

HY425 Lecture 01: Introduction

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People

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Lectures

Monday–Wednesday, 11:00–13:00, Rho Alpha 203

Friday 11:00–13:00, Rho Alpha 203 on a need basis

Course topics

Advanced Computer Architecture

- ▶ Principles of computer systems design (0.5 week)
- ▶ Basic pipelining (1 week)
- ▶ Instruction-level parallelism in HW (3 weeks)
- ▶ Instruction-level parallelism in SW (2 weeks)
- ▶ Memory hierarchies (2 weeks)
- ▶ Multiprocessors (2 weeks)
- ▶ Storage systems (1 week, but rarely make it)
- ▶ Interconnection networks (1 week, but rarely make it)
- ▶ Reserved for future use (0.5 week)

Assignments

Paper assignments

- ▶ Homework problems on paper
- ▶ Nominal load per homework: 3 hours
- ▶ Expect 3–5 paper assignments

Machine assignments

- ▶ Alternating with homework assignments
- ▶ Nominal load per assignment: 20 hours
- ▶ Expect 2–3 machine assignments

Assignments

Lab assignments

- ▶ Simulation of essential processor components
- ▶ Dynamic binary instrumentation (PIN)
- ▶ Projects are individual (no exceptions)
- ▶ Potential topics:
 1. Event counting
 2. Branch prediction
 3. Cache design, prefetching
 4. Multiprocessor coherence and consistency (optional)

Grading

Exams

- ▶ Midterm (Friday 29.10.10): 20%
- ▶ Final: 20%
- ▶ No threshold on exams, just get a 5.0+ average

Assignments

- ▶ Programming assignments 35%
- ▶ Homework assignments on paper 25%
- ▶ Do not cheat, we use moss

Other important course information

- ▶ Web page: <http://www.csd.uoc.gr/~hy425>
- ▶ Mailing list: hy425-list@csd.uoc.gr

Renaissance in computer design ending?

Microprocessors

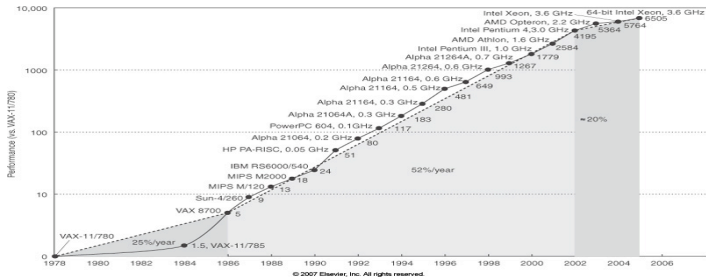
- ▶ RISC processors shifted focus to ILP, caches
- ▶ “Free” performance scaling with new technology
- ▶ Sustainable performance improvement until about 2002

New challenges

- ▶ Performance wall (lack of ILP, faster clocks, long latencies)
- ▶ Power wall (Performance per Watt drops)
- ▶ Reliability wall (more components, higher failure rates)

Renaissance in computer design ending?

Processor performance trends



Defining computer architecture

Instruction Set Architecture

- ▶ ISAs converged to a common RISC paradigm
 - ▶ CISC ISAs implemented on RISC pipelines
- ▶ Load-store architectures, general-purpose registers
- ▶ Aligned memory addressing, simple addressing modes
- ▶ Byte, word, double-word operands
- ▶ Arithmetic, logic, control operations
- ▶ Fixed-length encoding

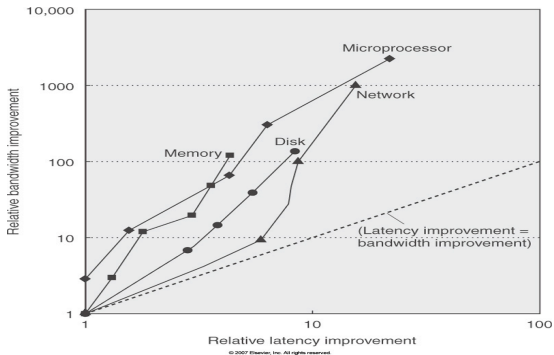
Defining computer architecture

Hardware Organization

- ▶ Processor architecture
 - ▶ Pipelining, hazards, ILP, HW/SW interface
- ▶ Memory hierarchies
- ▶ Interconnects
- ▶ I/O systems
- ▶ Hardware technology used (e.g. component size)
- ▶ Computer architecture focuses on **organization** and **quantitative principles** of design

What drives CA research and innovation?

Latency lags bandwidth



Transistor size trends and questions

Feature sizes, higher performance?

- ▶ Transistor size went down from 10 microns to 45 nanometers
- ▶ Quadratic increase in density, linear drop in feature size
- ▶ Linear increase in transistor performance

Where is the catch?

- ▶ Lower voltage to maintain safe operation
- ▶ Higher resistance and capacitance per unit of length
- ▶ Shorter wires but with higher resistance/capacitance
- ▶ **Wire delays improving poorly compared to transistors**

Power

The power equation

$$Power_{dynamic} = \frac{1}{2} \times Capacitive\ load \times Voltage^2 \times frequency$$

$$Energy_{dynamic} = Capacitive\ load \times Voltage^2$$

$$Power_{static} = Current_{static} \times Voltage$$

- ▶ Power due to switching more transistors increases
- ▶ Static power due to leakage current increasing

Measuring reliability

Reliability equations

MTTF = Mean Time To Failure

$$FIT = \text{Failures In Time (per billion hours)} = \frac{1}{MTTF}$$

MTTR = Mean Time to Repair

$$\text{Module availability} = \frac{MTTF}{MTTF + MTTR}$$

$$FIT_{\text{system}} = \sum_{i=1}^{\text{\#components}} FIT_i$$

Comparing design alternatives

Design X is n times faster than design Y

$$\frac{\text{Execution time}_Y}{\text{Execution time}_X} = n$$

- ▶ **Wall-clock time:** time to complete a task
- ▶ **CPU time:** time CPU is busy
- ▶ **Workload:** Mixture of programs (including OS) on a system
- ▶ **Kernels:** Common, important functions in applications
- ▶ **Microbenchmarks:** Synthetic programs trying to:
 - ▶ Isolate components and measure performance
 - ▶ Imitate workloads of real world in a controlled setting

Benchmarks

Desktop

- ▶ SPECintCPU (revised every few years)
- ▶ Real programs measuring processor-memory activity

Multi-core desktop/server

- ▶ SPECintCOMP, SPECintMPI (scientific), SPECintCapc (graphics)
- ▶ Focus on parallelism, synchronization, communication

Benchmarks

Client/Server

- ▶ SPECjbb, SPECjms, SPECjvm, SPECsfs, SPECmail, . . .
- ▶ Measuring throughput (how many tasks per unit of time)
- ▶ Measuring latency (how quickly does client get response)

Embedded systems

- ▶ EEMBC, MiBench
- ▶ Measuring performance, throughput, latency

Summarizing performance

Arithmetic mean of wall-clock time

- ▶ Biased by long-running programs
- ▶ May rank designs in non-intuitive ways:
 - ▶ Machine A: Program $P_1 \rightarrow 1000$ secs., $P_2 \rightarrow 1$ secs.
 - ▶ Machine B: Program $P_1 \rightarrow 800$ secs., $P_2 \rightarrow 100$ secs.
 - ▶ **What if machine runs P_2 most of the time?**

Measuring against a reference computer

$$n = \frac{SPEC_{ratio_A}}{SPEC_{ratio_B}} = \frac{\frac{Execution\ time_{reference}}{Execution\ time_A}}{\frac{Execution\ time_{reference}}{Execution\ time_B}} = \frac{Execution\ time_B}{Execution\ time_A} = \frac{Performance_A}{Performance_B}$$

Summarizing performance (cont.)

Example

	Computer A	Computer B	Computer C
Program P1 (secs)	1	10	20
Program P2 (secs)	1000	100	20
Total time (secs)	1001	110	40

Means

- ▶ Total time ignores program contribution to total workload
- ▶ Arithmetic mean biased by long programs
- ▶ Weighted arithmetic mean a better choice?
- ▶ How do we calculate weights?

Summarizing performance (cont.)

Weighted arithmetic mean

$$\sum_{i=1}^n \text{Weight}_i \times \text{Time}_i$$

Example, $W(1) = W(2) = 50$

	Computer A	Computer B	Computer C
Program P1 (secs)	1	10	20
Program P2 (secs)	1000	100	20
Total time (secs)	1001	110	40
Weighted mean	500.50	55.00	20.00

Summarizing performance (cont.)

Weighted arithmetic mean

$$\sum_{i=1}^n \text{Weight}_i \times \text{Time}_i$$

Example, $W(1) = 0.909$ $W(2) = 0.091$

	Computer A	Computer B	Computer C
Program P1 (secs)	1	10	20
Program P2 (secs)	1000	100	20
Total time (secs)	1001	110	40
Weighted mean	91.91	18.19	20.00

Summarizing performance (cont.)

Weighted arithmetic mean

$$\sum_{i=1}^n \text{Weight}_i \times \text{Time}_i$$

Example, $W(1) = 0.999$ $W(2) = 0.001$

	Computer A	Computer B	Computer C
Program P1 (secs)	1	10	20
Program P2 (secs)	1000	100	20
Total time (secs)	1001	110	40
Weighted mean	2.00	10.09	20.00

Summarizing performance (cont.)

Using ratios

- ▶ Ratios against reference machine are independent of mixture of programs

Geometric mean

$$\sqrt[n]{\prod_{i=1}^n \text{Execution time ratio}_i}$$

$$\frac{\text{Geometric mean}_A}{\text{Geometric mean}_B} = \text{Geometric mean}\left(\frac{A}{B}\right) \quad (1)$$

Pros and cons of geometric means

Pros

- ▶ Consistent rankings, independent of program frequencies
- ▶ Not influenced by peculiarities of any single machine

Cons

- ▶ **Geometric mean does not predict execution time**
 - ▶ Sensitivity to benchmark vs. machine remains
 - ▶ Encourages machine tuning for specific benchmarks
 - ▶ **Benchmarks can not be touched, but compilers can!**
- ▶ Any “averaging” metric loses information

Qualitative principles of design

Taking advantage of parallelism

- ▶ Use pipelining to overlap instructions
- ▶ Use multiple execution units
- ▶ Use multiple cores
- ▶ Use multiple processors to increase throughput

Locality

- ▶ Programs reuse instructions and data
- ▶ 90-10 rule
 - ▶ 90% of execution time spent running 10% of instructions
- ▶ Programs access data in nearby addresses

Qualitative principles of design (cont.)

Make the common case fast

- ▶ Trade-off's in design (e.g. performance vs. power/area)
- ▶ Provide efficient design for the common case
- ▶ Amdahl's Law

Using performance equations

- ▶ Processor performance equation =
f(cycle time, instruction count, stalls, instruction mix, ...)

Amdahl's Law

$$speedup = \frac{\text{performance with enhancement}}{\text{performance without enhancement}}$$

$$execution\ time_{new} = execution\ time_{old} \times \left((1 - fraction_{enhanced}) + \frac{fraction_{enhanced}}{speedup_{enhanced}} \right)$$

$$speedup_{overall} = \frac{execution\ time_{old}}{execution\ time_{new}} = \frac{1}{(1 - fraction_{enhanced}) + \frac{fraction_{enhanced}}{speedup_{enhanced}}} \Rightarrow speedup_{overall} \rightarrow \frac{1}{1 - fraction_{enhanced}}$$

Processor performance equation

CPU time

$$\text{CPU time} = \text{CPI} \times \text{cycle time}$$

$$\text{CPI} = \frac{\text{CPU clock cycles}}{\text{instruction count}} \Rightarrow$$

$$\text{CPU time} = \text{instruction count} \times \text{CPI} \times \text{cycle time} \Rightarrow$$

$$\text{CPU time} = \frac{\text{instructions}}{\text{program}} \times \frac{\text{clock cycles}}{\text{instruction}} \times \frac{\text{seconds}}{\text{clock cycle}}$$

Processor performance equation

How can CA help?

- ▶ Technology has been providing faster clock speeds
 - ▶ Main performance factor for almost 20 years
 - ▶ Trend seems to reverse
 - ▶ Limitations due to power consumption, reliability
- ▶ Architecture can pack more computing power in same area
- ▶ Architecture can improve CPI
- ▶ Algorithms and compilers can reduce instruction count

Processor performance equation

Instruction mixes

$$CPU\ time = CPI \times cycle\ time$$

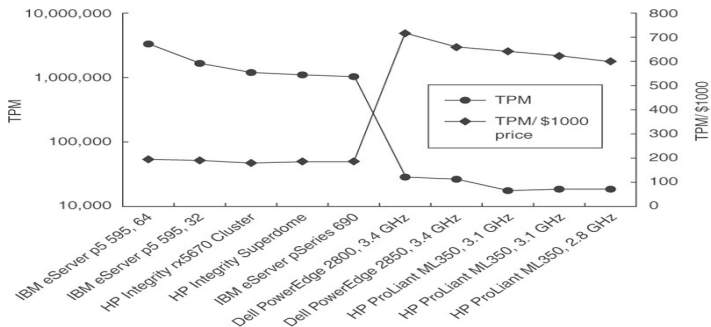
$$CPI = \frac{CPU\ clock\ cycles}{instruction\ count} \Rightarrow$$

$$CPU\ time = instruction\ count \times CPI \times cycle\ time \Rightarrow$$

$$cpu\ time = \frac{instructions}{program} \times \frac{clock\ cycles}{instruction} \times \frac{seconds}{clock\ cycle} \Rightarrow$$

$$CPU\ time = \left(\sum_{i=1}^n IC_i \times CPI_i \right) \times cycle\ time$$

Price/performance



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Concluding remarks

Fallacies and pitfalls

- ▶ Ignoring Amdahl's law
- ▶ Reliability is as good as that of the most faulty component
- ▶ Cost of processor dominates system cost?
 - ▶ Currently, on servers and laptops storage dominates cost!
- ▶ Benchmarks remain valid for long
 - ▶ Workloads evolve (Internet, laptops, handheld computers, sensors, controllers, actuators, ...)
 - ▶ Tuning for depreciated benchmarks undesirable

Concluding remarks (cont.)

Fallacies and pitfalls

- ▶ Reliability metrics ignoring lifetime of component
- ▶ Peak performance is expected performance
- ▶ Detecting but not correcting faults
 - ▶ Many components in the architecture non-critical for correct operation
 - ▶ Important to protect, check and duplicate critical components