

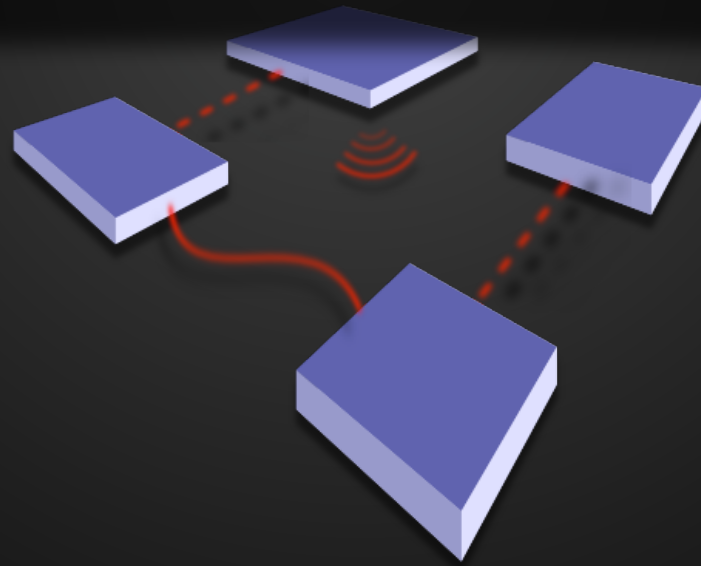
CS-435

spring semester 2020

Network Technology & Programming Laboratory

University of Crete
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CS-435

Lecture #13 preview

- 802.11e
- WEP, WPA
- 802.11i

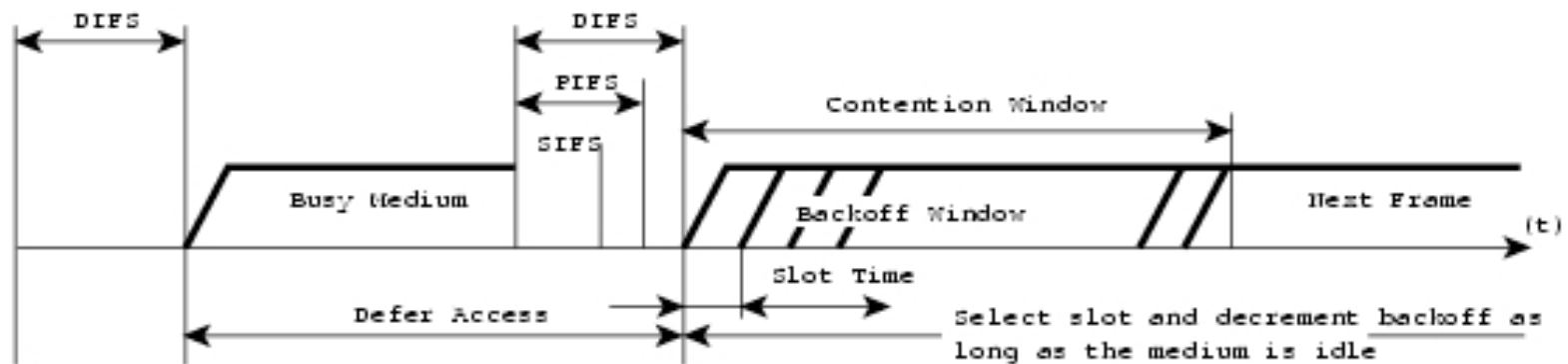
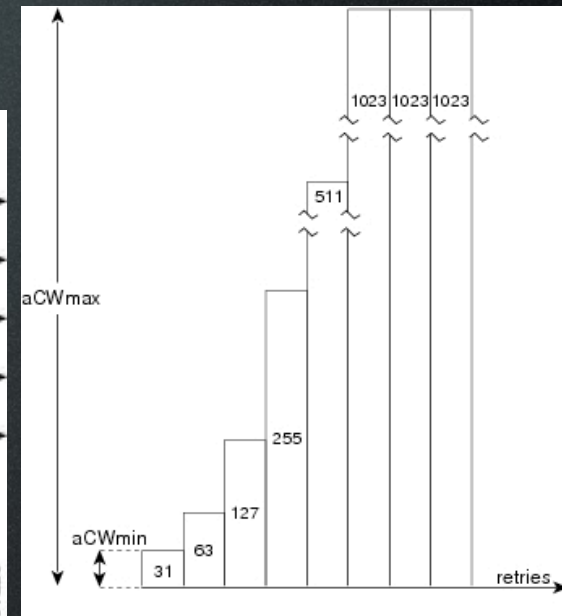
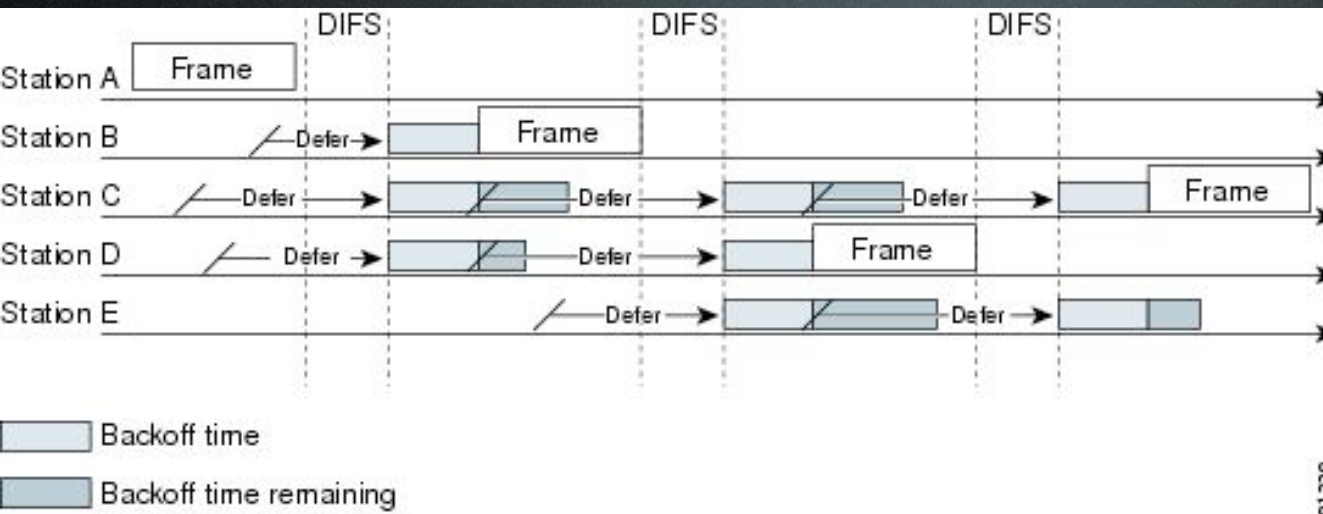
802.11 access methods

- **DCF** (Distributed Coordination Function)
- **PCF** (Point Coordination Function)
 - Time divided into Contention Period (**CP**) and Contention Free period (**CFP**)
 - During **CP**, transfers use **DCF**, i.e., Data-ACK, or RTS-CTS-Data-ACK, with exponential back-off etc.
 - During **CFP**, the AP controls all transmissions: which STA transmits to the AP and which STA receives data from the AP.
 - All STAs may transmit/receive packets during the CFP. The ability to transmit during the CFP is optional.
 - No RTS/CTS in CFP

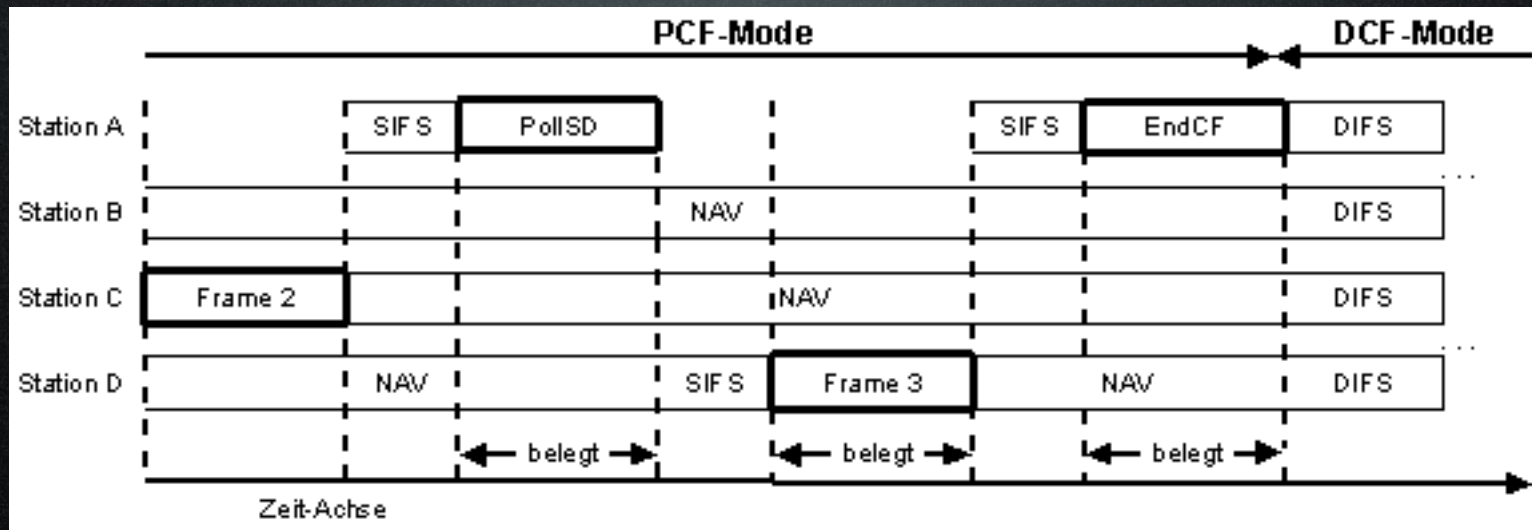
802.11 access methods

- **DCF** (Distributed Coordination Function)
 - DIFS: DCF Interframe space
 - SIFS: Shortest Interframe space
 - $\text{DIFS} = \text{SIFS} + 2 * \text{SlotTime}$
- **PCF** (Point Coordination Function)
 - PIFS: PCF Interframe Space

DCF



PCF



Polling during CFP

- STAs can only transmit data one SIFS after being polled & only one frame
- If the pollable STA transmits a frame in response to a CF-poll but does not receive an ACK, then it cannot retransmit until it is polled again or until the CP.
- A STA can set the MoreData bit is set
 - The more data bit is in the MAC header. The STA can only set it when responding to a CF-poll request.
- A STA should respond to a poll within SIFS. If not, the AP will regain control PIFS after the CF-poll was sent.
- If a STA has no data to respond with but the CF-poll had data (Data+CF-poll or Data+CF-ACK+CF-poll), then the STA responds with CF-ACK, with no data
- If the STA is polled without data (i.e., CF-poll or CF-ACK + CF-poll), then it should sent a null frame (no data). This ensures that the AP does not think that that STA missed the due to radio transmission error.

QoS Limitations of 802.11

- DCF
 - Only supports best-effort services
 - No guarantee in bandwidth, packet delay and jitter
- PCF
 - Unpredictable beacon frame delay due to incompatible cooperation between CP and CFP modes
 - Transmission time of the polled stations is unknown
 - The Point Coordinator (PC) does not know the QoS requirement of traffic

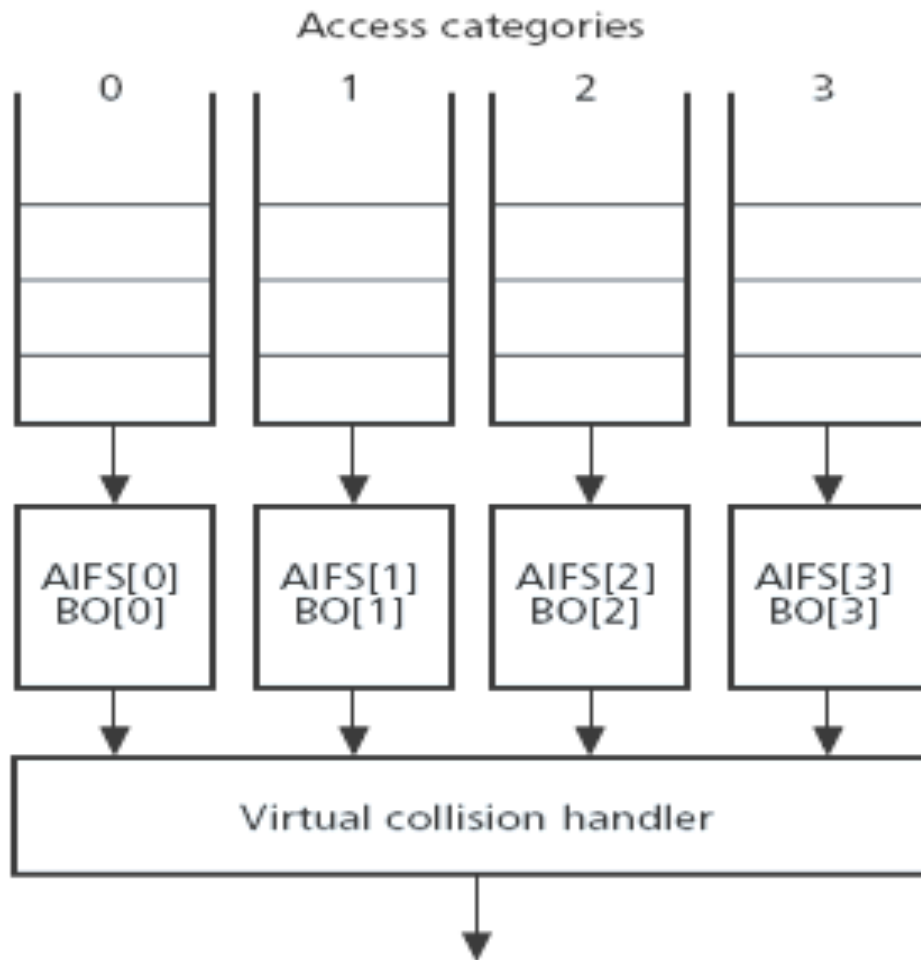
Overview of 802.11e

- Support QoS in WLAN
- Backwardly compatible with the DCF and PCF
- Hybrid Coordination Function (HCF) access method is added, including
 - Contention-Based channel access
 - Enhanced Distributed Channel Access (EDCA)
 - Controlled channel access
 - HCF Controlled Channel Access (HCCA)

Major Enhancements in 802.11e

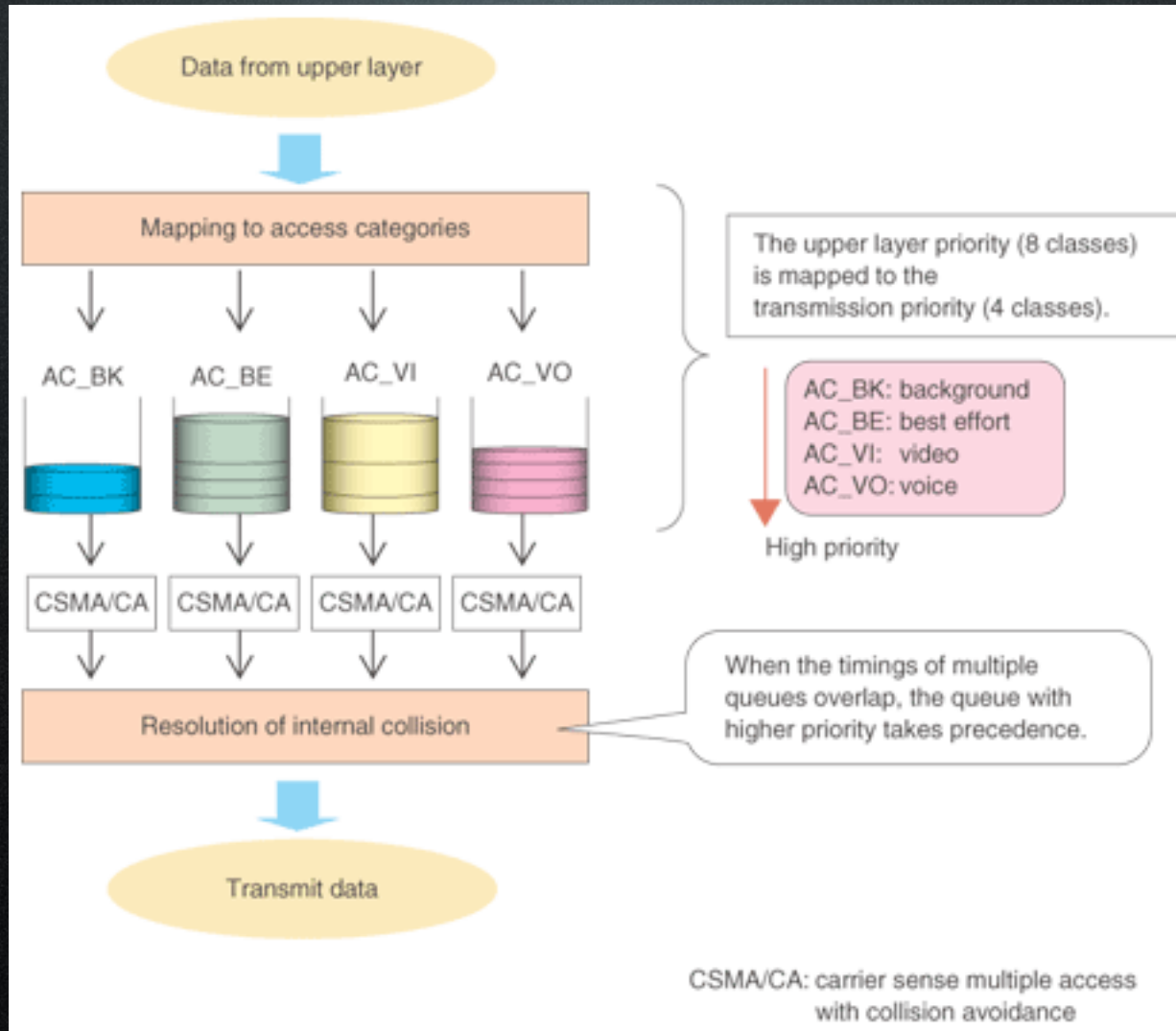
- Basic elements for QoS
 - Traffic Differentiation
 - Concept of Transmission Opportunity (TXOP)
- New Contention-based channel access
 - Enhanced Distributed Channel Access (EDCA)
- New Contention-free channel access
 - HCF Controlled Channel Access (HCCA)
- Other new mechanisms for higher throughput
 - Block Acknowledgement (Block Ack)
 - Direct Link Protocol (DLP)

Traffic Differentiation



Priority	AC	Designation
1	0	Best effort
2	0	Best effort
0	0	Best effort
3	1	Video probe
4	2	Video
5	2	Video
6	3	Voice
7	3	Voice

Traffic Differentiation



EDCA

Difference from original DCF

- Contention between ACs (Not STAs)
- New Inter-frame Space (IFS) for each AC:
Arbitration Inter frame Space (AIFS)
- Transmission Opportunity (TXOP)

Access Category (AC)

- In EDCA, media access is based on the AC of MSDU
- 4 AC's defined:
 - AC_BK (background)
 - AC_BE (best-effort)
 - AC_VI (video)
 - AC_VO (voice)
- In EDCA, the size of Contention-Window (CW) and Inter-frame space (IFS) is AC dependent

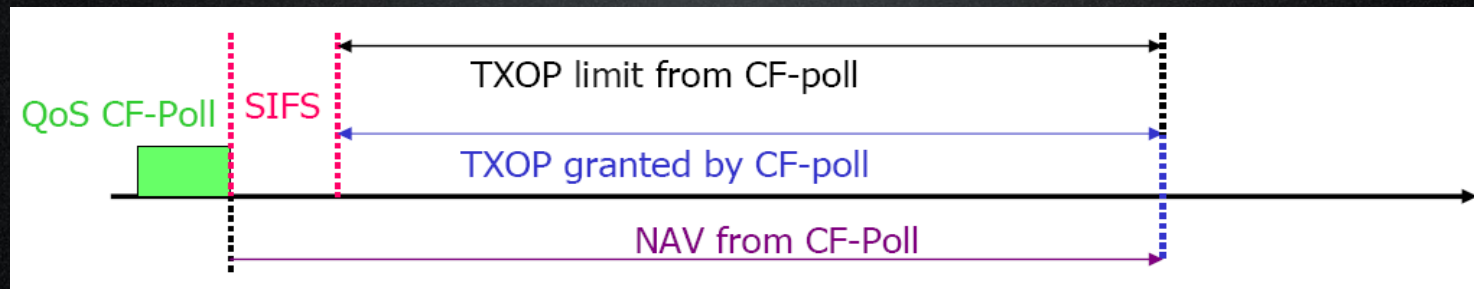
Arbitration Inter-frame Space (AIFS)

- QSTAs use AIFS to defer the contention window or transmission for each AC
- $AIF[AC] = AIFSN[AC] \times aSlotTime + aSIFSTime$
 - AIFSN for each AC is broadcast via beacon frame containing “EDCA Parameter Set” element
- $DIFS = 2 \times aSlotTime + aSIFSTime$

AC	CWmin	CWmax
AC_BK	aCWmin	aCWmax
AC_BE	aCWmin	aCWmax
AC_VI	$(aCWmin+1)/2-1$	aCWmin
AC_VO	$(aCWmin+1)/4-1$	$(aCWmin+1)/2-1$

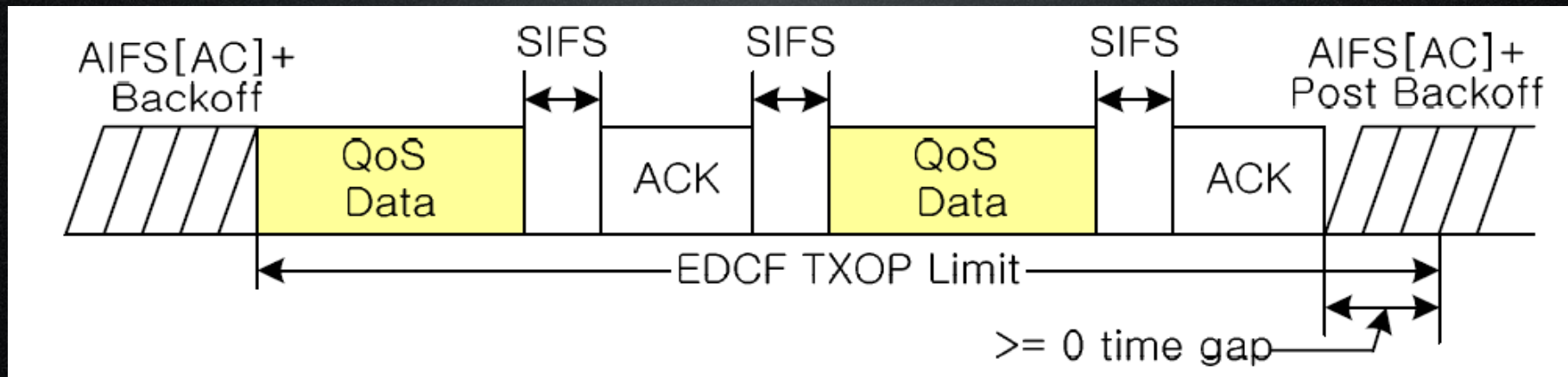
Transmission Opportunity (TXOP)

- TXOP: the duration a QSTA is enabled to transmit frame(s)
- When will a QSTA get a TXOP ?
 - Win a contention in EDCA during CP
 - Receive a CF-poll (“polled TXOP”) from HC

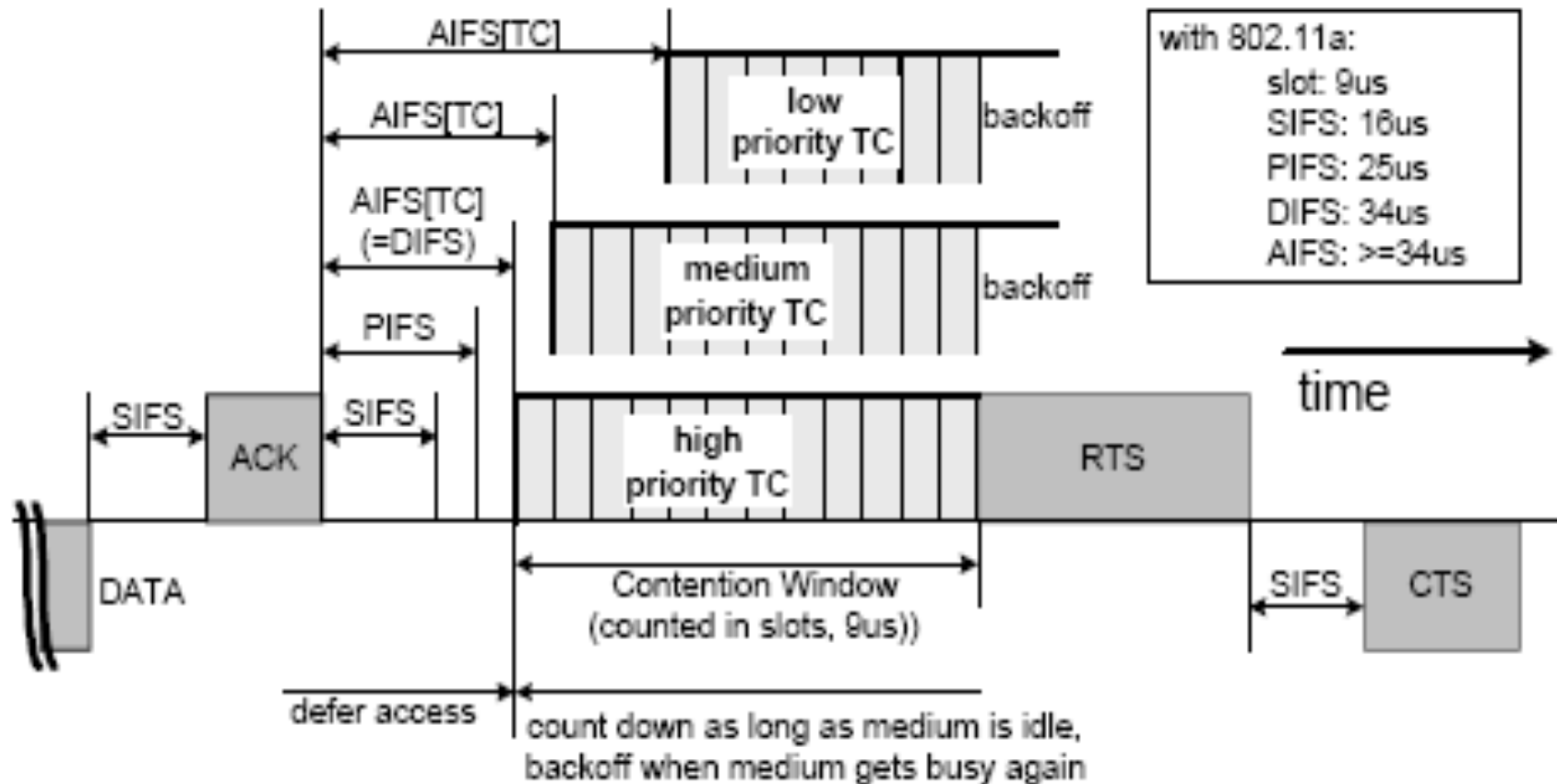


Transmission Opportunity (TXOP)

- In TXOP, frames exchange sequences are separated by SIFS



Multiple backoff of MSDU streams with different priorities



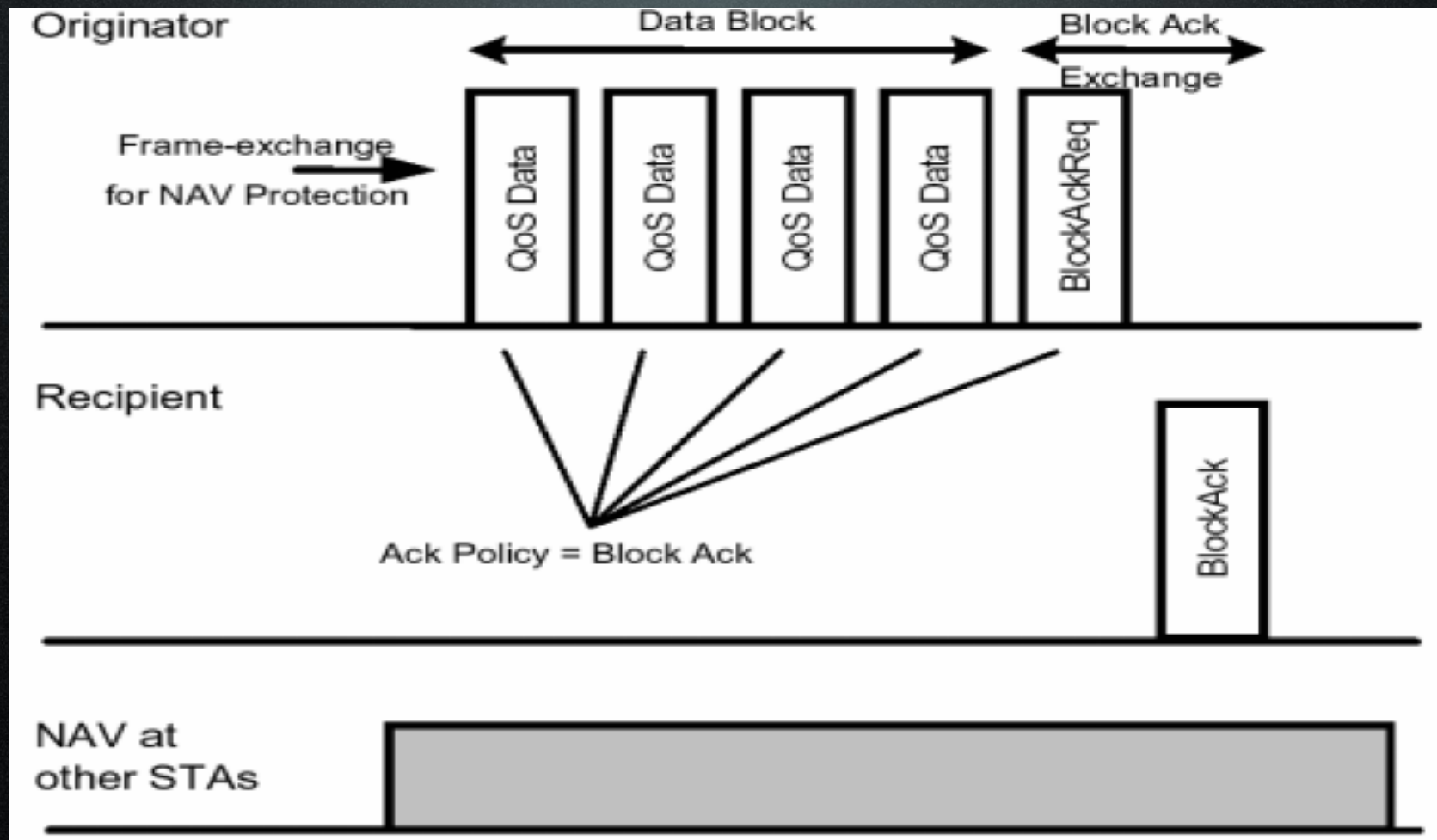
HCF Controlled Channel Access (HCCA)

- The procedure is similar to PCF
- Hybrid Coordinator (HC)
 - Operate at QAP
 - Control the iteration of CFP and CP
 - By using beacon and CF-End frame and NAV Mechanism (Same as PCF)
 - Use polling Scheme to assign TXOP to QSTA
 - Issue CF-poll frame to poll QSTA
 - Polling can be issued in both CFP & CP

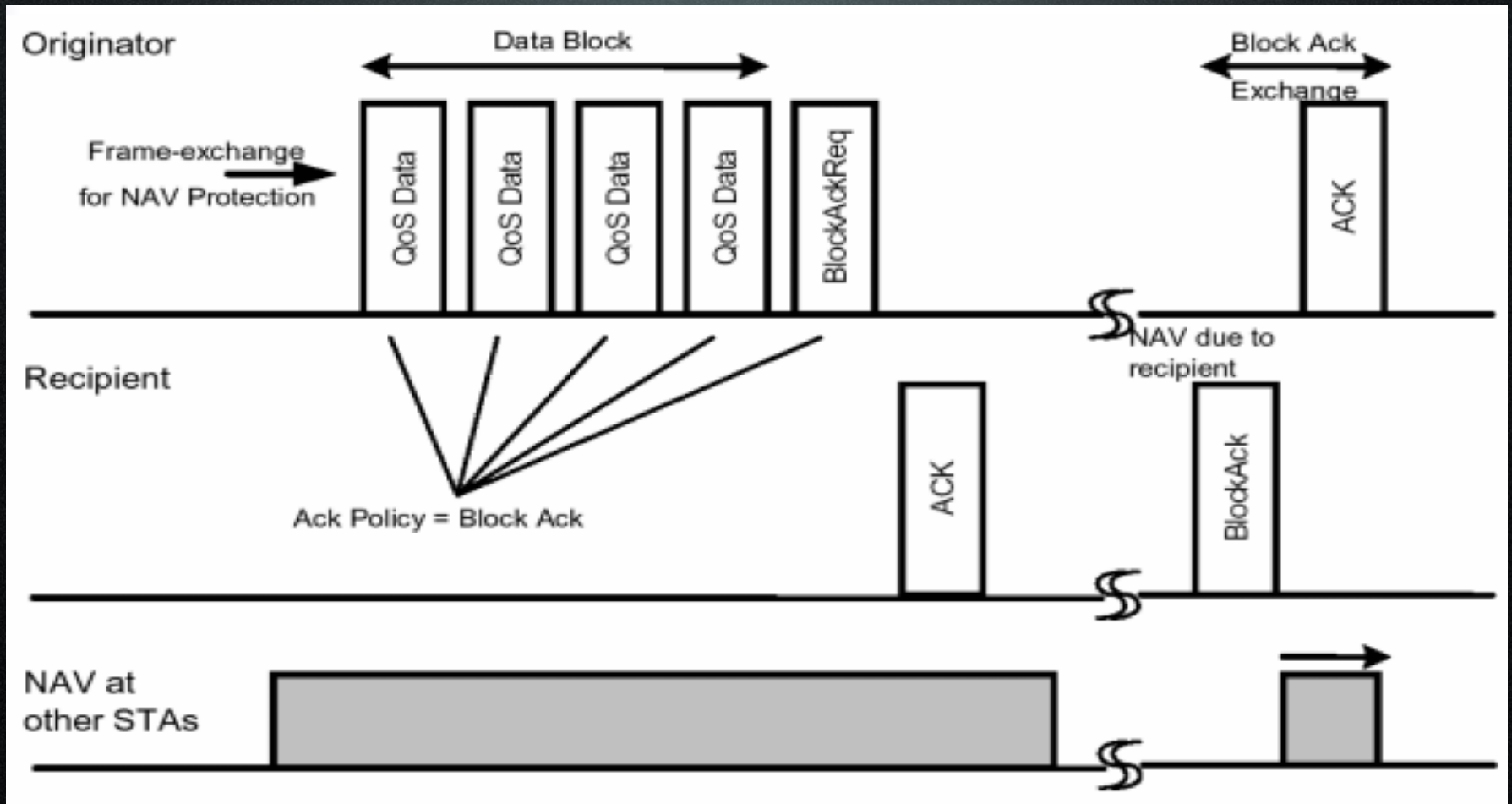
Brief of Block Ack

- Optional function
- Improve channel efficiency
 - By aggregating several acks into one frame
- Two types
 - Immediate Block Ack
 - Suitable for High-bandwidth, low latency traffic
 - Delayed Block Ack
 - Suitable for applications tolerating moderate latency

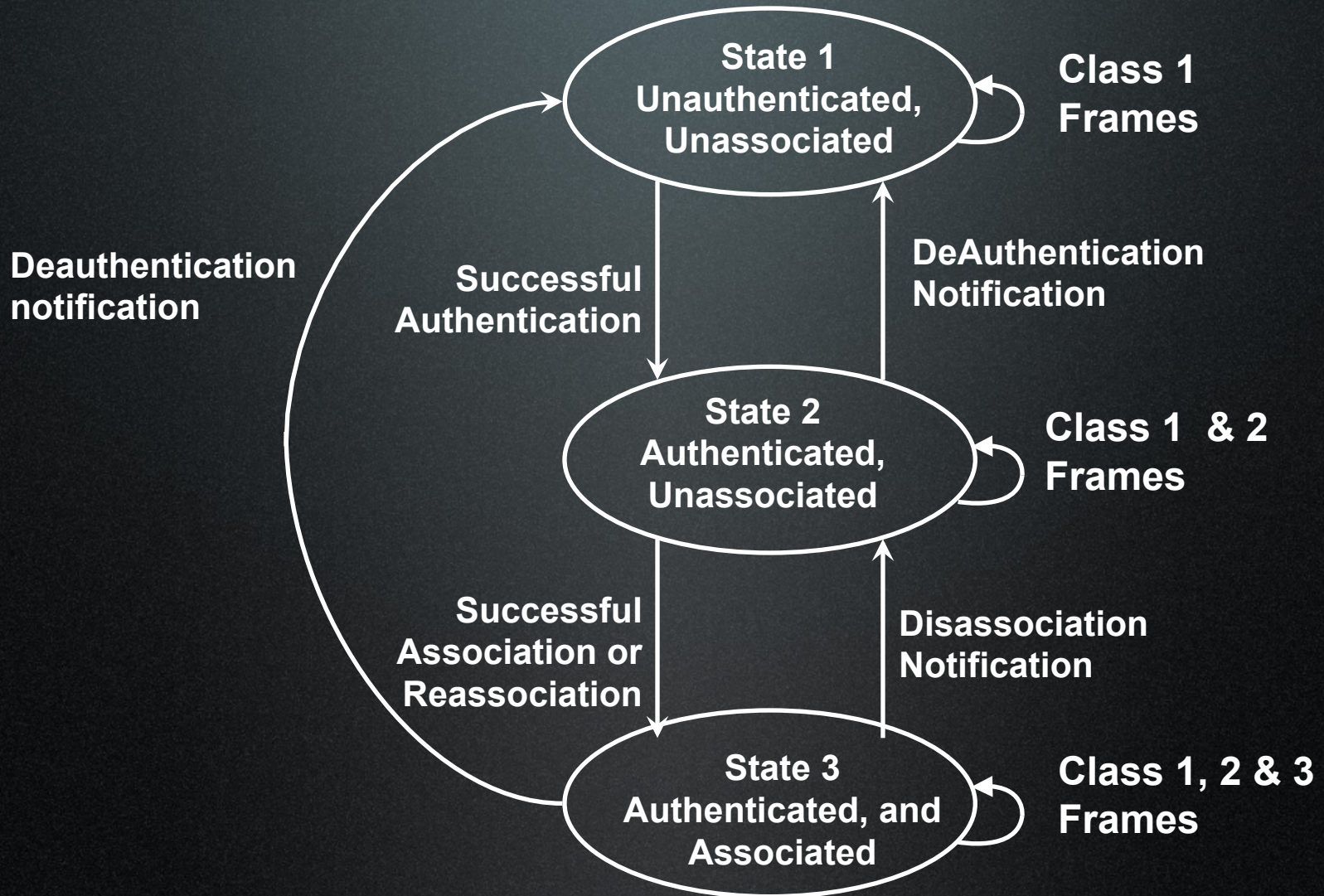
Immediate Block Ack



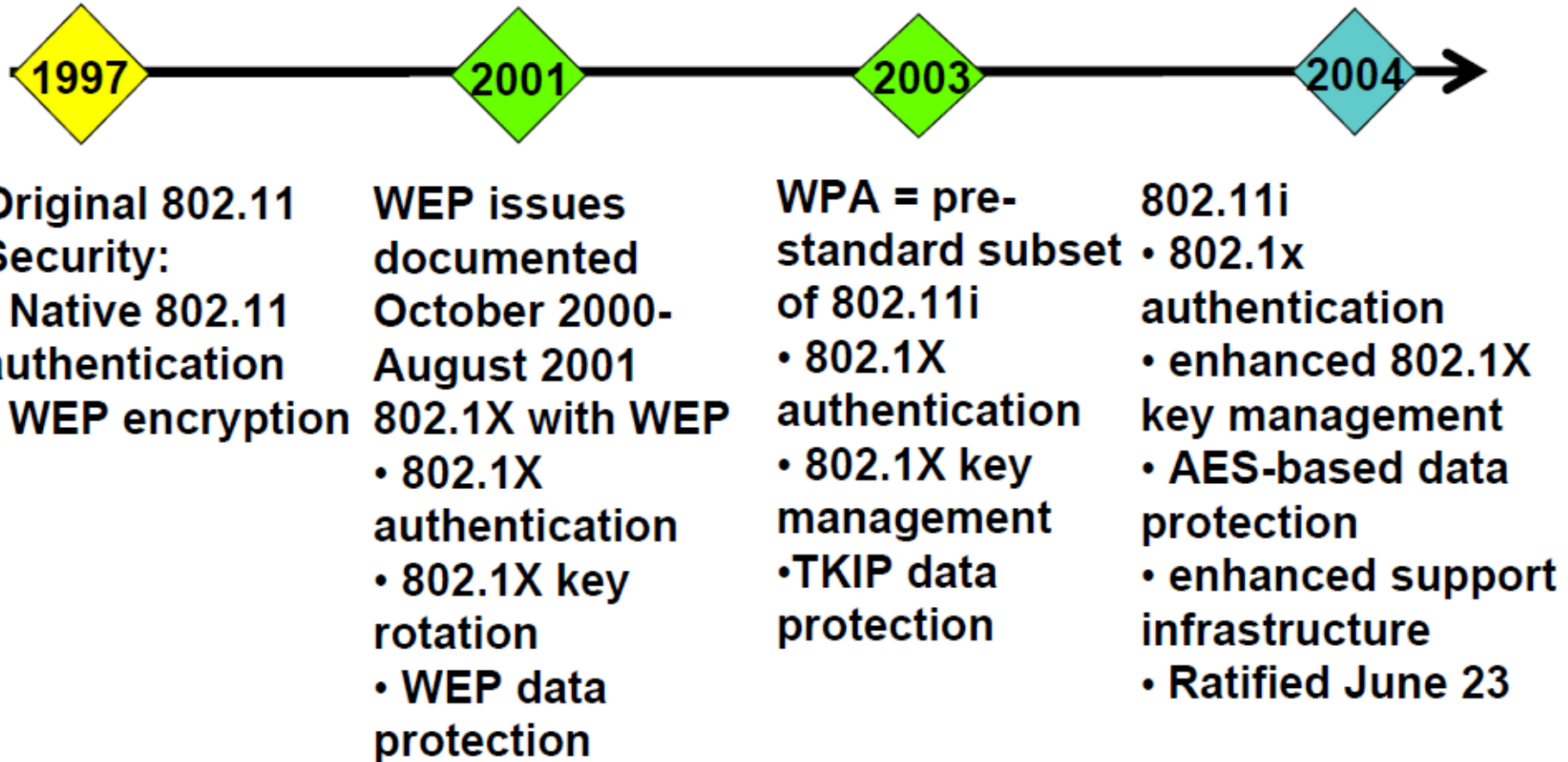
Delayed Block Ack



802.11 & security...



Chronology of Events



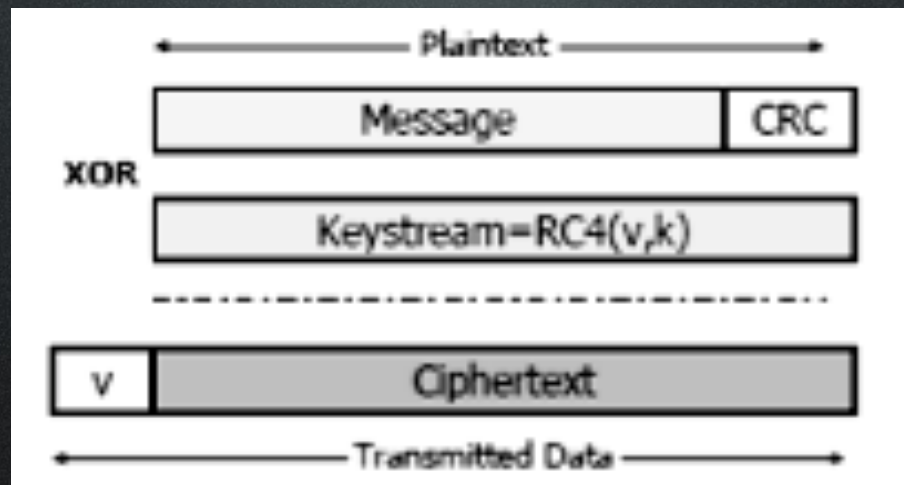
WEP?

- **WEP** relies on a secret key which is shared between the sender and the receiver.
- **SENDER**: Mobile station
- **RECEIVER**: Access Point
- **Secret Key** is used to encrypt packets before they are transmitted
- **Integrity Check** is used to ensure packets are not modified in transit.
- The standard does not discuss how shared key is established
- In practice, most installations use a **single key** which is shared between all mobile stations and access points

WEP?

- To send a **message M**:
- Compute a **checksum $c(M)$** (is not depend on secret key k)
- Pick an **IV v** and generate a **keystream $RC4(v,k)$**
- **XOR $\langle M, c(M) \rangle$** with the keystream to get the **ciphertext**
- Transmit **v** and **ciphertext** over a radio link
- When received a message **M**
- Use **transmitted v** and the **shared key k** to generate the **Keystream $RC4(v,k)$**
- **XOR** the **ciphertext** with **$RC4(v,k)$** to get **$\langle M', c' \rangle$**
- Check is **$c' = c(M')$**
- If it is, accept **M'** as the message transmitted

WEP?



RC4

- WEP uses the RC4 encryption algorithm known as “stream cipher” to protect the confidentiality of its data.
- Stream cipher operates by expanding a short key into an infinite pseudo-random key stream.
- Sender XORs the key stream with plaintext to produce the ciphertext.
- Receiver has the copy of the same key, and uses it to generate an identical key stream.
- XORing the key stream with the ciphertext yields the original message.

Two simple flaws...

- If an attacker flips a bit in ciphertext, then after decryption, that bit in the plaintext will be flipped.
- If an eavesdropper intercepts two ciphertexts encrypted with the same key stream, it is possible to obtain the XOR of the two plaintexts.

WEP vs. WPA

