



Presentation for use with the textbook, *Algorithm Design and Applications*, by M. T. Goodrich and R. Tamassia, Wiley, 2015

Sorting Lower Bound



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Comparison-Based Sorting



- Many sorting algorithms are comparison based.
 - They sort by making comparisons between pairs of objects
 - Examples: bubble-sort, selection-sort, insertion-sort, heap-sort, merge-sort, quick-sort, ...
- Let us therefore derive a lower bound on the running time of any algorithm that uses comparisons to sort n elements, x_1, x_2, \dots, x_n .

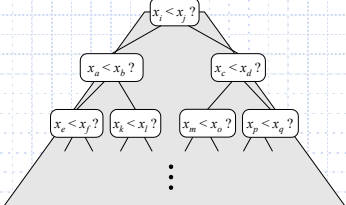
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graph TD
    A[ ] --> B{Is  $x_i < x_j$ ?}
    B -- yes --> C[ ]
    B -- no --> D[ ]
  
```

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Counting Comparisons

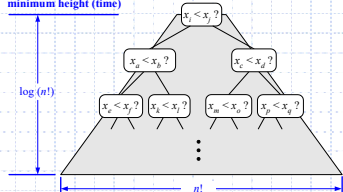
- Let us just count comparisons then.
- Each possible run of the algorithm corresponds to a root-to-leaf path in a **decision tree**



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Decision Tree Height

- The height of the decision tree is a lower bound on the running time
- Every input permutation must lead to a separate leaf output
- If not, some input ...4...5... would have same output ordering as ...5...4..., which would be wrong
- Since there are $n! = 1 \cdot 2 \cdot \dots \cdot n$ leaves, the height is at least $\log(n!)$



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The Lower Bound



- Any comparison-based sorting algorithm takes at least $\log(n!)$ time
- Therefore, any such algorithm takes time at least

$$\log(n!) \geq \log\left(\frac{n}{2}\right)^{\frac{n}{2}} = (n/2) \log(n/2).$$

- That is, any comparison-based sorting algorithm must run in $\Omega(n \log n)$ time.

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