



Fallacies of the Cost Based Optimizer

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Which Plan is better?

a)

<u>cost</u>	<u>card</u>	<u>operation</u>
2,979	446	SELECT STATEMENT
2,979	446	SORT ORDER BY
		FILTER
2,955	446	HASH JOIN
10	13,679	TABLE ACCESS FULL E
2,901	49,755	HASH JOIN
737	8,629	HASH JOIN
5	45	HASH JOIN
3	6	TABLE ACCESS FULL A
1	15	TABLE ACCESS FULL D
731	316,380	TABLE ACCESS FULL B
1,953	239,142	TABLE ACCESS FULL C

b)

<u>cost</u>	<u>card</u>	<u>operation</u>
792	1	SELECT STATEMENT
792	1	SORT ORDER BY
		FILTER
790	1	HASH JOIN
760	83	HASH JOIN
758	11	NESTED LOOPS
749	1	HASH JOIN
3	6	TABLE ACCESS FULL A
731	28,762	TABLE ACCESS FULL B
9	239,142	TABLE ACCESS BY INDEX ROWID C
4	239,142	INDEX RANGE SCAN C_IX0
1	15	TABLE ACCESS FULL D
10	13,679	TABLE ACCESS FULL E



Cost vs. Performance

Correlation between cost and performance?

Why not ?

Assumptions

- ❖ Uniform Distribution Assumption
 - ❖ Uniform Distribution over Blocks
 - ❖ Uniform Distribution over Rows
 - ❖ Uniform Distribution over Range of Values
- ❖ Predicate Independence Assumption
- ❖ Join Uniformity Assumption



Selectivity and Cardinality

$$\text{Selectivity} = \text{FF} = \text{card}_{\text{est}} / \text{card}_{\text{base}}$$

$$\text{card}_{\text{est}} = \text{FF} * \text{card}_{\text{base}}$$

The Makeup of Plan Costs

- ❖ The base table access cost is dependent on estimated # of blocks accessed which is - directly or indirectly - a function of the estimated row cardinality:
 - ❖ Table scan $nblks / k$
 - ❖ Unique scan $blevel + 1$
 - ❖ Fast full scan $leaf_blocks / k$
 - ❖ Index-only $blevel + FF * leaf_blocks$
 - ❖ Range scan $blevel + FF * leaf_blocks + FF * clustering_factor$

The Makeup of Plan Costs

Join cost is dependent on cardinality of row sources

- ❖ Nested Loop $\$_{outer} + \text{card}_{outer} * \$_{inner}$
- ❖ Sort-Merge $\$_{outer} + \$_{\text{sort}}_{outer} + \$_{inner} + \$_{\text{sort}}_{inner}$
- ❖ Hash $\$_{outer} + \$_{inner} + \$_{\text{hash}}$

Plan Costs Recap

Estimated cardinality = selectivity * base cardinality

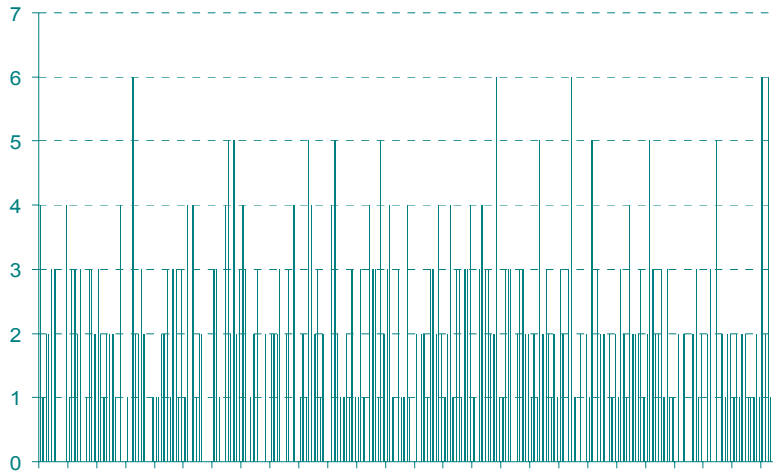
The cost of an access plan is a function of the estimated cardinalities of its components.

Incorrect estimates lead to incorrect plan component costs and sub-optimal or wrong access plans.

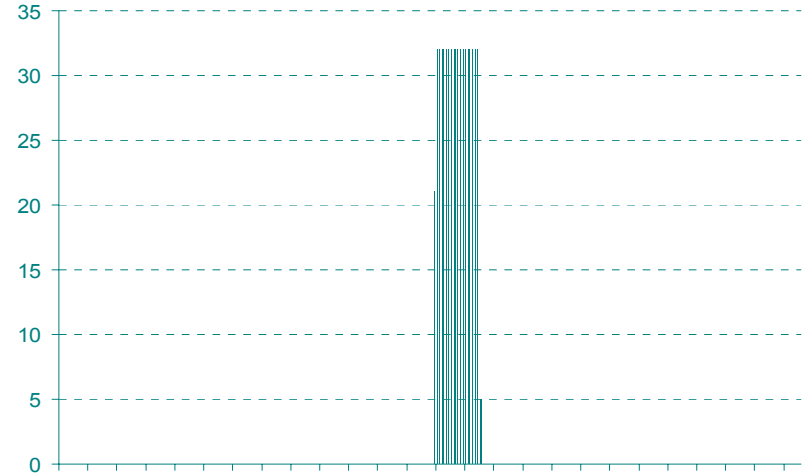
This is why accurate cardinality estimates are so important.

Distribution over blocks

uniform



clustered





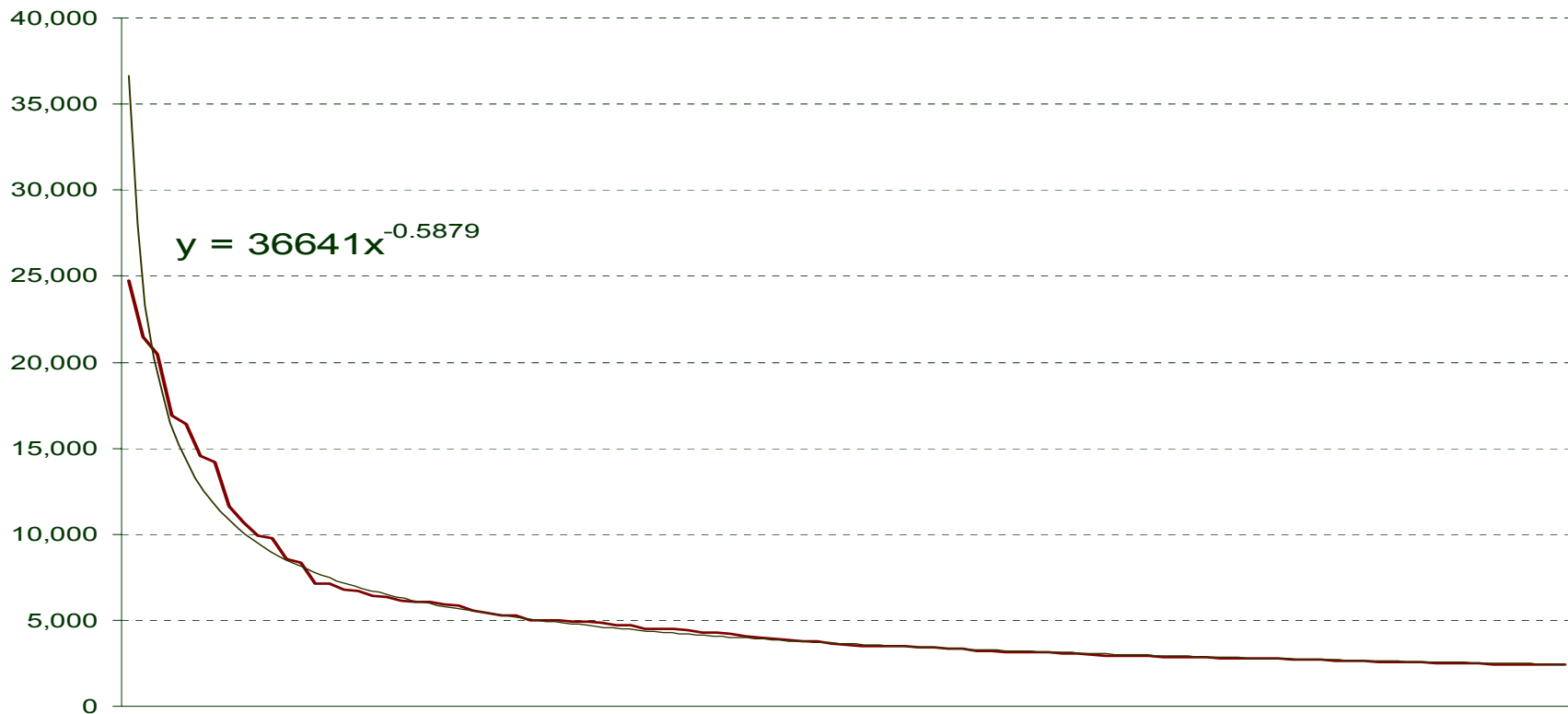
Distribution of Value Frequencies

“for an equality predicate (`last_name = 'Smith'`) the selectivity is set to the reciprocal of the number of distinct values of `last_name`, because the query selects rows that all contain one out of N distinct values.”*

* Oracle 9i/Performance Tuning Guide and Reference

Distribution of Value Frequencies

Power distribution



Distribution of Value Frequencies

column	NDV	density
-----	-----	-----
EMPLID	10,000	1.0000E-04
...		
COMPANY	200	5.0000E-03

```
Select emplid, jobcode, salary
from ps_job5 b where b.company = 'B01'
```

explain plan

card operation

```
-----
50 SELECT STATEMENT
50   TABLE ACCESS BY INDEX ROWID PS_JOB5
50     INDEX RANGE SCAN PSBJOB5
```

execution plan

card operation

```
-----
530 SELECT STATEMENT
530   TABLE ACCESS BY INDEX ROWID PS_JOB5
531     INDEX   GOAL: ANALYZED (RANGE SCAN) OF 'PSBJOB5' (NON-UNIQUE)
```

call	count	cpu	elapsed	disk	query	current	rows
-----	-----	-----	-----	-----	-----	-----	-----
Parse	1	0.47	0.47	21	359	5	0
Execute	1	0.00	0.00	0	0	0	0
Fetch	37	0.48	0.47	420	567	0	530
-----	-----	-----	-----	-----	-----	-----	-----
total	39	0.95	0.94	441	926	5	530

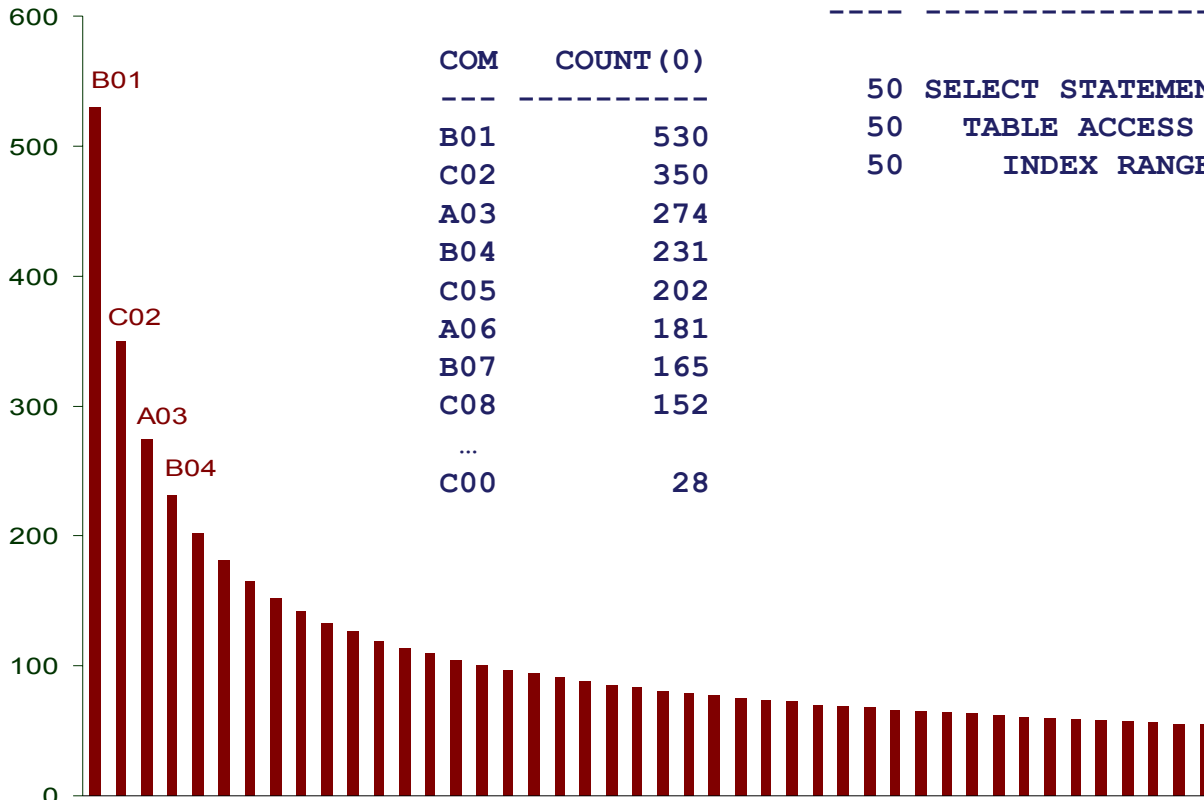
Distribution of Value Frequencies

column	NDV	density
EMPLID	10,000	1.0000E-04
...		
COMPANY	200	5.0000E-03

```
Select emplid, jobcode, salary
from ps_job5 b where b.company = 'B01'
```

explain plan

card operation



```
50 SELECT STATEMENT
50 TABLE ACCESS BY INDEX ROWID PS_JOB5
50 INDEX RANGE SCAN PSBJOB5
```

Distribution of Value Frequencies

With Histogram on company

```
Analyze table ps_job5 compute statistics for columns company [ size 75 ];
```

```
column          NDV          density          Select emplid, jobcode, salary
-----          -          -          from ps_job5 b where b.company = 'B01'
EMPLID          10,000      1.0000E-04
...
COMPANY          200         6.0644E-03
```

```
explain plan
```

```
card operation
```

```
-----
534 SELECT STATEMENT
534   TABLE ACCESS FULL PS_JOB5
```

```
execution plan
```

```
card operation
```

```
-----
530 SELECT STATEMENT GOAL: CHOOSE
530   TABLE ACCESS GOAL: ANALYZED (FULL) OF 'PS_JOB5'
```

call	count	cpu	elapsed	disk	query	current	rows
Parse	1	0.17	0.15	25	424	0	0
Execute	1	0.00	0.00	0	0	0	0
Fetch	37	0.24	0.22	912	943	15	530
total	39	0.41	0.37	937	1367	15	530

Distribution of Value Frequencies

With Histogram and bind Variable on company

column	NDV	density
-----	-----	-----
EMPLID	10,000	1.0000E-04
...		
COMPANY	200	6.0644E-03

```
Select emplid, jobcode, salary  
from ps_job5 b where b.company = :b1
```

explain plan

card operation

```
-----  
61 SELECT STATEMENT  
61 TABLE ACCESS BY INDEX ROWID PS_JOB5  
61 INDEX RANGE SCAN PSBJOB5
```

$$10,000 * 6.0644e^{-3} = 60.644 \text{ rounded up to } 61.$$

Distribution of Value Frequencies

With Histogram and bind Variable on company

```
Analyze table ps_job5 compute statistics for columns company size 10;
```

column	NDV	density
-----	-----	-----
EMPLID	10,000	1.0000E-04
...		
COMPANY	200	1.0870E-02

```
Select emplid, jobcode, salary  
from ps_job5 b where b.company = :b1
```

```
explain plan
```

```
card operation
```

```
-----  
109 SELECT STATEMENT  
109  TABLE ACCESS FULL PS_JOB5
```


Column Statistics and Histograms

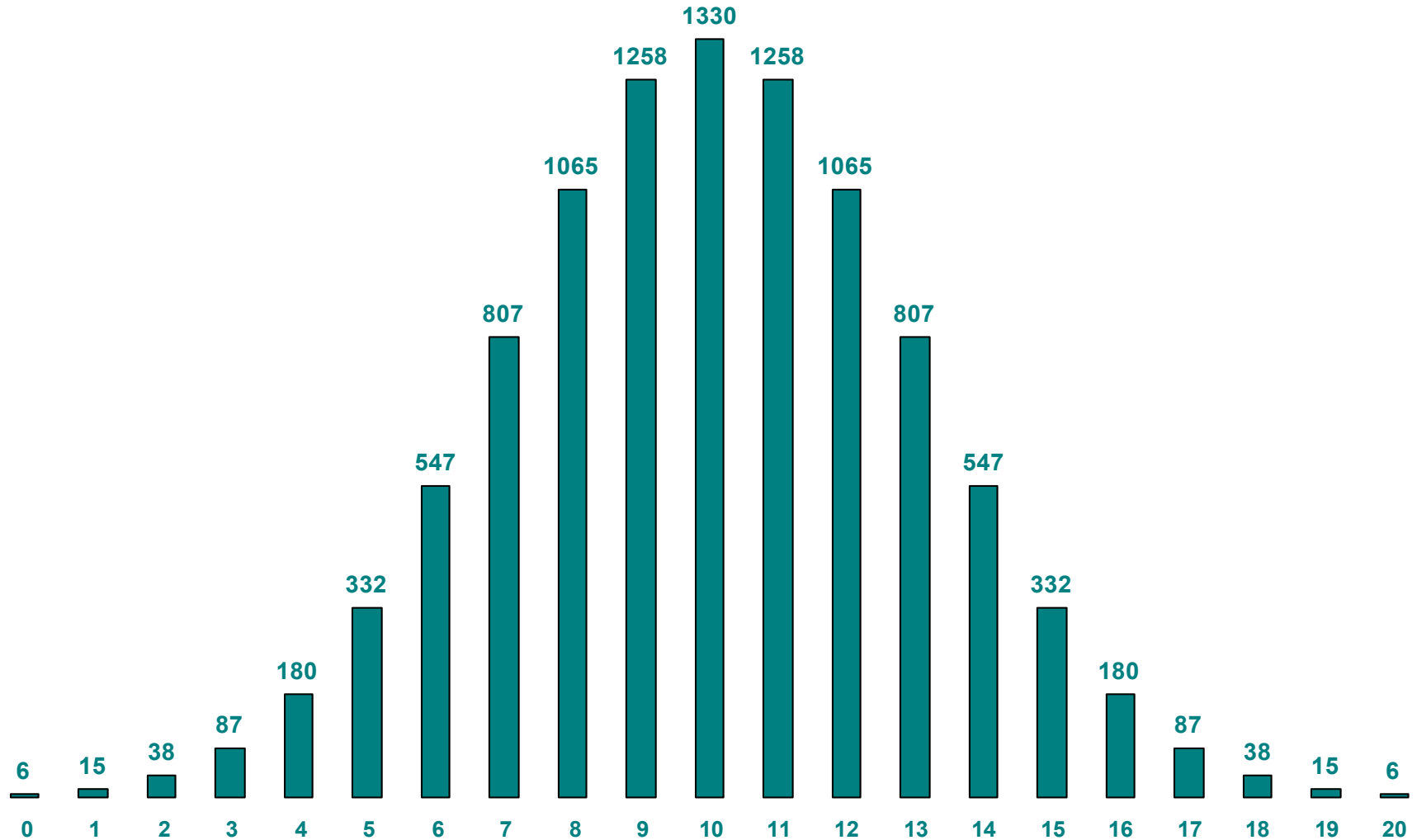
❖ Frequency Histogram

buckets = NDV

❖ Height Balanced Histogram

buckets < NDV

Histograms

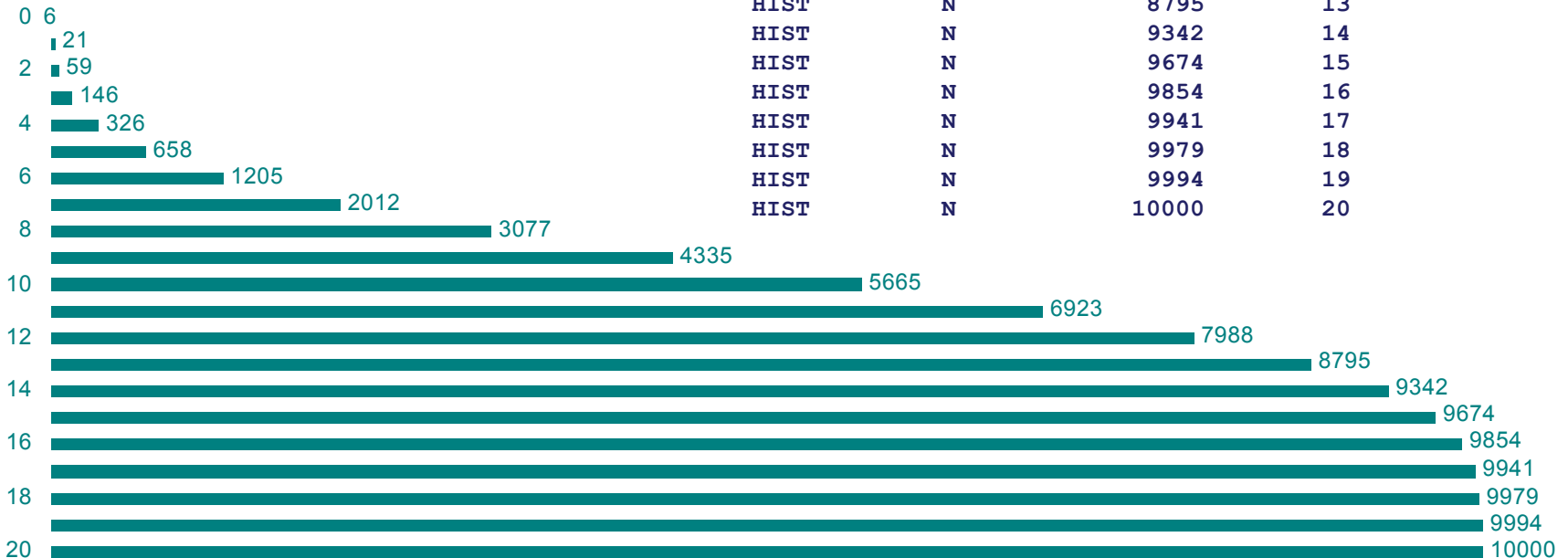


Frequency Histogram

analyze table hist compute statistics for columns n [size 75];

<u>table</u>	<u>column</u>	<u>NDV</u>	<u>density</u>	<u>lo</u>	<u>hi</u>	<u>bkts</u>
HIST	N	21	5.0000E-05	0	20	20

<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>	<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>
HIST	N	6	0	HIST	N	1205	6
HIST	N	21	1	HIST	N	2012	7
HIST	N	59	2	HIST	N	3077	8
HIST	N	146	3	HIST	N	4335	9
HIST	N	326	4	HIST	N	5665	10
HIST	N	658	5	HIST	N	6923	11
				HIST	N	7988	12
				HIST	N	8795	13
				HIST	N	9342	14
				HIST	N	9674	15
				HIST	N	9854	16
				HIST	N	9941	17
				HIST	N	9979	18
				HIST	N	9994	19
				HIST	N	10000	20



Frequency Histogram

- ❖ Predicate matches one of the values in the histogram

$$\text{selectivity} = \begin{matrix} < & & \leq & & = \\ \text{EP of prior row} & & \text{EP of matching row} & & \text{difference} \\ / \text{num_rows} & & / \text{num_rows} & & / \text{num_rows} \end{matrix}$$

Example:

<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>
HIST	N	1205	6
HIST	N	2012	7

$$N < 7 \quad \text{selectivity} = 1205 / \text{num_rows}$$

$$N \leq 7 \quad \text{selectivity} = 2012 / \text{num_rows}$$

$$N = 7 \quad \text{selectivity} = (2012 - 1205) / \text{num_rows}$$

Frequency Histogram

❖ Predicate does not match one of the values in the histogram

Since this is a value base histogram that should mean there are no rows in the table with that value for the column and therefore the selectivity should be 0.

However, the optimizer can not rely on the statistics being up-to-date and uses the density from the column statistics as selectivity.

❖ Bind Variable predicate

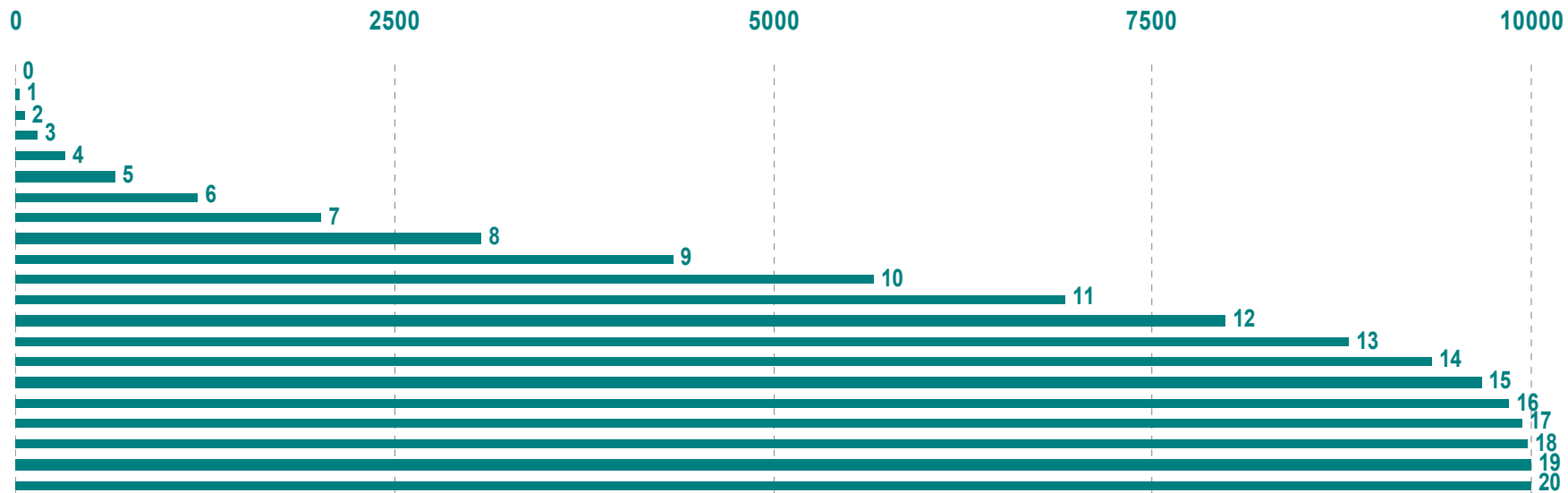
The selectivity is taken as $\max(1/\text{num_distinct}, \text{density})$, effectively ignoring the histogram.

Height Balanced Histogram

```
analyze table hist compute statistics for columns n size 4;
```

<u>table</u>	<u>column</u>	<u>NDV</u>	<u>density</u>	<u>lo</u>	<u>hi</u>	<u>bkts</u>
HIST	N	21	9.4128E-02	0	20	4

<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>
HIST	N	0	0
HIST	N	1	8
HIST	N	2	10
HIST	N	3	12
HIST	N	4	20

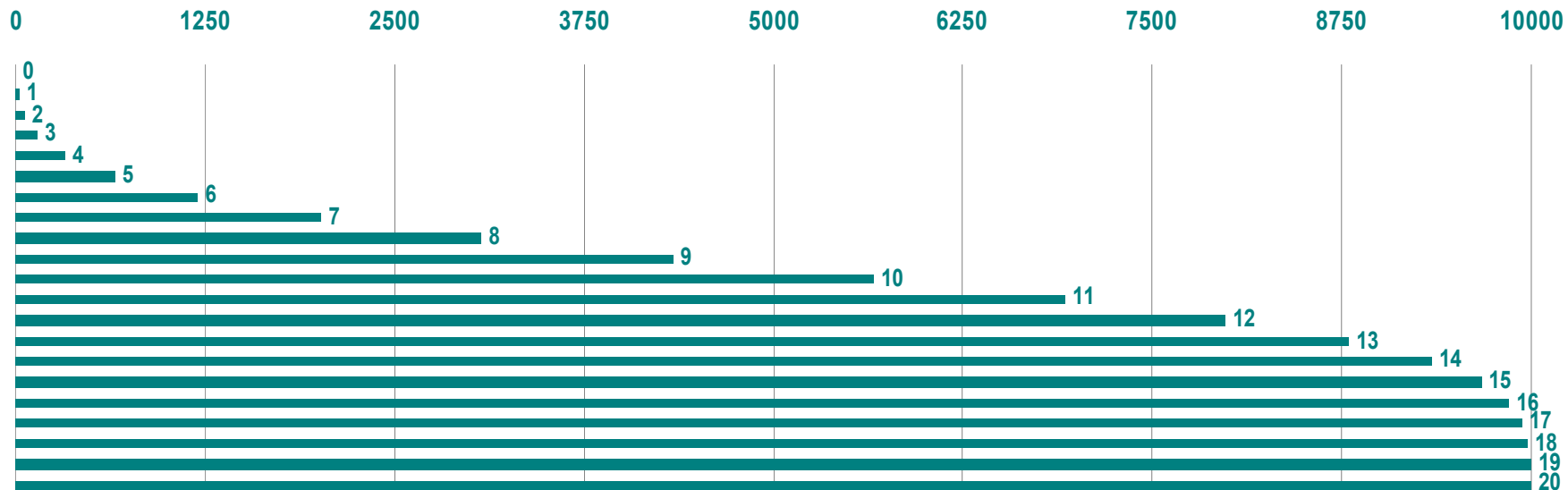


Height Balanced Histogram

```
analyze table hist compute statistics for columns n size 8;
```

<u>table</u>	<u>column</u>	<u>NDV</u>	<u>density</u>	<u>lo</u>	<u>hi</u>	<u>bkts</u>
HIST	N	21	6.2500E-02	0	20	8

<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>
HIST	N	0	0
HIST	N	1	7
HIST	N	2	8
HIST	N	3	9
HIST	N	4	10
HIST	N	5	11
HIST	N	6	12
HIST	N	7	13
HIST	N	8	20

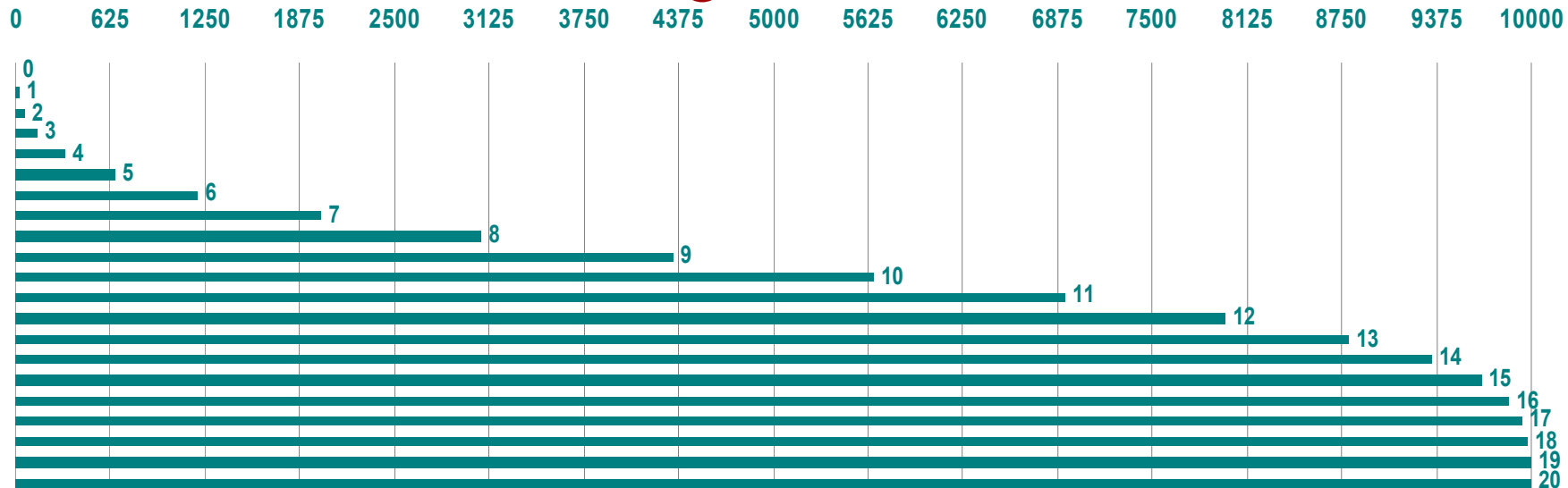


Height Balanced Histogram

analyze table hist compute statistics for columns n size 16

<u>table</u>	<u>column</u>	<u>NDV</u>	<u>density</u>	<u>lo</u>	<u>hi</u>	<u>bkts</u>
HIST	N	21	3.1250E-02	0	20	10

<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>
HIST	N	0	0
HIST	N	1	5
HIST	N	3	7
HIST	N	4	8
HIST	N	6	9
HIST	N	9	10
HIST	N	11	11
HIST	N	12	12
HIST	N	14	13
HIST	N	15	15
HIST	N	16	20



Height Balanced Histogram

Predicate does not match any of the values in the histogram

$< | <=$ $=$

$$\text{selectivity} = \frac{\text{EP of prior row}}{\text{buckets}} + \frac{(\text{value} - \text{value}_{\text{low}})}{(\text{value}_{\text{hi}} - \text{value}_{\text{low}})} \times \text{density}$$

Example:

<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>	<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>
HIST	N	0	0	HIST	N	11	11
HIST	N	1	5	HIST	N	12	12
HIST	N	3	7	HIST	N	14	13
HIST	N	4	8	HIST	N	15	15
HIST	N	6	9	HIST	N	16	20
HIST	N	9	10				

$N < 17$ selectivity = $15 / 16 + (17-15) / (20-15) / 16 = 9.6250E-01$

$N = 17$ selectivity = density = $3.1250E-02$

Height Balanced Histogram

Predicate matches a “non-popular” value in the histogram

< | <=

=

selectivity =

EP of prior row / buckets

density

Example:

<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>	<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>
HIST	N	0	0	HIST	N	11	11
HIST	N	1	5	HIST	N	12	12
HIST	N	3	7	HIST	N	14	13
HIST	N	4	8	HIST	N	15	15
HIST	N	6	9	HIST	N	16	20
HIST	N	9	10				

N < 15 selectivity = 14 / 16 = 8.7500E-01

N = 15 selectivity = density = 3.1250E-02

Height Balanced Histogram

Predicate matches a "popular" value in the histogram

$< \mid \leq$ $=$

selectivity = EP of prior row / buckets range / buckets

Example:

<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>	<u>table</u>	<u>column</u>	<u>EP</u>	<u>value</u>
HIST	N	0	0	HIST	N	11	11
HIST	N	1	5	HIST	N	12	12
HIST	N	3	7	HIST	N	14	13
HIST	N	4	8	HIST	N	15	15
HIST	N	6	9	HIST	N	16	20
HIST	N	9	10				

$N < 13$ selectivity = $12 / 16 = 7.5000E-01$

$N = 13$ selectivity = $2 / 16 = 1.2500E-02$

Histograms and Bind Variables

Density and cardinality estimate by # of buckets

```
Select emplid, jobcode, salary from ps_job5 b where b.company = :b1
```

<u>buckets</u>	<u>density</u>	<u>card</u>
10	1.0870E-02	109
25	8.5039E-03	86
50	7.4833E-03	75
75	6.0644E-03	61
90	5.5556E-03	56
100	5.0000E-03	50
150	3.3333E-03	50
199	2.5381E-03	50
200	5.0000E-05	50



Distribution over Range of Values

“The optimizer assumes that `employee_id` values are distributed evenly in the range between the lowest value and highest value.”*

* Oracle 9i Performance Tuning Guide and Reference

Distribution over Range of Values

<u>table</u>	<u>column</u>	<u>NDV</u>	<u>density</u>	<u>lo</u>	<u>hi</u>
PS_LEDGER	ACCOUNTING_PERIOD	15	6.6667E-02	0	999

Period 0 holds opening balances, periods 1-12 hold the ledger entries for the months, and periods 998 and 999 are used for special processing.

Distribution over Range of Values

<u>table</u>	<u>column</u>	<u>NDV</u>	<u>density</u>	<u>lo</u>	<u>hi</u>
PS_LEDGER	ACCOUNTING_PERIOD	15	6.6667E-02	0	999

accounting_period = n [n \in {1 .. 12}]

$$\Rightarrow \text{selectivity} = 1/\text{ndv} = 1/15 = 6.6667e^{-2}$$

accounting_period between 1 and 12

$$\Rightarrow \text{selectivity} = 12/(999-0) + 1/15 = 7.8679e^{-2}$$

accounting_period < 12

$$\Rightarrow \text{selectivity} = (12-0)/(999-0) = 1.2012e^{-2}$$

Distribution over Range of Values

Adjusting the high-value statistic

```
select sum(posted_total_amt) from ps_ledger
where accounting_period between 1 and 12
```

```
Column: ACCOUNTING Col#: 11 Table: PS_LEDGER Alias: PS_LEDGER
NDV: 15 NULLS: 0 DENS: 6.6667e-002 LO: 0 HI: 999
TABLE: PS_LEDGER ORIG CDN: 745198 CMPTD CDN: 58632
```

```
Column: ACCOUNTING Col#: 11 Table: PS_LEDGER Alias: PS_LEDGER
NDV: 15 NULLS: 0 DENS: 6.6667e-002 LO: 0 HI: 14
TABLE: PS_LEDGER ORIG CDN: 745198 CMPTD CDN: 684873
```

```
select sum(posted_total_amt) from ps_ledger
where accounting_period < 12
```

```
Column: ACCOUNTING Col#: 11 Table: PS_LEDGER Alias: PS_LEDGER
NDV: 15 NULLS: 0 DENS: 6.6667e-002 LO: 0 HI: 999
TABLE: PS_LEDGER ORIG CDN: 745198 CMPTD CDN: 49680
```

```
Column: ACCOUNTING Col#: 11 Table: PS_LEDGER Alias: PS_LEDGER
NDV: 15 NULLS: 0 DENS: 6.6667e-002 LO: 0 HI: 14
TABLE: PS_LEDGER ORIG CDN: 745198 CMPTD CDN: 638742
```


Predicate Independence Assumption

P1 AND P2 $S(P1 \& P2) = S(P1) * S(P2)$

P1 OR P2 $S(P1 | P2) = S(P1) + S(P2) - [S(P1) * S(P2)]$

```
select emplid, jobcode, salary
from ps_job1 b
where b.company = 'CCC'
      and b.paygroup = 'FGH';
```

250 rows selected.

Explain Plan

```
card operation
251 SELECT STATEMENT
251 TABLE ACCESS BY INDEX ROWID PS_JOB1
251 INDEX RANGE SCAN PSJOB1
```

```
select emplid, jobcode, salary
from ps_job2 b
where b.company = 'CCC'
      and b.paygroup = 'FGH';
```

2500 rows selected.

Explain Plan

```
card operation
251 SELECT STATEMENT
251 TABLE ACCESS BY INDEX ROWID PS_JOB2
251 INDEX RANGE SCAN PSJOB2
```

Predicate Independence Assumption

table	rows	blks	empty	chain	avg rl	table	rows	blks	empty	chain	avg rl
PS_JOB1	50,000	4,547	3	0	317	PS_JOB2	50,000	4,547	3	0	317

table	column	NDV	density	bkts	table	column	NDV	density	bkts
PS_JOB1	EMPLID	10,000	1.0000E-04	1	PS_JOB2	EMPLID	10,000	1.0000E-04	1
PS_JOB1	JOBCODE	198	5.0505E-03	1	PS_JOB2	JOBCODE	199	5.0251E-03	1
PS_JOB1	COMPANY	10	1.0000E-01	1	PS_JOB2	COMPANY	10	1.0000E-01	1
PS_JOB1	PAYGROUP	20	5.0000E-02	1	PS_JOB2	PAYGROUP	20	5.0000E-02	1
PS_JOB1	SALARY	49,597	2.0163E-05	1	PS_JOB2	SALARY	49,848	2.0061E-05	1

$$\begin{aligned}
 \text{card}_{\text{est}} &= \text{card}_{\text{base}} * \text{sel}(\text{company AND paygroup}) \\
 &= \text{sel}(\text{company}) * \text{sel}(\text{paygroup}) \\
 &= 50000 * 1.0000e^{-01} * 5.0000e^{-02} = 250
 \end{aligned}$$

index	column	NDV	#LB	index	column	NDV	#LB
PSBJOB1		200	400	PSBJOB2		20	449
	COMPANY	10			COMPANY	10	
	PAYGROUP	20			PAYGROUP	20	

Join Uniformity Assumption

join cardinality = $\text{card}_A * \text{card}_B * \text{join selectivity}$

join selectivity = $1/\max(\text{ndv}_A, \text{ndv}_B)$

“principle of inclusion”, i.e. each value of the smaller domain has a match in the larger domain – which is frequently true for joins between foreign keys and primary keys.

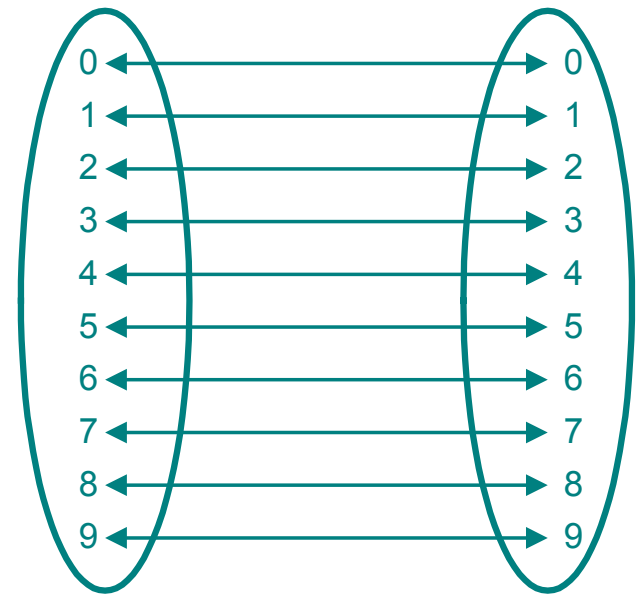
Join Uniformity Assumption

```
SQL> select 'A-'||a.n1, 'B-'||b.n1  
2 from t1 a, t1 b  
3 where a.n1 = b.n1;
```

```
10 SELECT STATEMENT  
10 HASH JOIN  
10 TABLE ACCESS FULL T1  
10 TABLE ACCESS FULL T1
```

```
A-0 B-0  
A-1 B-1  
A-2 B-2  
A-3 B-3  
A-4 B-4  
A-5 B-5  
A-6 B-6  
A-7 B-7  
A-8 B-8  
A-9 B-9
```

10 rows selected.



$$\begin{aligned}\text{Join cardinality} &= \text{card}_A * \text{card}_B * \text{join selectivity} \\ &= \text{card}_A * \text{card}_B * 1/\max(\text{ndv}_a, \text{ndv}_b) \\ &= 10 * 10 * 1/\max(10, 10) = 10\end{aligned}$$

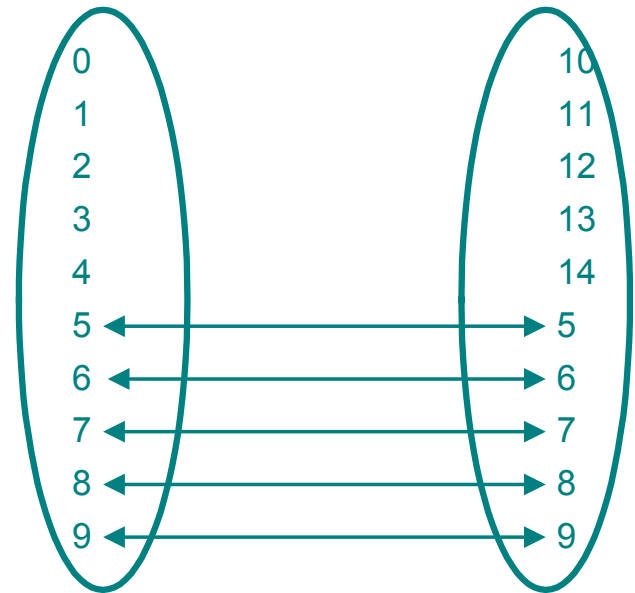
Join Uniformity Assumption

```
SQL> select 'A-'||a.n1, 'B-'||b.n1  
2 from t1 a, t2 b  
3 where a.n1 = b.n1;
```

```
10 SELECT STATEMENT  
10 HASH JOIN  
10 TABLE ACCESS FULL  
T2  
10 TABLE ACCESS FULL  
T2
```

```
A-5 B-5  
A-6 B-6  
A-7 B-7  
A-8 B-8  
A-9 B-9
```

5 rows selected.



$$\begin{aligned}\text{Join cardinality} &= \text{card}_A * \text{card}_B * \text{join selectivity} \\ &= \text{card}_A * \text{card}_B * 1/\max(\text{ndv}_a, \text{ndv}_b) \\ &= 10 * 10 * 1/\max(10, 10) = 10\end{aligned}$$

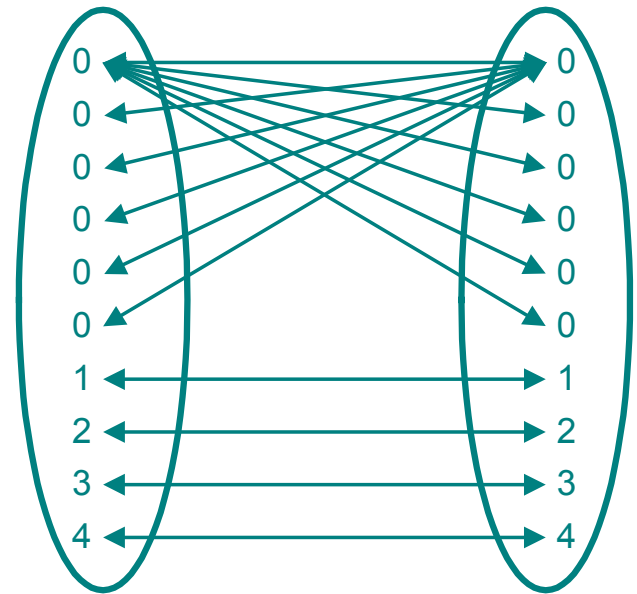
Join Uniformity Assumption

```
SQL> select 'A-'||a.n1, 'B-'||b.n1  
2 from t2 a, t2 b  
3 where a.n1 = b.n1;
```

```
20 SELECT STATEMENT  
20 HASH JOIN  
10 TABLE ACCESS FULL T2  
10 TABLE ACCESS FULL T2
```

```
A-0 B-0  
A-0 B-0  
A-0 B-0  
...  
A-0 B-0  
A-1 B-1  
A-2 B-2  
A-3 B-3  
A-4 B-4
```

40 rows selected.



$$\begin{aligned}\text{Join cardinality} &= \text{card}_A * \text{card}_B * \text{join selectivity} \\ &= \text{card}_A * \text{card}_B * 1/\max(\text{ndv}_a, \text{ndv}_b) \\ &= 10 * 10 * 1/\max(5, 5) = 20\end{aligned}$$

Join Selectivity and Cardinality

```
insert into t1(n1,n2)
select mod(rownum,10),mod(rownum,5)
from dba_objects where rownum <= 50;
```

<u>column</u>	<u>NDV</u>	<u>density</u>
N1	10	1.0000E-01
N2	5	2.0000E-01

```
select 'A.'||A.n1||'-B.'||B.n1
from t1 a, t2 b
where a.n1 = b.n1;
```

Explain Plan

<u>card</u>	<u>operation</u>
250	SELECT STATEMENT
250	HASH JOIN
50	TABLE ACCESS FULL T1
50	TABLE ACCESS FULL T2

Execution Plan

<u>Rows</u>	<u>Execution Plan</u>
0	SELECT STATEMENT GOAL: CHOOSE
250	HASH JOIN
50	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T1'
50	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T2'

Join Selectivity and Cardinality

```
select 'A.' || A.n1 || '-B.' || B.n1
from t1 a, t2 b
where a.n1 = b.n1
      and a.n2 = 5;
```

Explain Plan

<u>card</u>	<u>operation</u>
50	SELECT STATEMENT
50	HASH JOIN
10	TABLE ACCESS FULL T1
50	TABLE ACCESS FULL T2

Execution Plan

<u>Rows</u>	<u>Execution Plan</u>
0	SELECT STATEMENT GOAL: CHOOSE
50	HASH JOIN
10	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T1'
50	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T2'

Join Selectivity and Cardinality

```
select 'A.' || A.n1 || '-B.' || B.n1
from t1 a, t2 b
where a.n1 = b.n1
      and a.n1 = 5;
```

Explain Plan

<u>card</u>	<u>operation</u>
5	SELECT STATEMENT
5	HASH JOIN
5	TABLE ACCESS FULL T1
5	TABLE ACCESS FULL T2

Execution Plan

<u>Rows</u>	<u>Execution Plan</u>
0	SELECT STATEMENT GOAL: CHOOSE
25	HASH JOIN
5	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T1'
5	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T2'

Join Selectivity and Cardinality

```
Table stats      Table: T2      Alias: B
TOTAL :: CDN: 50  NBLKS: 1  TABLE_SCAN_CST: 1  AVG_ROW_LEN: 8
```

```
Table stats      Table: T1      Alias: A
TOTAL :: CDN: 50  NBLKS: 1  TABLE_SCAN_CST: 1  AVG_ROW_LEN: 8
```

SINGLE TABLE ACCESS PATH

```
Column:          N1  Col#: 1          Table: T1      Alias: A
NDV: 10          NULLS: 0          DENS: 1.0000e-001 LO: 0  HI: 9
TABLE: T1        ORIG CDN: 50  CMPTD CDN: 5
```

SINGLE TABLE ACCESS PATH

```
Column:          N1  Col#: 1          Table: T2      Alias: B
NDV: 10          NULLS: 0          DENS: 1.0000e-001 LO: 0  HI: 9
TABLE: T2        ORIG CDN: 50  CMPTD CDN: 5
```

...

```
Join cardinality: 5 = outer (5) * inner (5) * sel (2.0000e-001) [flag=0]
```

Transitive Closure

```
select
a.n1,a.n2,a.n3,b.n1,b.n2,b.n3,c.n1,c.n2,c.n3
from t4 a, t5 b, t6 c
where a.n1 = b.n1
      and b.n2 = c.n2
      and b.n1 = c.n1
```

cost	card	operation	Rows	Execution Plan
23	198	SELECT STATEMENT	0	SELECT STATEMENT GOAL: CHOOSE
23	198	HASH JOIN	202	HASH JOIN
6	100	HASH JOIN	100	HASH JOIN
1	20	TABLE ACCESS FULL T4	20	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T4'
4	100	TABLE ACCESS FULL T5	100	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T5'
16	500	TABLE ACCESS FULL T6	500	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T6'

```
select
a.n1,a.n2,a.n3,b.n1,b.n2,b.n3,c.n1,c.n2,c.n3
from t4 a, t5 b, t6 c
where a.n1 = b.n1
      and b.n2 = c.n2
      and b.n1 = c.n1
      and a.n1 = c.n1
```

cost	card	operation	Rows	Execution Plan
23	18	SELECT STATEMENT	0	SELECT STATEMENT GOAL: CHOOSE
23	18	HASH JOIN	202	HASH JOIN
6	100	HASH JOIN	100	HASH JOIN
1	20	TABLE ACCESS FULL T4	20	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T4'
4	100	TABLE ACCESS FULL T5	100	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T5'
16	500	TABLE ACCESS FULL T6	500	TABLE ACCESS GOAL: ANALYZED (FULL) OF 'T6'

Which Plan is better?

a)

<u>cost</u>	<u>card</u>	<u>operation</u>
2,979	446	SELECT STATEMENT
2,979	446	SORT ORDER BY
		FILTER
2,955	446	HASH JOIN
10	13,679	TABLE ACCESS FULL E
2,901	49,755	HASH JOIN
737	8,629	HASH JOIN
5	45	HASH JOIN
3	6	TABLE ACCESS FULL A
1	15	TABLE ACCESS FULL D
731	316,380	TABLE ACCESS FULL B
1,953	239,142	TABLE ACCESS FULL C

b)

<u>cost</u>	<u>card</u>	<u>operation</u>
792	1	SELECT STATEMENT
792	1	SORT ORDER BY
		FILTER
790	1	HASH JOIN
760	83	HASH JOIN
758	11	NESTED LOOPS
749	1	HASH JOIN
3	6	TABLE ACCESS FULL A
731	28,762	TABLE ACCESS FULL B
9	239,142	TABLE ACCESS BY INDEX ROWID C
4	239,142	INDEX RANGE SCAN C_IX0
1	15	TABLE ACCESS FULL D
10	13,679	TABLE ACCESS FULL E

Analysis of the Explain Plan

<u>cost</u>	<u>card</u>	<u>operation</u>
792	1	SELECT STATEMENT
792	1	SORT ORDER BY
		FILTER
790	1	HASH JOIN
760	83	HASH JOIN
758	11	NESTED LOOPS
749	1	HASH JOIN
3	6	TABLE ACCESS FULL A
731	28,762	TABLE ACCESS FULL B
9	239,142	TABLE ACCESS BY INDEX ROWID C
4	239,142	INDEX RANGE SCAN C_IX0
1	15	TABLE ACCESS FULL D
10	13,679	TABLE ACCESS FULL E



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