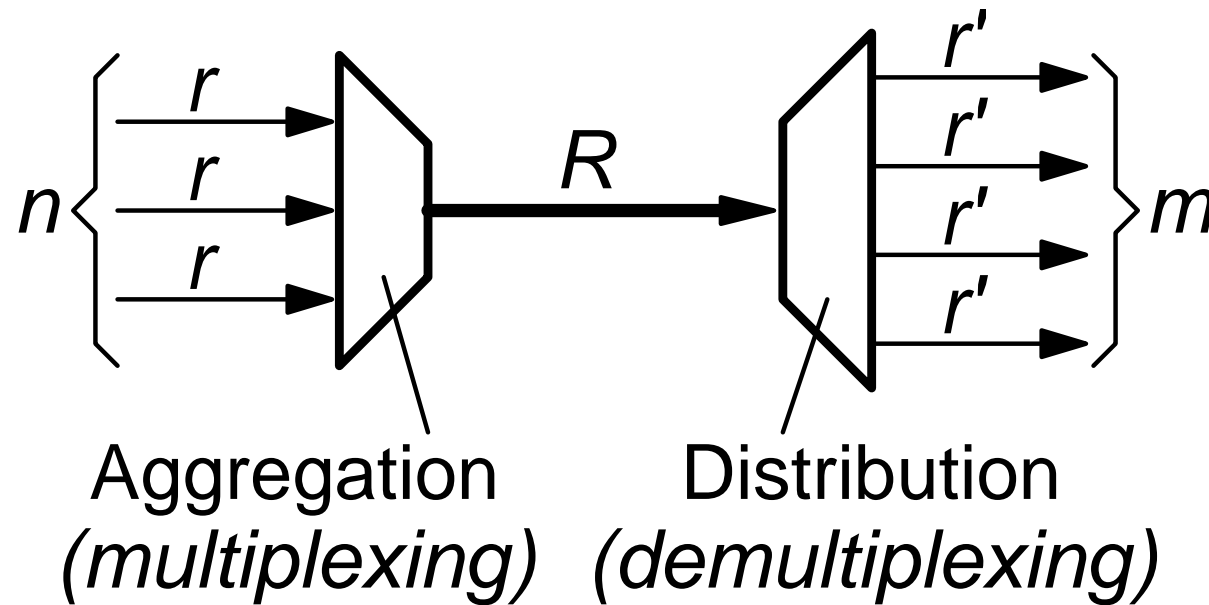
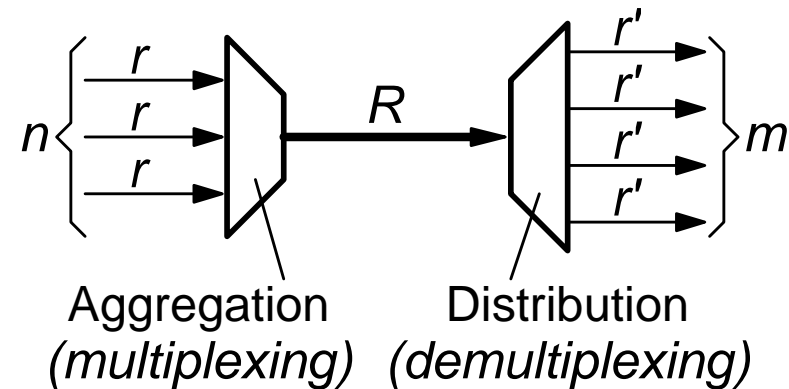


## 1.3 Multiplexing, Time-Switching, Point-to-Point versus Buses



- Simplest Networking, like simplest programming:  
Sequential !

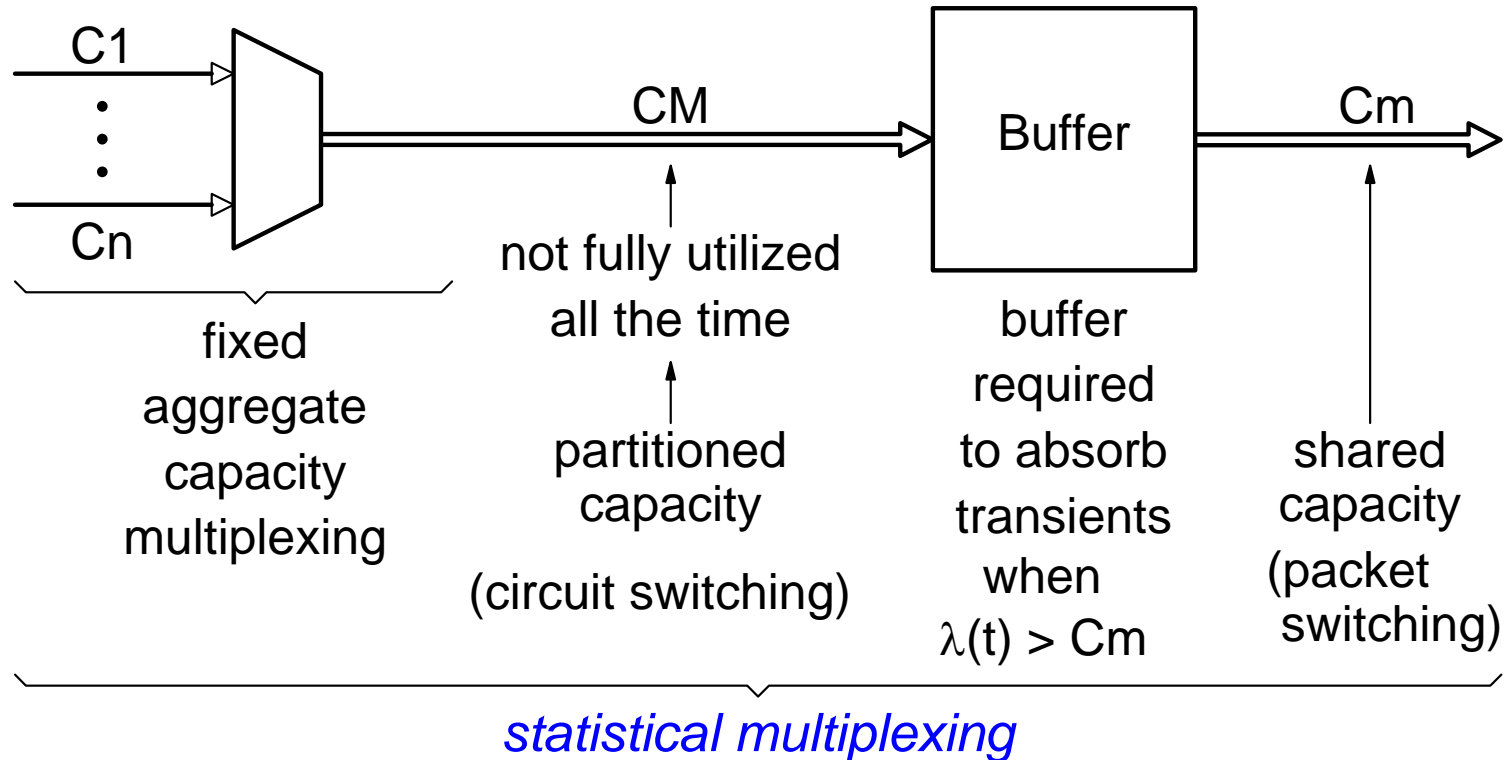
## Time-Switching



- Shared Medium Communication
- Demultiplexor determines where each piece of information is routed to by selecting the time at which each output receives information from the shared medium
- Simplicity: one thing at a time – no parallelism
- Non-scalable! – cannot increase  $R$  indefinitely
- Full-capacity ( $R = n \cdot r$ ) or Statistical Multiplexing ( $R < n \cdot r$ )
- Implementation issues, point-to-point links versus bidirectional, shared physical medium

# Statistical Multiplexing

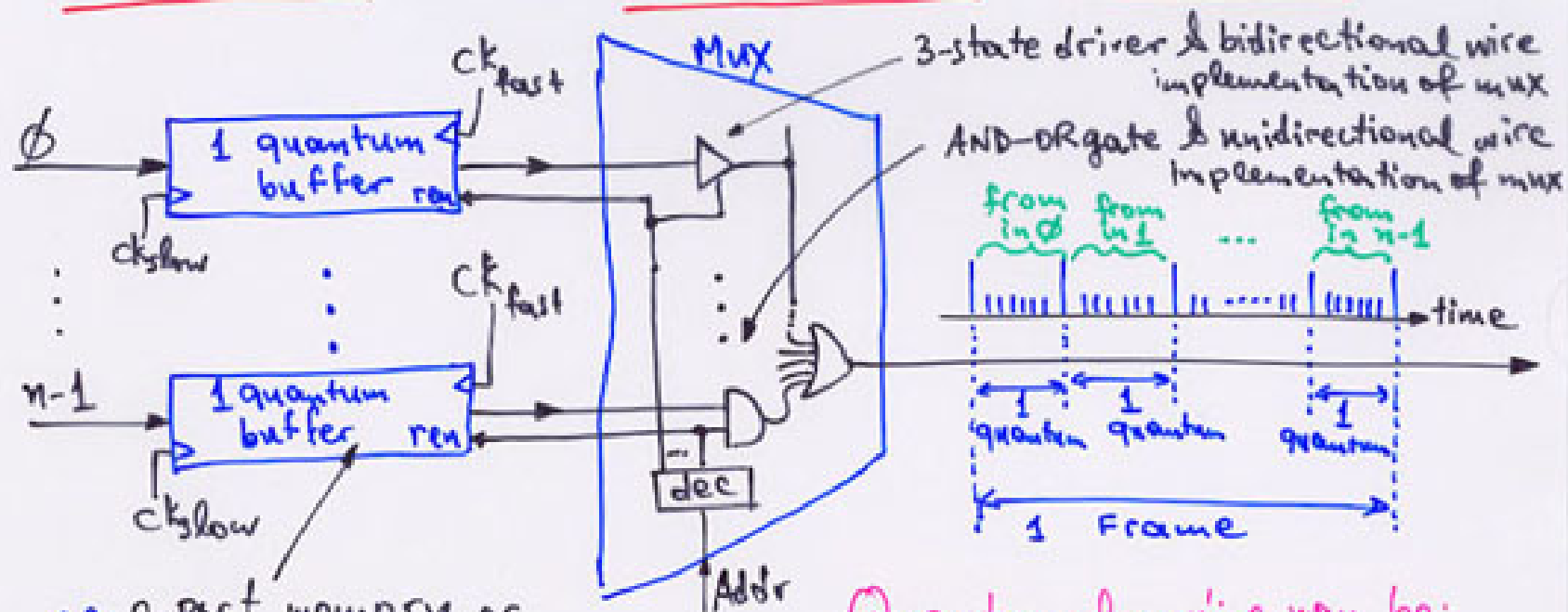
$$C_1 + \dots + C_n = C_M > C_m$$



Capacity  $C_m$  is better utilized than capacity  $C_M$

The average aggregate source rate must not exceed  $C_m$  over time windows determined by the size of the buffer memory

# Same-Width (per link), Faster-Clock Multiplexor:



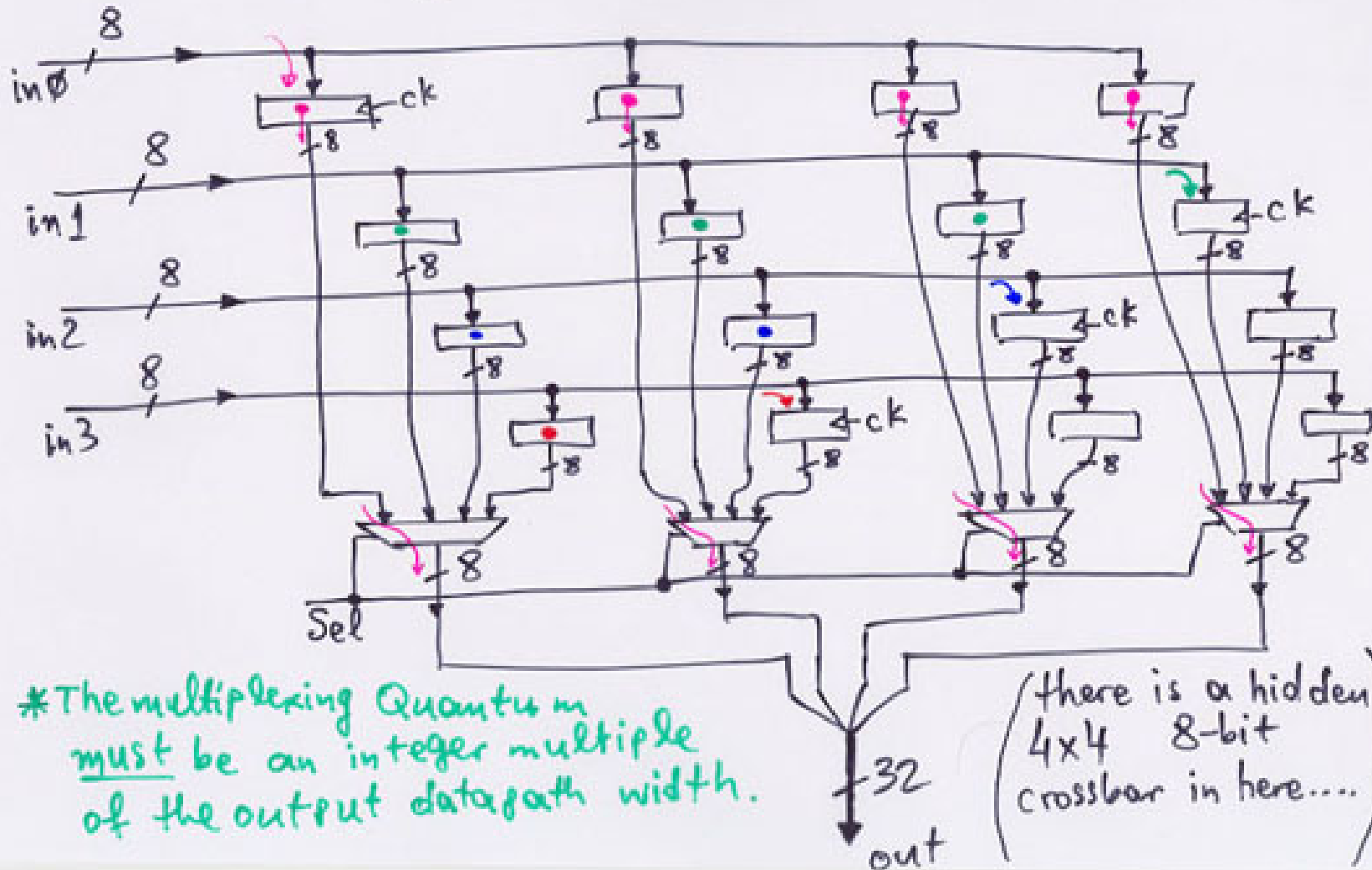
e.g. 2-port memory, or two shift registers using double buffering  
 (properly "stew-synchronized" input  $\Rightarrow$  buf.size = 1 quantum;  
 non-synchronized input  $\Rightarrow$  buf.size = 2 quanta (double buffering needed))

## Quantum of mux'ing may be:

- 1 bit {
  - parallel to serial converter
  - digital telephony - DS2, DS3
- 1 Byte . digital telephony - {DS1, SONET
- 1 Cell or 1 Block {
  - ATM
  - bus-oriented switches
  - other (high-throughput) switches, internally
- 1 Packet . packet switching

## Same-Clock, Wider-Datapath Multiplexor:

Example: 4-to-1 mux, 8-bit inputs, 32-bit output, common clock multiplexing quantum = 32 bits, skew-synchronized inputs

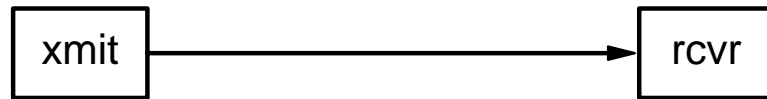


\*The multiplexing Quantum must be an integer multiple of the output datapath width.

(there is a hidden 4x4 8-bit crossbar in here....)

# Point-to-Point versus Shared Bidirectional Links

## Point-to-Point:

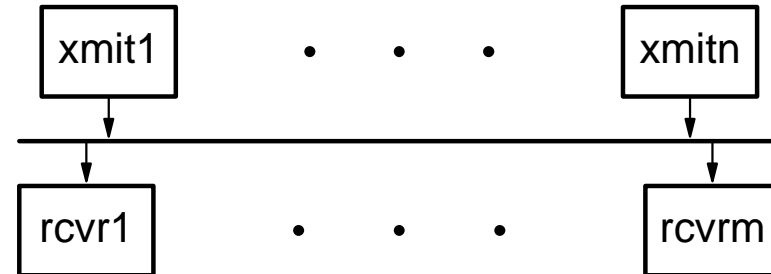


(e.g. high-speed copper or fiber links)

### ⇒ Higher Performance:

- no turn-around delay
- no arbitration overhead
- increased parallelism  
when used with switches

## Shared Medium:

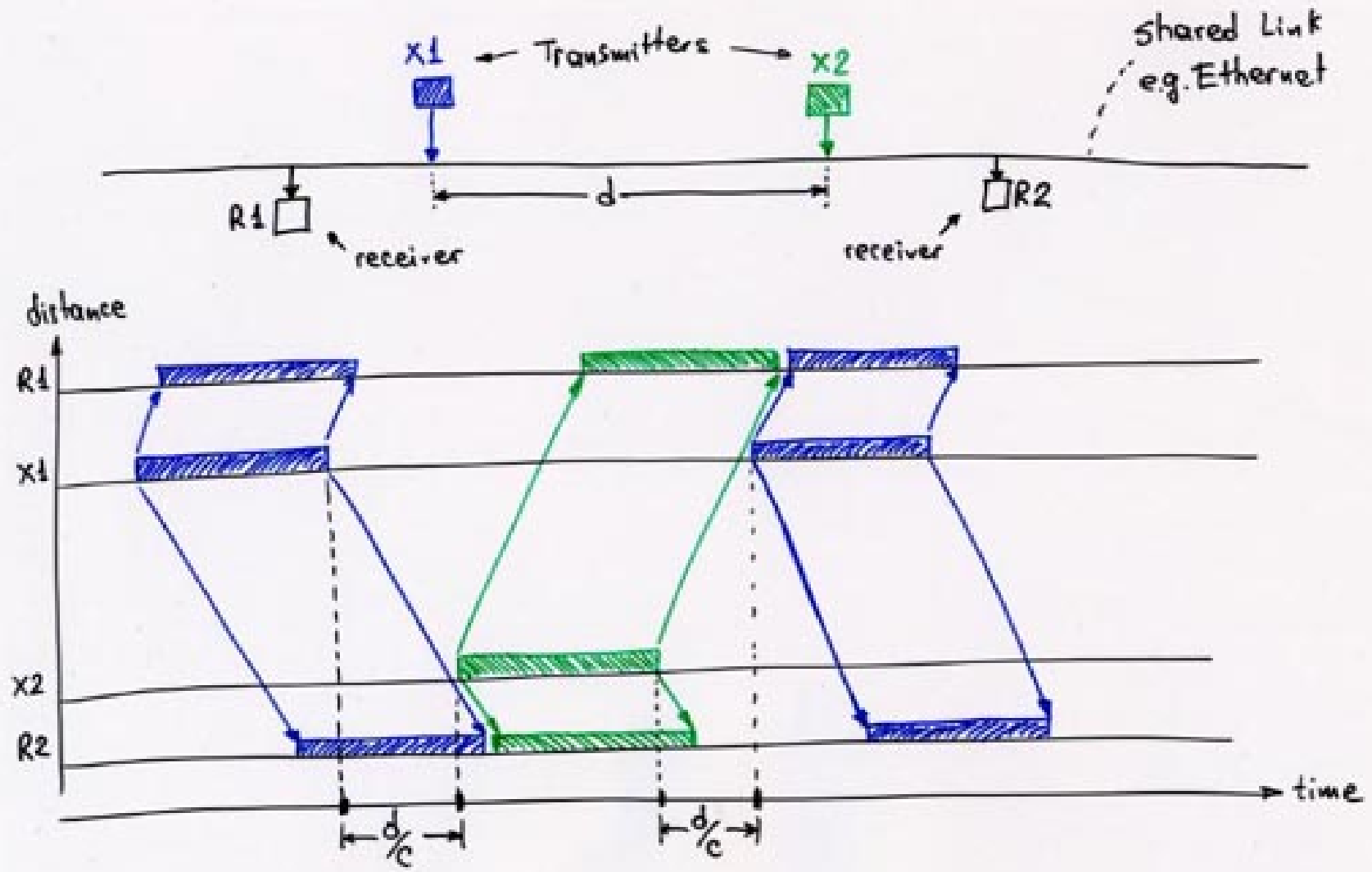


(e.g. wireless links, old ethernet, buses)

### ⇒ Lower Cost:

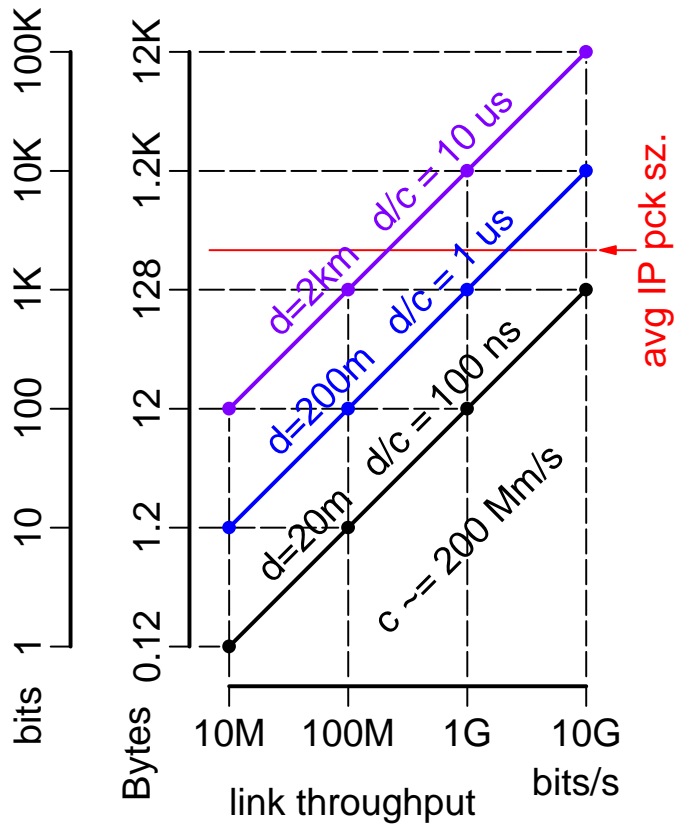
- broadcast & select: switching is inherent in the medium
- natural in some environments  
--e.g. wireless without directional antennas

# Shared Medium Link: Turn-Around Overhead



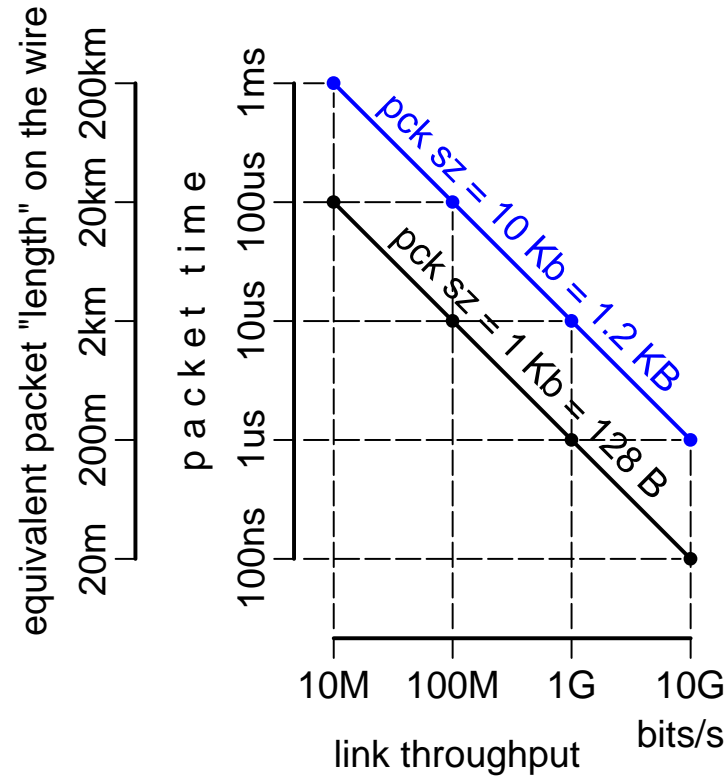
## How many bits is the length of the wire?

Turn-around delay expressed as a lost opportunity to transmit an amount of information equivalent to:



## How long is a packet on the wire?

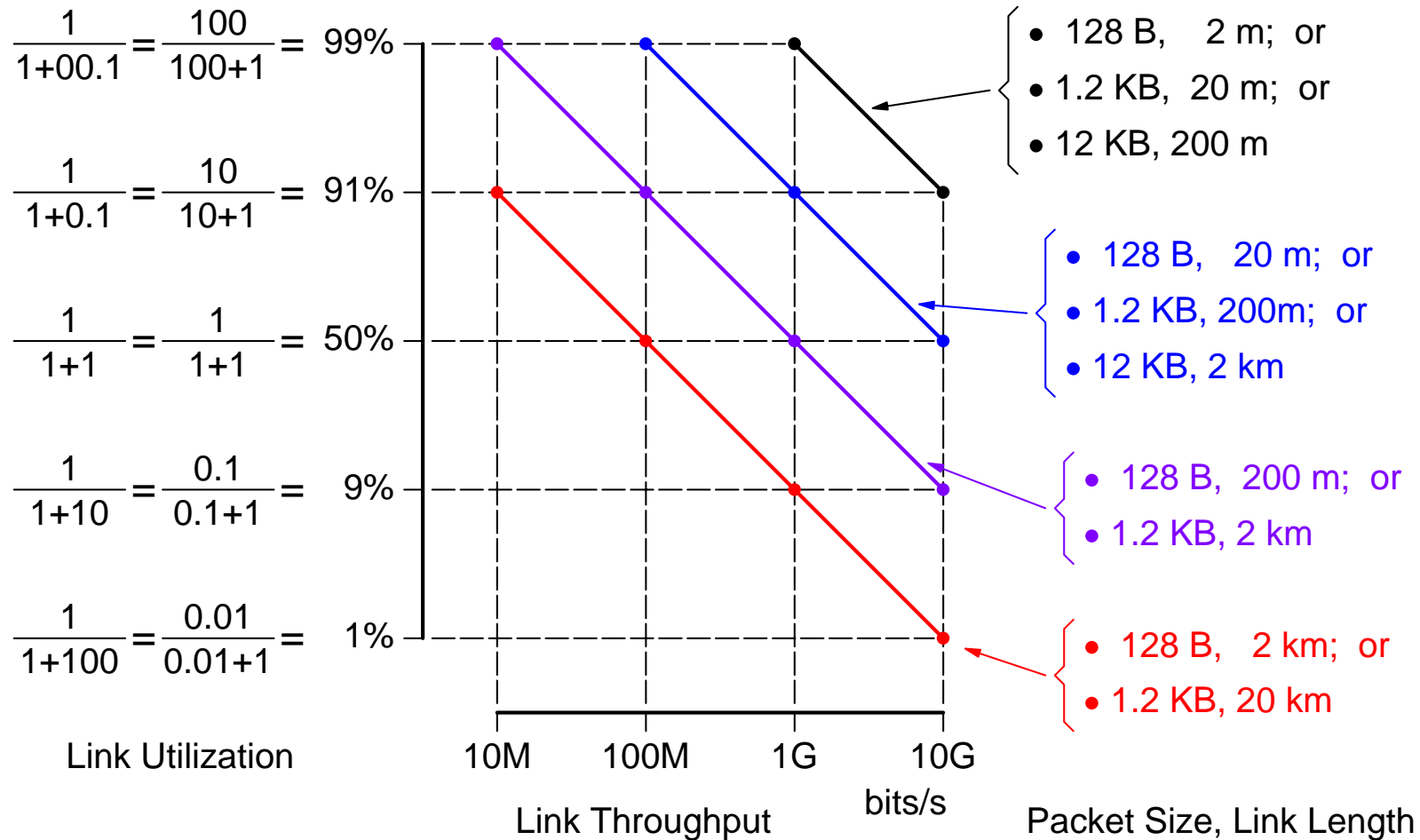
Duration of each transmit session (assume one packet), in time or in equivalent packet "length" on the wire:



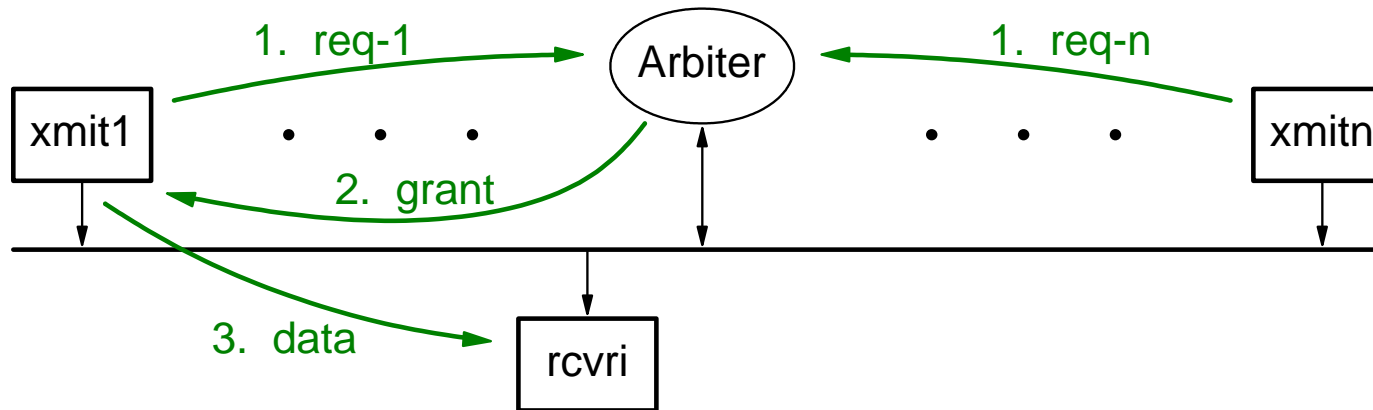


# Link Utilization = f(Packet Length, Wire Length, Throughput)

$$\text{Link Utilization} = \frac{\text{packetSize}}{\text{packetSize} + \text{turnAroundBitEquiv}} = \frac{\text{packetTime}}{\text{packetTime} + \text{turnAroundDelay}}$$



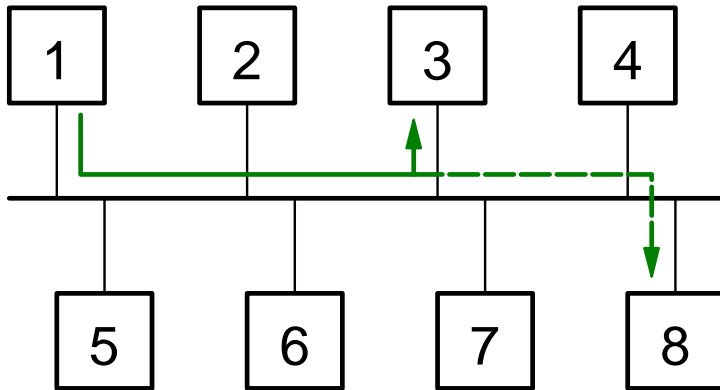
## Arbitration Overhead in Shared Media



- Separate medium for requests and grants?
  - increased media cost, increased latency.
- Shared medium for all of request, grant, and data?
  - reduced throughput, increased latency.
- Optimistic arbitration (CDMA/ethernet style) ?
  - limited peak throughput, very high latency at high loads.

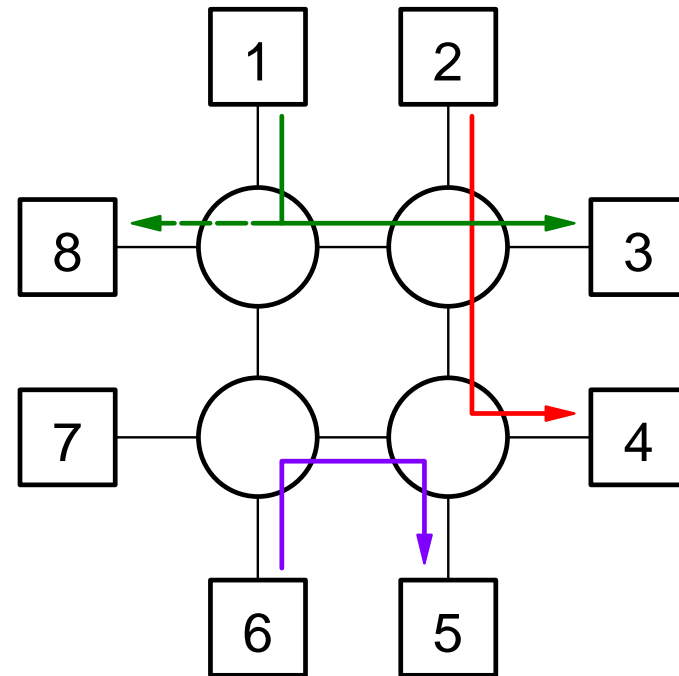
# Sequential versus Parallel Transmissions

Shared Medium:



Single transmission at a time

Point-to-Point Links + Switches:



Multiple transmissions in parallel

## This Course: Point-to-Point Links

- Throughput in shared media is rather low
  - time switching (simpler), protocols usually in software, with little, if any, & simple H/W support
- Shared media in modern networking: almost expelled, except in wireless
  - wireless MAC protocols are a major topic in other courses, and differ significantly from the hardware architecture of high-speed switches