Outline

• Objectives of OCL
• Why to use OCL
• Presentation of OCL
• Assertions and Programming Languages
What is OCL (Object Constraint Language)?

- A **formal language** for specifying **constraints** on o-o models
- It is **declarative** (it describes **what** rather than **how**)
- it is a **typed** language
  - and more … user friendly (comparing it with other formal languages)

Constraints?
- Some constraints can be expressed graphically (e.g. multiplicity of an association, partition subclasses).
- Some other cannot, e.g.:
  - constraints involving >1 associated classes
  - constraints involving attribute values (and their combination)
  - pre/post-conditions on operations

OCL can be used to express constraints formally

Why to write OCL constraints?

*Why to express explicitly such constraints?*

They make the models more precise

- so that to understand them better
- so that the programmers can implement them (correctly)
- so that to allow a **formal validation** of the model prior to implementation

They can be “translated” to assertions in programming languages
  - some CASE tools offer these validation and translation facilities
Class Diagrams are not very precise

- Can a minor (underage) work for a company?
- Can a company hire a person already hired?

Class Diagrams are not very precise (II)

- Can a person start a job before his/her birth?
- Can a promotion lower the salary of an employee?
- Is there any lower bound for the salaries of those working for more than 10 years?
Class Diagrams are not very precise (III)

... The importance of Background Knowledge

- What if you had to build a system whose class diagrams were in Spanish?

Class Diagrams are not very precise (IV)

Valid object diagram
Object Constraint Language (OCL)

- OCL is a formal language used to describe expressions on UML models.
- OCL expressions typically specify invariant conditions that must hold for the system being modeled.
- They also specify queries over objects described in a model.
- OCL is a typed language, so that each OCL expression has a type. To be well formed, an OCL expression must conform to the type conformance rules of the language. For example, you cannot compare an Integer with a String.
- When OCL expressions are evaluated, they do not have side effects; i.e. their evaluation cannot alter the state of the corresponding executing system. However, OCL expressions could be used to specify operations / actions that, when executed, do alter the state of the system.

Where to use OCL?

OCL can be used for a number of different purposes:
- To specify invariants on classes and types in the class model
- To describe pre- and post conditions on Operations and Methods
- To specify derivation rules for attributes for any expression over a UML model.
- To describe Guards in State Diagrams
- To specify target (sets) for messages and actions
- To specify type invariant for Stereotypes
- As a query language

UML modelers can use OCL to
- to specify application-specific constraints in their models.
- to specify queries on the UML model, which are completely programming language independent
The main types of OCL Constraints

- **Invariants on classes**
  - conditions to be true always by all instances of a class
    - e.g. salary > 1000 Euro

- **pre-conditions on operations**
  - conditions to be true before the execution of an operation
    - e.g. the operation “fire” can be executed only on a hired person

- **post-conditions on operations**
  - conditions to be true after the execution of an operation
    - e.g. after “withdraw(amount)” the balance of the bank account should be reduced by “amount”.

How we can specify a constraint?

- Declaration of the context of a constraint by referencing the model element that a constraint applies to

- Declaration of the type of a constraint (inv, pre, post)

- Expressing the desired condition by referencing properties of model elements and using various operations that are supported.
Context Declaration

- Specifies the element the constraint applies to.
- The context can be
  - a class (for invariants)
  - an operation (for pre/post-conditions)

**Example:**

<table>
<thead>
<tr>
<th>Employee</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
</tr>
<tr>
<td>age</td>
</tr>
<tr>
<td>salary</td>
</tr>
<tr>
<td>SetAge(a)</td>
</tr>
<tr>
<td>SetSalary(s)</td>
</tr>
</tbody>
</table>

**Context** Employee** inv: self.salary > 1000**

**Context** Employee::SetSalary(salary) **pre: salary > 1000**

---

Constraint names and comments

<table>
<thead>
<tr>
<th>Employee</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
</tr>
<tr>
<td>age</td>
</tr>
<tr>
<td>salary</td>
</tr>
<tr>
<td>SetAge(a)</td>
</tr>
<tr>
<td>SetSalary(s)</td>
</tr>
</tbody>
</table>

**Context** Employee::SetAge (age) **pre: age > 0**

**Context** Employee::SetAge (age) **pre positive_age: age > 0**

Optional constraint name
Allowing the constraint to be referenced by name.

**Context** Employee::SetAge (age) **pre positive_age: age > 0**

-- the age should always be positive

Comment
Everything immediately following the two dashes up to and including the end of line is part of the comment.
self

In most cases, the keyword self can be dropped because the context is clear. As an alternative for self, a different name can be defined playing the part of self.

Selectors (how we reference elements)

self.age // returns the age of a particular person
self.employment // returns the employer (company) of a person
self.employer // as before

self.employment // returns the set of all employees of a company
self.president // returns the singleton with the president of a company
self.stockPrice() // returns the value this method would return
Selectors (how we reference elements)

Because the multiplicity of the role manager is one, self.manager is an object of type Person. This happens when the multiplicity of the association-end has a maximum of one ("0..1" or "1")

Selectors Referencing Association Classes

The salary should be > 1000

Context Person inv: self.employment.salary > 1000

We use dot and the name of the association class starting with a lowercase letter
Selectors
Referencing Recursive Associations

The age of a children should be less than the age of its parents.

Here the name of the association class alone is not enough. We need to distinguish the direction in which the association is navigated.

To make the distinction, the rolename of the direction in which we want to navigate is added to the association class name, enclosed in square brackets.

Context Person inv: self.hasParent[parent].age > self.age

self.hasParent.age is invalid

Selectors
Referencing Recursive Associations

Let c be a company. The name of the company that owns c should be different than the name of c.

Context Company inv: self.owns[owner].name <> self.name
Operations

- **Boolean Operations**
  - and  // ∧
  - or  // ∨
  - not  // ¬
  - implies  // →
  - xor
- **Comparison operations**
  - <, >, <=, >=, <>, ==
- **Arithmetic**
  - +, -, *, /, abs(), div, floor(), round()
- **String operations**
  - concat(s1, s2), toUpper(s),

- **Nil**
  - if an attribute attr of an object obj has no value then obj.attr returns nil
- **Empty**
  - if there are no associated objects to an object obj through an association assoc then obj.assoc returns the empty bag {}.
- **Nil <> Empty**

---

**Referring to enumerations**

- **Enumerations**
  - Person
    - salary
    - sex
  - JobType
    - admin
    - programmer
    - secretary

**Context** Person inv: self.job=JobType::admin implies self.salary > 10.000

**Context** Person inv: self.name=“Yannis” implies self.sex::Male
Collections in OCL

- Allow us to refer to the objects that are referred using associations (typically in those with upper multiplicity > 1)

```
<table>
<thead>
<tr>
<th>Person</th>
<th>works</th>
<th>Company</th>
</tr>
</thead>
<tbody>
<tr>
<td>age</td>
<td>0,*</td>
<td>name</td>
</tr>
</tbody>
</table>
```

```
p1: Person
p2: Person
c1:Company
c2:Company
c3:Company
```

c2.works c2.works

Collections in OCL (II)

Single navigation of an association results in a Set, combined navigations in a Bag, and navigation over associations adorned with {ordered} results in an OrderedSet.

```
Polygon definedBy 3..* {ordered} Point
```

Collection is an abstract type, with the concrete collection types (Set, Sequence, and Bag) as its subtypes.
Objects and Collections

- **Objects**
  - are instances of classes, including the predefined ones (e.g. Integer)

- **Sets**
  - a “set” of objects
  - example: Set { p1, p2}

- **Bag**
  - duplicates allowed
  - example: Bag { p1, p1, p1, p2, p1}

- **Sequence**
  - is a bag of ordered elements
  - example: Sequence {p1, p2, p3, p1 } // <p1, p2, p3, p1>

Collection Operations

The type Collection defines a large number of predefined operations to enable the modeler to manipulate collections.

As OCL is an expression language, collection operations never change collections (rather than changing the original collection they project the result into a new one).

- c1->Size()       // number of elements of c1
- c1->count(elem)  // counts the number of occurrences of elem in c1
- c1->includes(elem) // checks if elem is member of c1
- c1->includesAll(coll) // checks if coll is contained in c1
- c1->excludes(elem) // returns True if elem is not member of in c1
- c1->isEmpty()    // checks if c1={}  

- c1->forAll(expr)  // returns True if expr is true for all elements of c1
- c1->exists(expr)  // returns True if expr is true for at least one element of c1
- c1->select(expr)  // returns the elements of c1 that satisfy expr
- c1->reject(expr)  // returns the elements of c1 that do not satisfy expr

- SET OPERATIONS:
  - c1->union(c2), c1->intersection(c2), c1-c2
Collection Operations: Examples

- p1.works->size() is 2
- p1.works->count(c3) is 0
- p1.works->includes(c2)
  - is True
- p1.works->includes(c3) is False
- c2.works->includesAll(c1.works)
  - is True
- c1.works->includesAll(c2.works)
  - is False
- c3.works->isEmpty() is True

Collection Operations: Examples (II)

- c2.works->forAll{ x | x.age>20 and x.age < 70}
  - is False
- c2.works->exists{ x | x.age>20 and x.age < 70}
  - is True
- c2.works->select{ x | x.age>20 and x.age < 70}
  - will return {p2}
- p1.works->intersection(p2.works)
  - will return {c2}
- p1.works - p2.works
  - will return {c1}
A single object can be treated as a singleton

A single object can be used as a Set as well. It then behaves as if it is a Set containing the single object. The usage as a set is done through the arrow followed by a property of Set.

context Company
inv: self.manager->size() = 1

Select / Reject

The reject operation is available in OCL for convenience, because each reject can be restated as a select with the negated expression. Therefore, the following two expressions are identical:

collection->reject(v : Type | boolean-expression-with-v)
collection->select(v : Type | not(boolean-expression-with-v))

The collection of all the employees who are not married is empty:

context Company
inv: self.employee->reject(isMarried)->isEmpty()
Collect operation

The select and reject operations always result in a sub-collection of the original collection.

When we want to specify a collection which is derived from some other collection, but which contains different objects from the original collection (i.e., it is not a sub-collection), we can use a collect operation.

The collect operation uses the same syntax as the select and reject and is written as one of:

- `collection->collect( v : Type | expression-with-v )`
- `collection->collect( v | expression-with-v )`
- `collection->collect( expression )`

The value of the reject operation is the collection of the results of all the evaluations of expression-with-v.

An example: specify the collection of birthDates for all employees in the context of a company. This can be written in the context of a Company object as one of:

- `self.employee->collect( birthDate )`
- `self.employee->collect( person | person.birthDate )`
- `self.employee->collect( person : Person | person.birthDate )`

Shorthand for Collect

Because navigation through many objects is very common, there is a shorthand notation for the collect that makes the OCL expressions more readable.

Instead of

```
self.employee->collect(birthdate)
```

we can also write:

```
self.employee.birthdate
```

In general, when we apply a property to a collection of Objects, then it will automatically be interpreted as a collect over the members of the collection with the specified property. For any property name that is defined as a property on the objects in a collection, the following two expressions are identical:

```
collection.propertyname
```

and so are these if the property is parameterized:

```
collection.propertyname (par1, par2, ...)
collection->collect (propertyname(par1, par2, ...))
```
Collect (3)

When the source collection is a **Set** the resulting collection is not a Set but a **Bag**.

If the source collection is a **Sequence** or an **OrderedSet**, the resulting collection is a **Sequence**.

When more than one employee has the same value for birthDate, this value will be an element of the resulting Bag more than once.

The Bag resulting from the collect operation always has the same size as the original collection.

It is possible to make a Set from the Bag, by using the **asSet** property on the Bag. Example:

```
self.employee->collect(birthDate)->asSet()
```

Results in the Set of different birthDates from all employees of a Company.

Examples with Bags and other operations

- **employment.age** is a bag
- **employment.income** is a bag
- **employment.income->asSet()** returns all distinct incomes of the employees
Examples of Invariants (using collection operations)

<table>
<thead>
<tr>
<th>Person</th>
<th>Company</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>name</td>
</tr>
<tr>
<td>age</td>
<td></td>
</tr>
</tbody>
</table>

employment

All persons should have positive age

**Context** Person  **inv**: self.age > 0

All persons that work for a company should be adults

**Context** Company  **inv**: self.employment->forall( x | x.age > 18)

---

Examples of Invariants (using collection operations)

A person can be a manager of only one company

<table>
<thead>
<tr>
<th>Person</th>
<th>Company</th>
</tr>
</thead>
<tbody>
<tr>
<td>name</td>
<td>name</td>
</tr>
<tr>
<td>age</td>
<td></td>
</tr>
</tbody>
</table>

manager

1..* employee

0..*

manages

0..1 employer

0..*

All companies should have managers that are not employers of other companies

**Context** Company  **inv**: not (self.manager->exists(x | x.employer->exists(y | y<>self))

**Context** Company  **inv**: self.manager.employer->forall(x | x = self)
Another example

**Context** Person

\[
\text{inv: self.parent->excludes(self) and self.children->excludes(self)}
\]

**Forall**

- **context** Company
- \text{inv: self.employee->forAll( age <= 65 )}
- \text{inv: self.employee->forAll( p | p.age <= 65 )}
- \text{inv: self.employee->forAll( p : Person | p.age <= 65 )}

These invariants evaluate to true if the age property of each employee is less or equal to 65.

The forAll operation has an extended variant in which more than one iterator is used. Both iterators will iterate over the complete collection. Effectively this is a forAll on the Cartesian product of the collection with itself.

- **context** Company \text{inv: self.employee->forAll( e1, e2 : Person | e1 <> e2 implies e1.forename <> e2.forename) }

This expression evaluates to true if the forenames of all employees are different. It is semantically equivalent to:

- **context** Company \text{inv: self.employee->forAll( e1 | self.employee->forAll( e2 | e1 <> e2 implies e1.forename <> e2.forename))}
Examples of Constraints (using collection operations)

Pre/Post-Conditions

Context Person::hire(c:Company)

pre: not employment->includes(c)
post: employment->includes(c)

Context Person::fire(c:Company)

pre: employment->includes(c)
post: not employment->includes(c)

Context Person::increaseAge ()

post: age = age@pre +1

Context Person::Promote (inc)

post: self.income = income@pre * (1+inc)

@Pre

When the pre-value of a property evaluates to an object, all further properties that are accessed of this object are the new values (upon completion of the operation) of this object.

a.b@pre.c -- takes the old value of property b of a, say object18,
-- and then the new value of c of object18.

a.b@pre.c@pre -- takes the old value of property b of a, say object18
-- and then the old value of c of object18.

The ‘@pre’ postfix is allowed only in OCL expressions that are part of a Postcondition.

Asking for a current property of an object that has been destroyed during execution of the operation results in OclUndefined. Also, referring to the previous value of an object that has been created during execution of the operation results in OclUndefined.
Post-conditions

Result, out-parameters

The reserved word **result** denotes the result of the operation, if there is one.

```
Context Person::getIncome(d:Date): Integer
post: result = 1000
```

The right-hand-side of this definition may refer to the operation being defined (i.e., the definition may be recursive) as long as the recursion is not infinite.

When the operation has no **out** or **in/out** parameters (like in this example), then the type of result is the return type of the operation (here Integer).

When the operation has **out** or **in/out** parameters, the return type is a **Tuple**.

The postcondition for the income operation with an **out** parameter **bonus** could be:

```
Context Person::getIncome(d:Date, bonus:Integer): Integer
post: result = Tuple{bonus=300, result=1000}
```

The return type of operation calls is **Tuple(bonus: Integer, result: Integer)**.

---

Post-conditions

Result, out-parameters (2)

```
Context Person::getIncome(d:Date, bonus:Integer): Integer
post: result = Tuple{bonus=300, result=1000}
```

The out parameters need not be included in the operation call (we have to provide values only for the **in** or **in/out** parameters).

Let **Yannis** be an object of the class Person, and let **d1** be a Date.

Then, **Yannis.getIncome(d1)** is a valid operation call.

The type of the result of this operation call is **Tuple(bonus: Integer, result: Integer)**.

We can access these values using the names of the out parameters, and the keyword **result**, for example:

```
Yannis.getIncome(d1).bonus = 300 and
Yannis.getIncome(d1).result = 1000
```
An OCL expression may be used to indicate the result of a query operation. The expression must conform to the result type of the operation. Like in the pre/post-conditions, the parameters may be used in the expression. Pre/post-conditions, and body expressions may be mixed together after one operation context.

### Context

**Person::getCompany():Company**

**pre:** self.employment->size()>0

**body:** self.employment

---

### Context

**Person::getCurrentSpouse():Person**

**pre:** self.isMarried = true

**body:** self.marriages->select( m | m.ended = false).spouse
Initial and Derived Attributes

An OCL expression may be used to indicate the initial or derived value of an attribute or association end.

\[
\text{context Typename::attributeName: Type} \\
\text{init: -- some expression representing the initial value}
\]

\[
\text{context Typename::assocRoleName: Type} \\
\text{derive: -- some expression representing the derivation rule}
\]

The expression must conform to the result type of the attribute. If the context is an association end the expression must conform to the classifier at that end when the multiplicity is at most one, or Set or OrderedSet when the multiplicity may be more than one. Initial, and derivation expressions may be mixed together after one context.

**Context**

Person::income: Integer

\[
\text{init: parents.income->sum()^1} \quad \text{-- pocket allowance}
\]

\[
\text{derive: if underAge} \\
\text{then parents.income->sum()^1} \quad \text{-- pocket allowance}
\]

\[
\text{else job.salary} \quad \text{-- income from regular job}
\]

The let expression allows one to define a variable which can be used in the constraint.

**Let Expressions**

Sometimes a sub-expression is used more than once in a constraint. The let expression allows one to define a variable which can be used in the constraint.

\[
\text{context Person inv:} \\
\text{let income : Integer = self.job.salary->sum() in} \\
\text{if isUnemployed then} \\
\text{\quad income < 100} \\
\text{else} \\
\text{\quad income >= 100} \\
\text{endif}
\]

A let expression may be included in any kind of OCL expression. It is only known within this specific expression.
**«definition» expressions**

The Let expression allows a variable to be used in one OCL expression. To enable reuse of variables/operations over multiple OCL expressions we can use the stereotype «definition».

All variables and operations defined in the «definition» constraint are known in the same context as where any property of the Classifier can be used.

The syntax of the attribute or operation definitions is similar to the Let expression, but each attribute and operation definition is prefixed with the keyword 'def'.

```plaintext
context Person
def: income : Integer = self.job.salary->sum()
def: nickname : String = 'Little Red Rooster'
def: hasTitle(t : String) : Boolean = self.job->exists(title = t)
```

The names of the attributes / operations in a let expression may not conflict with the names of respective attributes/associationEnds and operations of the Classifier.

**Re-typing or casting**

In some circumstances, it is desirable to use a property of an object that is defined on a subtype of the current known type of the object. Because the property is not defined on the current known type, this results in a type conformance error.

When it is certain that the actual type of the object is the subtype, the object can be re-typed using the operation `oclAsType(OclType)`. This operation results in the same object, but the known type is the argument OclType.

When there is an object obj of type Type1 and Type2 is a subtype of Type1, then it is allowed to write:

```plaintext
obj1.oclAsType(Type2) --- evaluates to object with type Type2
```
Accessing overridden properties of supertypes

Whenever properties are redefined within a type, the properties of the supertypes can be accessed using the `oclAsType()` operation. Whenever we have a class B as a subtype of class A, and a property p1 of both A and B, we can write:

**context B inv:**

- `self.oclAsType(A).p1` -- accesses the p1 property defined in A
- `self.p1` -- accesses the p1 property defined in B

In this model fragment there is an ambiguity with the OCL expression on Dependency:

**context Dependency inv:**

- `self.source <> self`

This can either mean normal association navigation, which is inherited from ModelElement, or it might also mean navigation through the dotted line as an association class. Both possible navigations use the same role-name, so this is always ambiguous. Using `oclAsType()` we can distinguish between them with:

**context Dependency**

- `inv: self.oclAsType(Dependency).source->isEmpty()`
- `inv: self.oclAsType(ModelElement).source->isEmpty()`

---

Predefined properties on all objects

There are several properties that apply to all objects, and are predefined in OCL.

- **oclIsTypeOf (t: OclType) : Boolean**
  - returns true if the type of `self` and `t` are the same, e.g.
    - **context Person**
    - `inv: self.oclIsTypeOf(Person)` -- is true
    - `inv: self.oclIsTypeOf(Company)` -- is false

- **oclIsKindOf (t: OclType) : Boolean**
  - The `oclIsTypeOf` deals with the direct type of an object. The `oclIsKindOf` property determines whether `t` is either the direct type or one of the supertypes of an object.

- **oclInState (s: OclState) : Boolean**
  - will be discussed later on

- **oclIsNew () : Boolean**
  - It returns true if, used in a postcondition, the object is created during performing the operation. i.e., it didn’t exist at precondition time.

- **oclAsType (t : OclType) : instance of OclType**
  - we have discussed this already
All properties discussed until now in OCL are properties on instances of classes. The types are either predefined in OCL or defined in the class model. In OCL, it is also possible to use features defined on the types/classes themselves. These are, for example, the class-scoped features defined in the class model. Furthermore, several features are predefined on each type.

A predefined feature on classes, interfaces and enumerations is `allInstances()`, which results in the Set of all instances of the type in existence at the specific time when the expression is evaluated.

Example
We want to make sure that all instances of Person have unique names:

```ocl
class Person
inv: allInstances().forall(p1, p2 | p1 <> p2 implies p1.name <> p2.name)
```

The `Person.allInstances()` is the set of all persons that exist in the system at the time that the expression is evaluated and is of type `Set(Person)`.

Type conformance rules:

- **Type conformance rules:**
  - `Type1` conforms to `Type2` when they are identical or when `Type1` is a subtype of `Type2` (standard rule for all types).
  - `Collection(Type1)` conforms to `Collection(Type2)`, when `Type1` conforms to `Type2`. This is also true for `Set(Type1)`, `Set(Type2)`, `Sequence(Type1)`, `Sequence(Type2)`, `Bag(Type1)` and `Bag(Type2)`.
  - The types `Set(X)`, `Bag(X)` and `Sequence(X)` are all subtypes of `Collection(X)`.

Type conformance is transitive: if `Type1` conforms to `Type2`, and `Type2` conforms to `Type3`, then `Type1` conforms to `Type3` (standard rule for all types).

For example, if Bicycle and Car are two separate subtypes of Transport:

- `Set(Bicycle)` conforms to `Set(Transport)`
- `Set(Bicycle)` conforms to `Collection(Bicycle)`
- `Set(Bicycle)` conforms to `Collection(Transport)`

However

- `Set(Bicycle)` **does not conform** to `Bag(Bicycle)`, nor the other way around.
- They are both subtypes of `Collection(Bicycle)` at the same level in the hierarchy.
Use of OCL expressions in UML models (apart from class diagrams)

**OCL and State Diagrams**

Diagram:

```
Active --> Idle
```

Transition labels: Event[Condition]Action

- all three are optional

- **Event:**
  - if nil then when the task is completed we continue

- **Condition**
  - logical condition (transition occurs if its value is True)
  - the guards of transitions from a state must be mutually exclusive so that to have a unique next state

- **Action**
  - processes that occur quickly and are not interruptible

**OCL expression**

An OCL expression acting as value of a guard is of type Boolean.
The expression is evaluated at the moment that the transition attached to the guard is attempted.
OCL and State Diagrams (II)

OCL and State Diagrams (III)
Predefined properties on all objects
oclInState (s : OclState) : Boolean

This operation returns true if the object is in the state s.
Values for s are the names of the states in the statemachine(s) attached to the Classifier of object. For nested states the statenames can be combined using the double colon.

Here the values for s can be
- On
- Off
- Off::Standby
- Off::NoPower.

If the classifier of object has the above associated statemachine valid OCL expressions are:
- object.oclInState(On)
- object.oclInState(Off)
- object.oclInState(Off::Standby)
- object.oclInState(Off::NoPower)

If there are multiple statemachines attached to the object’s classifier, then the statename can be prefixed with the name of the statemachine containing the state and the double colon ‘::’, as with nested states.
**OCL and Interaction Diagrams**

OCL can be used for expressing the conditions under which a message (in a sequence or communication diagram) is sent.

![Interaction Diagram](image)

**Which UML CASE Tools support OCL and how?**

- We can attach OCL constraints to our diagrams using an appropriate stereotype and a dashed line should connect it to its contextual element.
- OCL constraints are exchanged using XMI.
- Tools that support OCL:
  - ArgoUML allows expressing them.
  - OCL Evaluator (a tool for editing, syntax checking & evaluating OCL).
  - Octopus OCL 2.0 Plug-in for Eclipse.
Assertions and Programming Languages

• Assertion techniques (preconditions, postconditions, invariants)
• History of assertion techniques:
  – Hoare 1972
  – Meyer 97a (he proposed the idea Design by Contract)
• Assertions support in Programming Language:
  – Eiffel supports them
  – In Java it is also possible (e.g. using JAF)
Techniques for adding Assertion Support in a PL

- **Built in**
  - Syntactic correctness of assertions is checked by the compiler
  - The runtime environment performs the runtime assertion checks

- **Preprocessing**
  - Formulate assertions separate from the program or include the assertions as comments. A preprocessor translates the assertions to program code
  - Pros: separation (separation of programmatic logic from contracts)
  - Cons: the original program code is modified (e.g. the line numbers of compiler errors do not fit the line numbers of the program)

- **Metaprogramming**
  - Traditionally this is possible only in dynamically typed and interpreted languages
  - Programs that have the possibility to reason about themselves have so-called **reflective capabilities** (Java has a reflection API)
  - The main advantage of metaprogramming approaches is that no specialized preprocessor has to be used but the native compiler. Nevertheless a specialized runtime environment has to be used to enable assertion checking

Assertions and Java

- "An assertion is a statement containing a boolean expression that the programmer believes to be true at the time the statement is executed".

- It is a facility provided within the Java programming language to test the correctness or assumptions made by your program. Assertions are checks provided within the system to ensure the smooth running of the program.

- **Why Assertions?**
  - **Why we need another level of checking when exceptions can do the job?**
  - Exceptions are primarily used to handle unusual (abnormal) conditions arising during program execution.
  - They do not guarantee smooth or correct execution of the program.
  - Assertions are used to specify conditions that a programmer assumes are true.
  - If a programmer can swear that the value being passed into a particular method is positive no matter what a calling client passes, it can be documented using an assertion to state it. Assertions help state scenarios that ensure the program is running smoothly. Assertions can be efficient tools to ensure correct execution of a program. They also improve the confidence about the program.
  - We can turn them off
Java Assertion Facility (JAF)

Syntax

```java
assert expression1;
```

The expression is the one we wish to assert as true. If the assumption fails, the expression evaluates to be false which means the assertion failed. In case the expression succeeds the program continues normally.

When an assertion fails the program throws an AssertionError on to the stack trace.

Examples:

```java
assert i<0;
assert (!myString.equals('""'));
```

Java Assertion Facility (JAF)
(builtin since Java 1.4)

Syntax

```java
assert expression1 : expression2;
```

The first argument takes a Boolean expression, while the second expression would be the resulting action to be taken if the assertion fails. The `Expression2` should be a value and can also be a result of executing a function. The compiler would throw an error if the second expression returns a void value.

When an assertion fails the program throws an AssertionError on to the stack trace. The program creates an object AssertionError with the return type of `Expression2`. The overloaded AssertionError constructor would then convert the returned data type into String and dump it on the stack trace with a meaningful message.

Examples:

```java
assert age>0 : "The value of age cannot be negative" +age;
assert ((i/2*23-12)>0):checkArgumentValue();
assert isParameterValid():throwIllegalParameterError();
```

In the last example the method checkArgumentValue() must return a value.
**OCL Constraints and Java**

**Context** Account: withdraw(amount: Real)
- **pre:** amount <= balance
- **post:** balance = balance@pre - amount

```java
class Account {
    private float balance = 0;
    public void withdraw(float amount) {
        assert amount <= balance;
        balance = balance - amount;
    }

    public float getBalance() {
        return balance;
    }
}
```

**Context** Account: getBalance(): Real
- **post:** result = balance

**Context** Employee: SetAge (age)
- **pre:** age > 0

```java
class Employee {
    public void SetAge(int age) {
        assert age > 0;
        this.age = age;
    }
}
```

**OCL Constraints in Java (2)**

**Context** Employee: SetAge (age)
- **pre:** age > 0

```java
class Employee {
    public void SetAge(int age) {
        assert age > 0;
        this.age = age;
    }
}
```

**Design by contract**

```java
class Employee {
    public void SetAge(int age) throws ArgumentException {
        if (age <= 0) {
            throw new ArgumentException("negative age");
        }
        this.age = age;
    }
}
```

**Defensive Programming**

```java
class Employee {
    public void SetAge(int age) throws ArgumentException {
        if (age <= 0) {
            throw new ArgumentException("negative age");
        }
        this.age = age;
    }
}
```
Example: Rectangle and Square

<table>
<thead>
<tr>
<th>Rectangle</th>
</tr>
</thead>
<tbody>
<tr>
<td>itsWidth</td>
</tr>
<tr>
<td>itsHeight</td>
</tr>
<tr>
<td>setWidth(w)</td>
</tr>
<tr>
<td>setHeight(h)</td>
</tr>
<tr>
<td>getArea()</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Square</th>
</tr>
</thead>
<tbody>
<tr>
<td>setWidth(w)</td>
</tr>
<tr>
<td>setHeight(h)</td>
</tr>
</tbody>
</table>

```cpp
class Rectangle {
public:
    virtual void setWidth(double w) {itsWidth=w;}
    virtual void setHeight(double h){itsHeight=h;}
    double getArea() {return itsHeight * itsWidth;}
private:
    double itsWidth;
    double itsHeight;
};

class Square: public Rectangle {
public:
    virtual void setWidth(double w);
    virtual void setHeight(double h);
};

void Square::setWidth(double w)
{Rectangle::setWidth(w); Rectangle::setHeight(w); } ;
void Square::setHeight(double h)
{Rectangle::setWidth(h); Rectangle::setHeight(h); };
```

Example: Rectangle and Square (II)

```cpp
class Rectangle {
public:
    virtual void setWidth(double w) {itsWidth=w;}
    virtual void setHeight(double h){itsHeight=h;}
    double getArea() {return itsHeight * itsWidth;}
private:
    double itsWidth;
    double itsHeight;
};

class Square: public Rectangle {
public:
    virtual void setWidth(double w);
    virtual void setHeight(double h);
};

void g(Rectangle* r)
{
r->setWidth(5);
r->setHeight(4);
assert(r->getArea()==20);
}
```

It will function correctly if \( r \) is a rectangle.
It will not function correctly if \( r \) is a square

The class Rectangle actually violates an “invariant” of the class Rectangle, specifically the *width-height independence*. 

```cpp
void g(Rectangle* r)
{
r->setWidth(5);
r->setHeight(4);
assert(r->getArea()==20);
}
```
Example: Rectangle and Square (III)

This could be expressed in OCL with
a post condition of setWidth: \textit{i.e. the height is the old value of height};
and a post condition of setHeight \textit{i.e. the width is the old value of width}.

\begin{align*}
\text{Context Rectangle:setWidth}(w) & \quad \text{post: } \text{itsWidth} = w \quad \text{and} \\
& \qquad \\text{itsHeight} = \text{itsHeight}@\text{pre} \\
\text{Context Rectangle:setHeight}(w) & \quad \text{post: } \text{itsHeight} = h \quad \text{and} \\
& \qquad \\text{itsWidth} = \text{itsWidth}@\text{pre}
\end{align*}

\[\text{[Meyers]:} \]
When we override a method A with a method B
the precondition of B should be that of A or a weaker condition, and
the postcondition of B should be that of A or a stronger (more strict) condition.

This reveals the problem in our example: the postcondition of Square:setWidth is weaker
(although it should be stronger according to the above rule).
So, if for example we had copied the postconditions of the Rectangle’s methods to the
methods of Square, we would have seen the problem while testing the class Square.

Reading and References

- How to download the current (UML 2.0 OCL) specification
- Tools
  - OCL Evaluator (a tool for editing, syntax checking & evaluating OCL)
  - Octopus OCL 2.0 Plug-in for Eclipse