Introduction to Scala

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Parallel Programming

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Introduction

- Part 1: Introduction to Scala
- Part 2: Concurrency in Scala
What is Scala?

- Scala is a statically typed language
  - Combines Object-Oriented Programming and Functional Programming
  - Developed in EPFL, lead by Martin Odersky
  - Influenced by Java, ML, Haskell, Erlang, and other languages

- Many high-level language abstractions
  - Uniform object model
  - Higher-order functions, pattern matching
  - Novel ways to compose and abstract expressions

- Managed language runtime
  - Runs on the Java Virtual Machine
  - Runs on the .NET Virtual Machine
Goals of Scala

- Create a language with better support for component software
- Hypotheses:
  - Programming language for component software should be scalable
    - The same concepts describe small and large parts
    - Rather than adding lots of primitives, focus on abstraction, composition, decomposition
  - Language that unifies OOP and functional programming can provide scalable support for components
Why use Scala?

- Runs on the JVM
  - Can use any Java code in Scala
  - Almost as fast as Java
- Much shorter code
  - Odersky reports 50% reduction in most code
  - Local type inference
- Fewer errors
  - No NullPointerException errors
- More flexibility
  - As many public classes per source file as you want
  - Operator overloading
- All of the above, for .NET too
Why learn Scala?

- Creating a trend in web service programming
  - LinkedIn
  - Twitter
  - Ebay
  - Foursquare
  - List is growing
Features of Scala (1)

- Both functional and object-oriented
  - Every value is an object
  - Every function is a value (including methods)
- Scala is statically typed
  - Includes local type inference system

**Java 1.5**

```java
Pair p = new Pair<Integer, String>(1, "Scala");
```

**Scala**

```scala
val p = new Pair(1, "Scala");
```
Features of Scala (2)

- Supports lightweight syntax for anonymous functions, higher-order functions, nested functions, currying
- ML-style pattern matching
- Integration with XML
  - Can write XML directly in Scala program
  - Can convert XML DTD into Scala class definitions
- Support for regular expression patterns
- Allows defining new control structures without using macros, and while maintaining static typing
- Any function can be used as an infix or postfix operator
- Can define methods named +, <= or ::
Features of Scala (3)

- Actor-based programming, distributed, concurrent
- Embedded DSLs, usable as scripting language
- Higher-kindred types, first class functions, closures
- Delimited continuations
- Abstract Types, Generics

Warning: Scala is the gateway drug to ML, Haskell, ...
An Example Class ...

**Java**

```java
public class Person {
    public final String name;
    public final int age;
    Person(String name, int age) {
        this.name = name;
        this.age = age;
    }
}
```

**Scala**

```scala
class Person(val name: String, val age: Int) {}
```
... and its use

Java

```java
import java.util.ArrayList;
...
Person[] people;
Person[] minors;
Person[] adults;
{
    ArrayList<Person> minorsList = new ArrayList<Person>();
    ArrayList<Person> adultsList = new ArrayList<Person>();
    for (int i = 0; i < people.length; i++)
        (people[i].age < 18 ? minorsList : adultsList).add(people[i]);
    minors = minorsList.toArray(people);
    adults = adultsList.toArray(people);
}
```

Scala

```scala
val people: Array[Person] = Array(
    new Person("Joe", 24),
    new Person("William", 23),
    new Person("Jack", 22),
    new Person("Averell", 21))
val (minors, adults) = people partition(_.age < 18)
```
Scala unifies class hierarchies and abstract data types (ADTs)
- Introduces pattern matching for objects
- Uses concise manipulation of immutable data structures
Example: Pattern matching

**Class hierarchy for binary trees**

abstract class Tree[T]

```scala
  case object Empty extends Tree[Nothing]
  case class Binary[T](elem: T, left: Tree[T], right: Tree[T]) extends Tree[T]
```

**In-order traversal**

```scala
def inOrder[T](t: Tree[T]): List[T] = t match {
  case Empty => List()
  case Binary(e, l, r) => inOrder(l) ::: List(e) ::: inOrder(r)
}
```

- Extensibility
- Encapsulation: only constructor params exposed
- Representation independence
Functions and Collections

- First-class functions make collections more powerful
- Especially immutable ones

**Container operations**

```scala
people.filter(_.age >= 18)
  .groupBy(_.surname)
  .values
  .count(_.length >= 2)
```
The Scala Object System

- Class-based
- Single Inheritance
- Can define singleton objects easily
- Subtyping is nominal: it is a subtype if declared to be a subtype
- Traits, compound types, views
  - Flexible abstractions
trait Nat;

object Zero extends Nat {
    def isZero: Boolean = true;
    def pred: Nat =
        throw new Error("Zero.pred");
}

class Succ(n: Nat) extends Nat {
    def isZero: Boolean = false;
    def pred: Nat = n;
}
Traits

- Similar to interfaces in Java
- They may have implementations of methods
- And can contain state!
- Can have multiple inheritance
Example: Traits

```
trait Similarity {
    def isSimilar(x: Any): Boolean;
    def isNotSimilar(x: Any): Boolean = !isSimilar(x);
}

class Point(xc: Int, yc: Int) extends Similarity {
    var x: Int = xc;
    var y: Int = yc;
    def isSimilar(obj: Any) =
        obj.isInstanceOf[Point] &&
        obj.asInstanceOf[Point].x == x;
}
```
Mixin Class Composition (1)

- **Mixin**: “A class which contains a combination of methods from other classes.”
- Basic inheritance model is single inheritance
- But mixin classes allow more flexibility

```scala
class Point2D(xc: Int, yc: Int) {
    val x = xc;
    val y = yc;
    // methods for manipulating Point2Ds
}
class ColoredPoint2D(u: Int, v: Int, c: String) extends Point2D(u, v) {
    var color = c;
    def setColor(newCol: String): Unit = color = newCol;
}
class Point3D(xc: Int, yc: Int, zc: Int) extends Point2D(xc, yc) {
    val z = zc;
    // code for manipulating Point3Ds
}
class ColoredPoint3D(xc: Int, yc: Int, zc: Int, col: String)
    extends Point3D(xc, yc, zc) with ColoredPoint2D(xc, yc, col);

// ERROR: cannot mixin classes with classes, only traits
```
Mixin Class Composition (2)

- Fix: extract the code to be added, into a trait
- Mixin the trait selectively into subclasses

```scala
class Point2D(xc:Int, yc:Int) {
    val x = xc;
    val y = yc;
    // methods for manipulating Point2Ds
}
trait Color {
    var color: String = null;
    def setColor(c: String): Unit = color = c;
}
class ColoredPoint2D(u:Int, v:Int, c:String) extends Point2D(u,v) with Color {
    color = c;
}
class Point3D(xc:Int, yc:Int, zc:Int) extends Point2D(xc,yc) {
    val z = zc;
    // code for manipulating Point3Ds
}
class ColoredPoint3D(xc:Int, yc:Int, zc:Int, col:String)
    extends Point3D(xc, yc, zc) with Color;
```
Mixin Class Composition (3)

- Mixin composition adds members explicitly defined in ColoredPoint2D (members that were not inherited).
- Mixing a class C into another class D is legal only as long as D’s superclass is a subclass of C’s superclass.
- i.e., D must inherit at least everything that C inherited.
- Why?
Mixin Class Composition (3)

- Mixin composition adds members explicitly defined in `ColoredPoint2D` (members that were not inherited).
- Mixing a class `C` into another class `D` is legal only as long as `D`’s superclass is a subclass of `C`’s superclass.
- *i.e.*, `D` must inherit at least everything that `C` inherited.
- Why?
- Remember that only members explicitly defined in `ColoredPoint2D` are mixin inherited.
- So, if those members refer to definitions that were inherited from `Point2D`, they had better exist in `ColoredPoint3D`.
  - They do, since `ColoredPoint3D` extends `Point3D` which extends `Point2D`
Views (1)

- Defines an *implicit coercion* from one type to another
- Similar to conversion operators in C++ and C#

```
trait Set[T] {
  def extend(x: T): Set[T]
  def contains(x: T): Boolean
}

// ...
implicit def list2set[T](list: List[T]) : Set[T] = new Set[T] {
  def extend(x: T): Set[T] = list2set(x :: list)
  def contains(x: T): Boolean =
    !list.isEmpty && ((list.head == x) || (list.tail contains x))
}
```
Implicit views are inserted automatically by the Scala compiler.

If \( e \) is of type \( T \) then a view is applied to \( e \) if:

- Expected type of \( e \) is not \( T \) (or a supertype)
- A member selected from \( e \) is not a member of \( T \)

Compiler uses only views in scope.
Lazy Views

- Many containers have lazy views
- Do not compute until absolutely necessary
- Different meaning but same name with implicit views (!)

```scala
scala> (1 to 1000000000).filter(_%2 ==0).take(10).toList
java.lang.OutOfMemoryError: GC overhead limit exceeded
  at java.lang.Integer.valueOf(Integer.java:832)
  at scala.runtime.BoxesRunTime.boxToInteger(BoxesRunTime.java:69)
  at scala.collection.immutable.Range.foreach(Range.scala:166)
  at scala.collection.TraversableLike.filterImpl(TraversableLike.scala:258)
  at scala.collection.TraversableLike.filter(TraversableLike.scala:270)
  at scala.collection.AbstractTraversable.filter(Traversable.scala:104)
... 26 elided

scala> (1 to 1000000000).view.filter(_%2 ==0).take(10).toList
res19: List[Int] = List(2, 4, 6, 8, 10, 12, 14, 16, 18, 20)
```
Variance Annotations (1)

class Array[A] {
    def get(index: Int): A
    def set(index: Int, elem: A): Unit
}

- **Array[String]** is not a subtype of **Array[Any]**
- If it were, we could do the following:

```scala
val x = newArray[String](1);
val y : Array[Any] = x;
y.set(0, new FooBar());
// just stored a FooBar in a String array!
```
Variance Annotations (2)

- Covariance is OK with functional data structures
- ... because they are immutable

```scala
trait GenList[+T] {
  def isEmpty: Boolean;
  def head: T;
  def tail: GenList[T]
}
object Empty extends GenList[Any] {
  def isEmpty: Boolean = true;
  def head: Any = throw new Error("Empty.head");
  def tail: GenList[Any] = throw new Error("Empty.tail");
}
class Cons[+T](x: T, xs: GenList[T]) extends GenList[T] {
  def isEmpty: Boolean = false;
  def head: T = x;
  def tail: GenList[T] = xs
}
```
Can also have contravariant type parameters
- Useful for an object that can only be written to

Scala checks that variance annotations are sound
- Covariant positions: Immutable field types, method results
- Contravariant: method argument types
- Type system ensures that covariant parameters are only used covariant positions
- (similar for contravariant)

If no variance specified, then Invariant
- Neither superclass, nor subclass
Functions are Objects

- Every function is a value
  - Values are objects, so functions are also objects
- The function type $S \Rightarrow T$ is equivalent to the class type `scala.Function1[S, T]

```
trait Function1[-S, +T] {
  def apply(x: S): T
}
```

- For example, the anonymous successor function $(x: \text{Int}) \Rightarrow x + 1$ or in shorter code $(\_ + 1)$ expands to

```
new Function1[Int, Int] {
  def apply(x: Int): Int = x + 1
}
```
Arrays are Objects

- Arrays (mathematically): Mutable functions over integer ranges

**Syntactic Sugar**

\[
a(i) = a(i) + 2 \text{ for a.update(i, a.apply(i) + 2)}
\]

**Example**

```scala
final class Array[T](_length: Int)
  extends java.io.Serializable
  with java.lang.Cloneable {
    def length: Int = ...
    def apply(i: Int): T = ...
    def update(i: Int, x: T): Unit = ...
    override def clone: Array[T] = ...
}
```
Partial Functions

- Functions that are defined only for some objects
- Test using `isDefinedAt`

**Example**

```scala
trait PartialFunction[-A, +B] extends (A => B) {
  def isDefinedAt(x: A): Boolean
}
```

- Blocks of pattern-matching cases are instances of partial functions
- This lets programmers write control structures that are not easy to express otherwise
Automatic Closure Construction

- Allows programmers to make their own control structures
- Can tag the parameters of methods with the modifier `=>`
- When method is called, the actual `=>` parameters are not evaluated and a no-argument function is passed
Example: Custom loop construct

object TargetTest1 {
  def loopWhile(cond: => Boolean)(body: => Unit): Unit =
    if (cond) {
      body;
      loopWhile(cond)(body);
    }

  def main(args: Array[String]) {
    var i = 10;
    loopWhile (i > 0) {
      Console.println(i);
      i = i - 1;
    }
  }
}
Types as Class Members

abstract class AbsCell {
  type T;
  val init: T;
  private var value: T = init;
  def get: T = value;
  def set(x: T): Unit = { value = x }
}
def createCell() : AbsCell =
  new AbsCell { type T = Int; val init = 1 }

Clients of `createCell` cannot rely on the fact that T is `Int`, since this information is hidden from them.
Scala Parallel Collections

```scala
val list = (1 to 10000).toList
list.map(_ + 42)
```

- Sequential map, addition
Scala Parallel Collections

```scala
val list = (1 to 10000).toList
list.par.map(_ + 42)
```

- Parallel list
- Many data structures available
  - ParArray
  - ParVector
  - mutable.ParHashMap
  - mutable.ParHashSet
  - immutable.ParHashMap
  - immutable.ParHashSet
  - ParRange
  - ParTrieMap
val lastNames = List("Smith","Jones","Frankenstein","Bach","Jackson","Rodin").par
lastNames.map(_.toUpperCase)

val parArray = (1 to 10000).toArray.par
parArray.fold(0)(_ + _)

val lastNames = List("Smith","Jones","Frankenstein","Bach","Jackson","Rodin").par
lastNames.filter(_.head >= 'J')
import scala.collection.parallel.immutable.ParVector
val pv1 = new ParVector[Int]
val pv2 = Vector(1,2,3,4,5,6,7,8,9).par
Parallel Collections

- Side-effecting operations can lead to non-determinism
  - side effects are reordered or concurrent
- Non-associative operations lead to non-determinism
  - order of operations changes
Example: Race!

```scala
var sum = 0
val list = (1 to 1000).toList.par
list.foreach(sum += _);
sum
// something

var sum = 0
list.foreach(sum += _);
sum
// something else
```
Example: Associativity

```scala
val list = (1 to 1000).toList.par
list.reduce(_-_) // some result
list.reduce(_-_) // some other result
list.reduce(_-_) // yet another result, depending on what subtraction runs first
```
The Actor Model

- A model of concurrent computation
- Introduced in 1973 (Lisp, Simula)
- Main idea: *Everything is an Actor*
  - Similar to OO idea that *Everything is an Object*
- An actor can:
  - Send messages to other actors
  - Create new actors
  - React to messages it receives
- There is no constraint on order between these
  - Can occur in parallel across actors, also for any actor
  - Parallel computation and communication
A c t o r s i n S c a l a

- Initial built-in implementation
- Language primitives
- Built into the language
  - Obsolete now
- Integration with Akka library
  - Akka: library with distributed actors
  - Concurrency
  - Scalability
  - Fault-tolerance
  - Single unified programming model
  - Managed runtime (contained into the library)
  - Open Source
Actors in Akka

- Goal: Program at very high level of abstraction
- Do not think of shared state, threads, state visibility, locks, collections, etc.
- Only think how messages flow into the system
- Runtime system does the rest
  - High CPU utilization
  - Low latency
  - Scalability
  - Built-in support for error detection and recovery
Parallel and Distributed

- Akka actors are distributable by design
  - Designed to scale up (more threads) and scale out (more nodes)
  - Same program, different deployments
  - Perfect for cloud deployment
    - Elastic, dynamic
    - Fault-tolerant, self-healing
    - Adaptive load-balancing, migration
    - Loosely coupled, allows dynamic changes at runtime
What is an Actor

- Unit of code organization in Akka
- Actors help create concurrent, scalable and fault-tolerant applications
- Like Java-EE Servlets and session beans, Actors help organize code to keep “policy” and “business logic” separate
- Used in telecom systems with “9 nines” uptimes
- Abstraction intuitively: Virtual Machines in the Cloud (but faster)
  - Encapsulated, decoupled, black boxes
  - Manage their own memory and behavior
  - Communicate asynchronously, non-blocking messages
  - Can grow and shrink on demand, add new actors, stop some
  - Hot-deploy: change behavior at runtime, add new components, new code
  - Actors are the same, but for a single application
Actor uses

May be alternative to:

» Thread
» Object instance, component
» Callback Listener
» Singleton, service
» Load-balancer, router, thread pool
» Jave EE Session Bean, Message-Driven Bean
» Out-of-process service
» FSM
Theoretical definition

- Fundamental unit of computation that embodies:
  - Processing
  - Storage
  - Communication

- 3 axioms - When an actor receives a message, it can:
  - Create new actors
  - Send messages to actors it knows
  - Designate how it should handle the next message received
Core Actor operations

- Define
- Create
- Send
- Become
- Supervise
import akka.actor._

class Summer extends Actor {
  var sum = 0

  def receive = {
    case ints: Array[Int] =>
      sum += ints.reduceLeft((a, b) => (a+b) % 7)
    case "print" => println("Sum:" + sum)
  }
}
Create an Actor

- Create an instance of an Actor
- Very lightweight in Akka: 2.7 million actors per GB RAM
- Very strong encapsulation:
  - state
  - behavior
  - message queue
- State and behavior are indistinguishable
- Only way to observe state: send a message, see reaction
import akka.actor._
class Summer extends Actor {
    var sum = 0

    def receive = {
        case ints: Array[Int] =>
            sum += ints.reduceLeft((a, b) => (a+b) % 7)
        case "print" => println("Sum:" + sum)
    }
}

val system = ActorSystem("SummerSystem")
val summer = system.actorOf(Props[Summer], name = "summer")
Actors form Hierarchies

- **system** is “guardian actor”
- Can create actors with `context.actorof()`, guarded by creating actor
- Hierarchies can be tall trees
- Name resolution works like a file system: Actor `/summer/someother`
Send Messages

- Asynchronous and non-blocking: “Fire and Forget”
- Everything happens Reactively
  - An Actor is passive until a message is sent to it
  - Messages are “kinetic energy” in Actor System
  - But light messages may trigger heavy reactions
- Everything is asynchronous and lockless
- Lightweight: single machine can handle millions of messages per second
import akka.actor._
class Summer extends Actor {
    var sum = 0
    def receive = {
        case ints: Array[Int] =>
            sum += ints.reduceLeft((a, b) => (a+b) % 7)
        case "print" => println("Sum:" + sum)
    }
}
val system = ActorSystem("SummerSystem")
val summer = system.actorOf(Props[Summer], name = "summer")
summer tell (1 to 10).toArray
summer ! (1 to 20).toArray
import akka.actor._

class SomeActor extends Actor {
    def receive = {
        case User(name) =>
            sender tell ("Hi " + name)
    }
}
Remote Deployment

```scala
akka {
    actor {
        provider = akka.remote.RemoteActorRefProvider
        deployment {
            /Summer {
                remote = akka://SummerSystem@machine42:31337
            }
        }
    }
}
```
Dynamically redefine actor behavior
Triggered reactively by receiving a message
Type system analogy: Object changes type
  ▶ change interface, protocol, implementation
Actor will now react differently to messages
Behaviors are stacked, can be pushed and popped
Why?

- Let an actor with high contention become load-balancer, distribute work “behind”
- Implement FSM
- Graceful degradation
- Generic Worker easy spawn, becomes whatever is needed etc.
- Very useful once you get used to it
context become {
  case NewMessage =>
    ...
}

Pratikakis (CSD)  Scala  CS342, 2018  59/66
Example: load balancing

```scala
val router = 
  system.actorOf( 
    Props[SomeActor].withRouter( 
      RoundRobinRouter(nrOfInstances = 5) 
    )
  )
```

Example: load balancing++

```scala
val resizer = DefaultResizer(lowerBound = 2, upperBound = 15)

val router = 
  system.actorOf(
    Props[SomeActor].withRouter(
      RoundRobinRouter(resizer = Some(resizer))
    )
  )
```
Failure Management, Traditionally

- Single thread of control
- If thread blows up, we’re $#%@ed
- Must do explicit error handling within thread
- Errors do not propagate between threads
  - No way to find out if something broke
- Leads to defensive programming
  - if(printf()) ...
    - Error handling tangled with business logic
    - Error checking salted all over the codebase
- Things shouldn’t be that bad
Supervise

- Manage another Actor’s failures
- Error handling in actors by letting Actors monitor (supervise) each other for failure
- If an Actor crashes, notification will be sent to supervisor
- Clean separation of processing and error handling
- Every actor has default supervisor strategy, usually sufficient
Example: Supervision

class Supervisor extends Actor {
  override val supervisorStrategy =
    OneForOneStrategy(maxNrOfRetries = 10, withinTimeRange = 1 minute)
  {
    val worker = context.actorOf(Props[Worker])

    def receive = {
      case n: Int => worker forward n
    }
  }
}
class Supervisor extends Actor {
  override val supervisorStrategy =
    AllForOneStrategy(maxNrOfRetries = 10, withinTimeRange = 1 minute)
  {
    val worker = context.actorOf(Props[Worker])

    def receive = {
      case n: Int => worker forward n
    }
  }
}
class Worker extends Actor {
    ...
    override def preRestart(reason: Throwable, message: Option[Any]) {
        // Clean up before restart
    }

    override def postRestart(reason: Throwable) {
        // Initialize after restart
    }
}