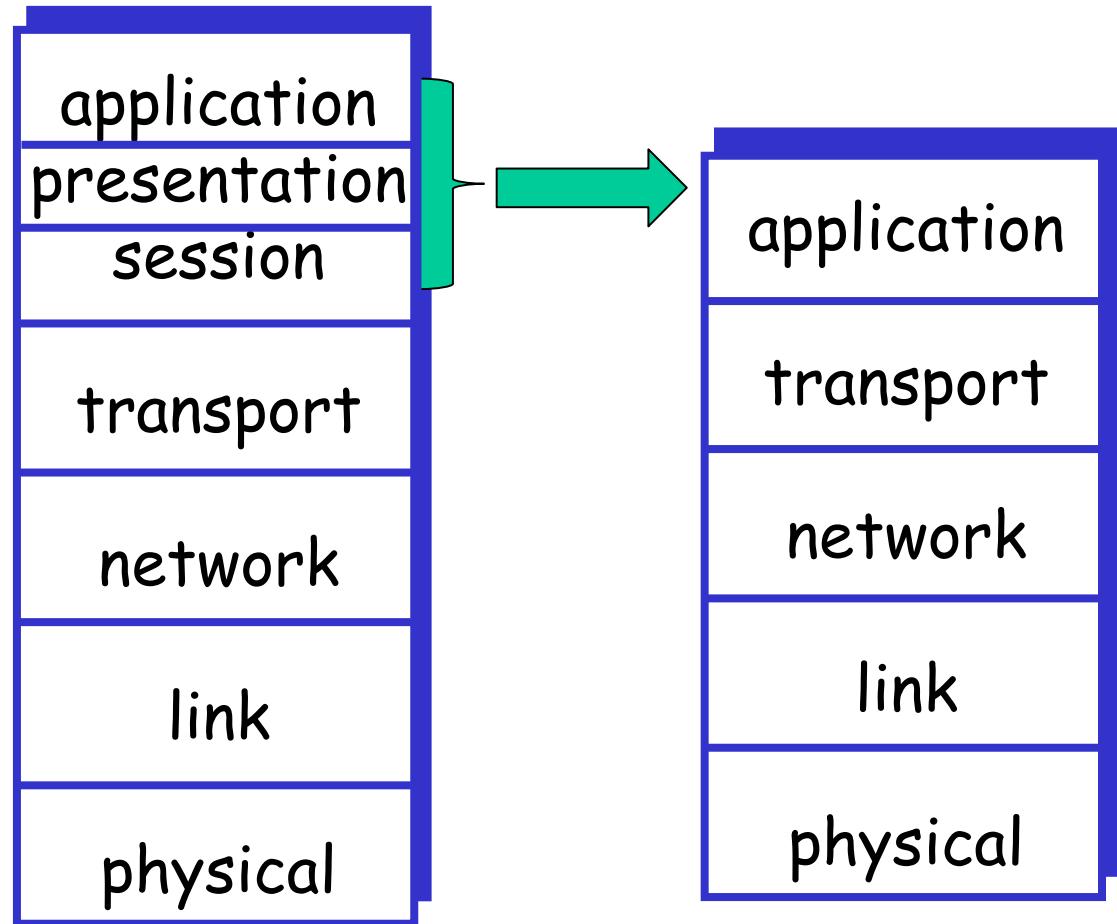


- Internet protocol stack
- Encapsulation
- Connection oriented VS connectionless services
- Circuit Switching
- Packet Switching
- Store-and-forward switches
- Multiplexing: TDM, FDM, Statistical multiplexing, CDMA

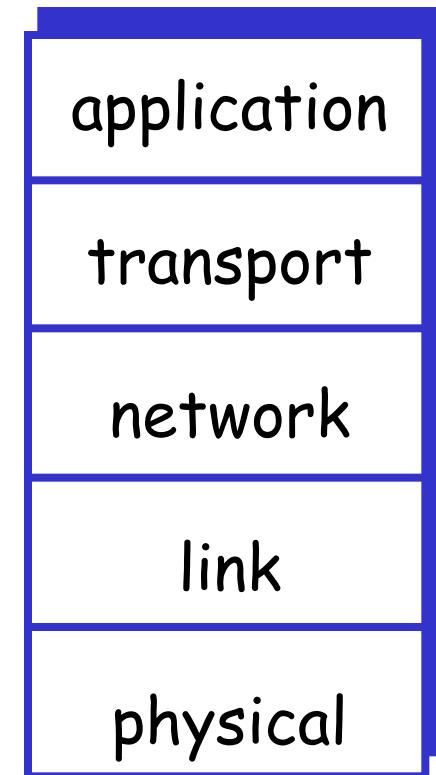
OSI VS Internet protocol stack



- Internet stack “missing” these layers!
 - ❖ these services, *if needed*, must be implemented in application
 - ❖ needed?

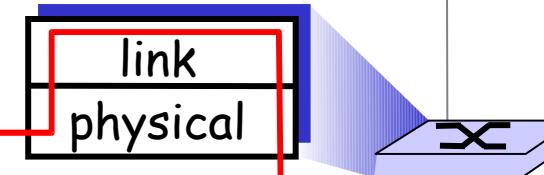
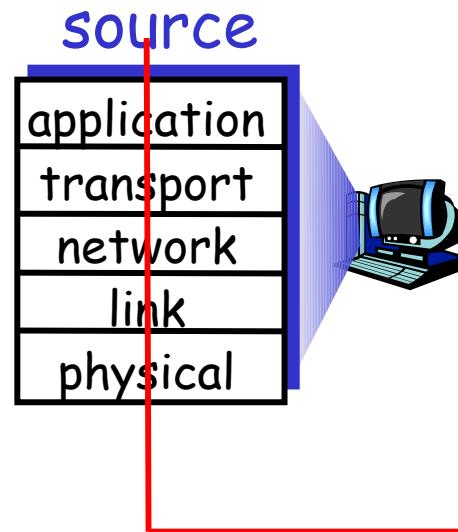
Internet protocol stack

- **application:** supporting network applications
 - ❖ FTP, SMTP, HTTP, DNS protocols
- **transport:** process-process data transfer
 - ❖ TCP, UDP protocols
- **network:** routing of datagrams from source to destination
 - ❖ IP, routing protocols
- **link:** data transfer between neighboring network elements
 - ❖ PPP, Media Access Control (Ethernet, DSL, ISDN, FDDI)
- **physical:** bits "on the wire"

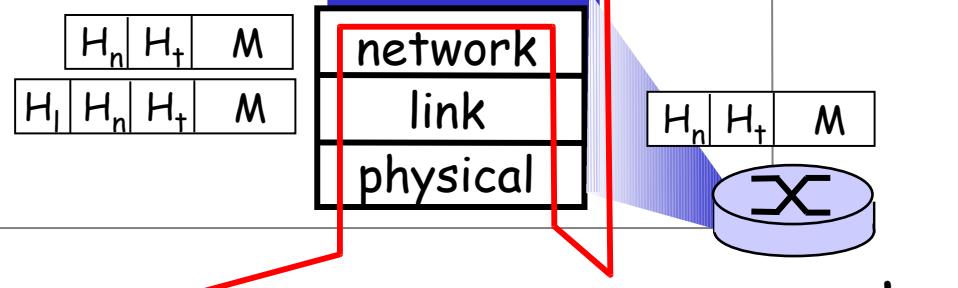
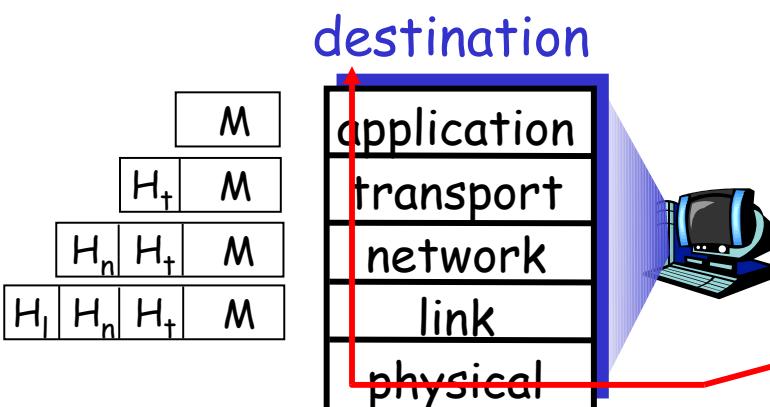


Encapsulation

message	M
segment	H _t M
datagram	H _n H _t M
frame	H _l H _n H _t M



switch



router

Why layering?

Dealing with complex systems:

- Abstraction
- explicit structure allows identification, relationship of complex system's pieces
 - ❖ layered **reference model** for discussion
- modularization eases maintenance, updating of system
 - ❖ change of implementation of layer's service transparent to rest of system
 - ❖ e.g., change in gate procedure doesn't affect rest of system

Connection Oriented services

- Establish end to end logical or physical connection before any data are sent
- Involves handshaking
- Reliable data transfer may be involved (e.g. TCP)

Data link layer examples:

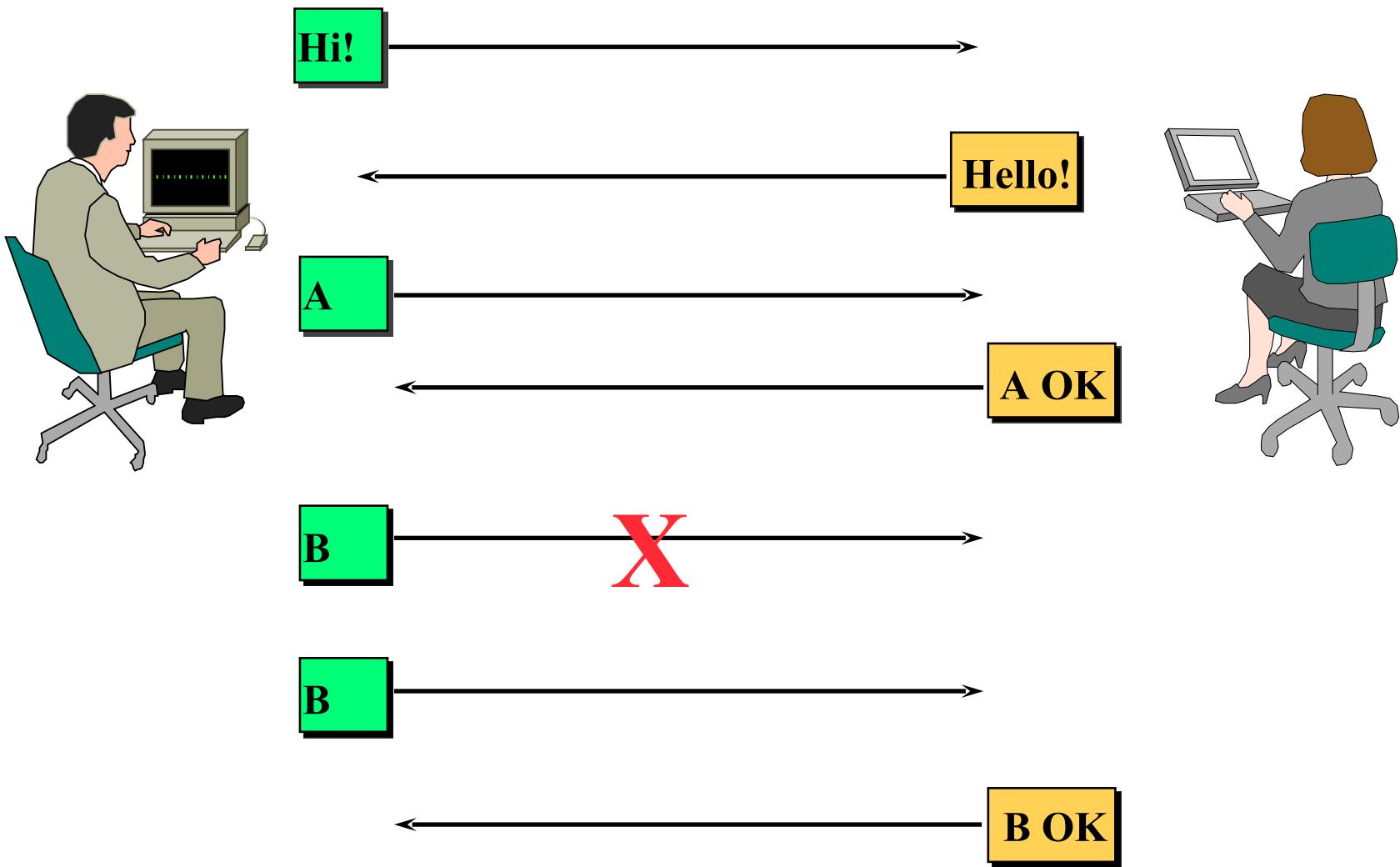
- Circuit mode communication
- Virtual Circuits (packet switching). Same path! We just need a VCI.

Transport layer examples:

- TCP

Transmission Control Protocol (TCP)

(belongs to transport layer)



Connectionless services

- No handshaking!
- Each data packet carries information about the destination address (datagram)

Network layer example:

- IP protocol

Transport layer examples:

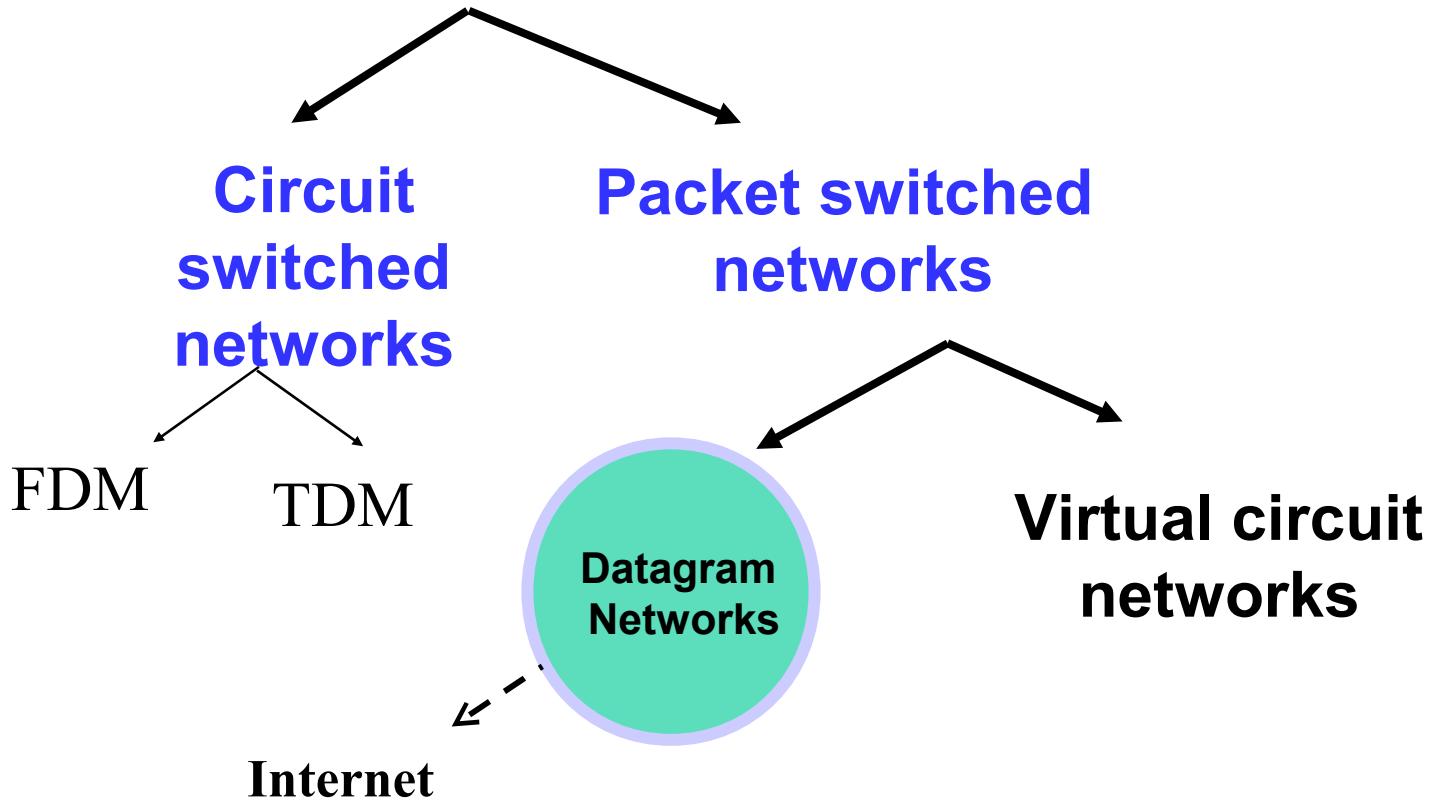
- UDP

Connection oriented VS connectionless services

- The distinction takes place in several layers
- Packet switching examples in both categories
- Connection oriented service on connectionless service??

TCP/IP

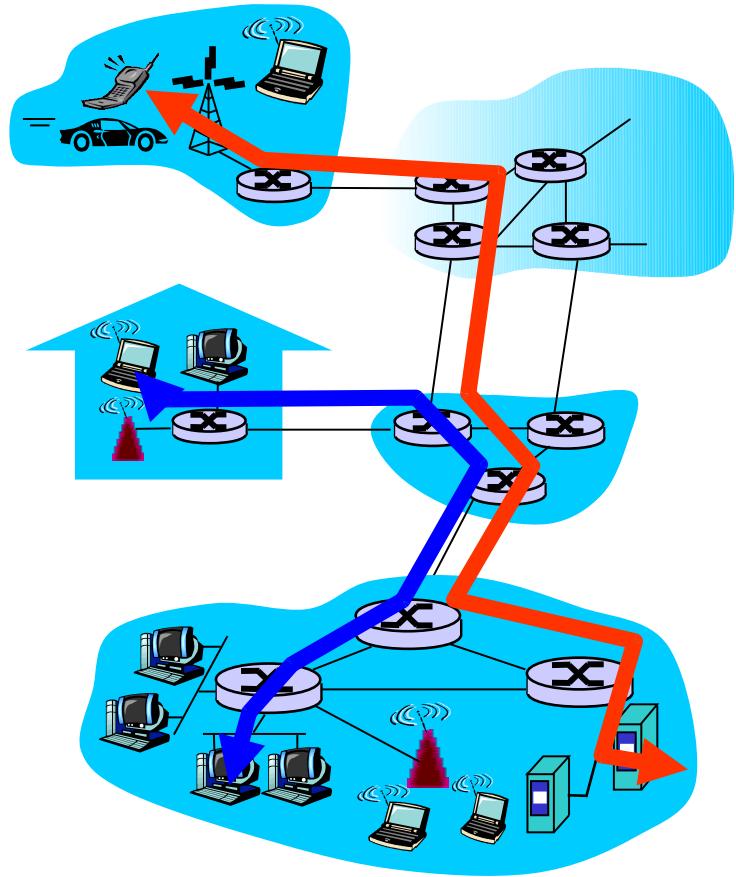
Switched networks



Circuit Switching

End-end resources reserved for "call"

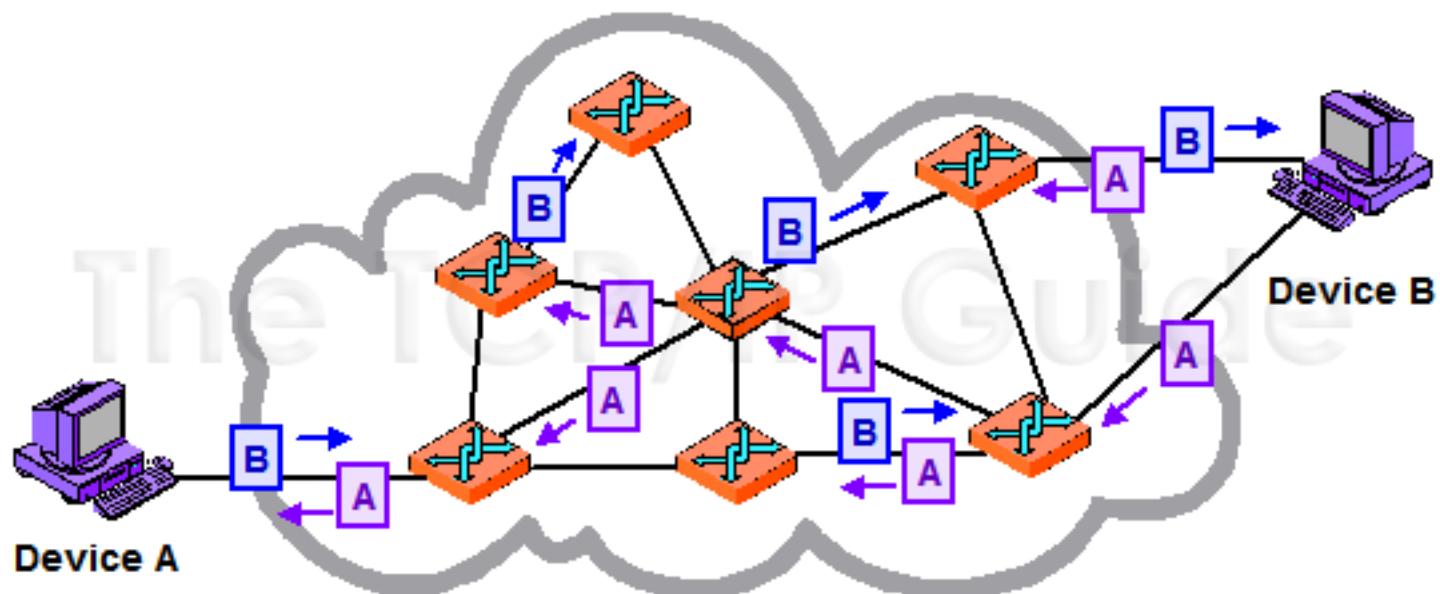
- link bandwidth, switch capacity
- dedicated resources: no sharing (??)
- circuit-like (guaranteed) performance
- call setup required



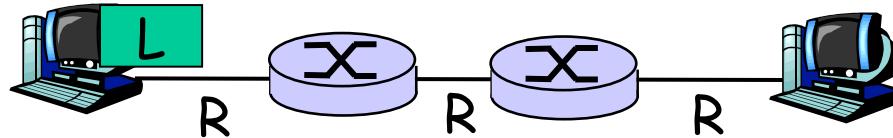
Packet Switching

each end-end data stream divided into *packets*

- user A, B packets *share* network resources
- each packet uses full link bandwidth
- resources used as needed
- Same route??



Packet-switching: store-and-forward



- takes L/R seconds to transmit (push out) packet of L bits on to link at R bps
- Store and forward:*** entire packet must arrive at router before it can be transmitted on next link
- delay = $3L/R$ (assuming zero propagation delay)

Example:

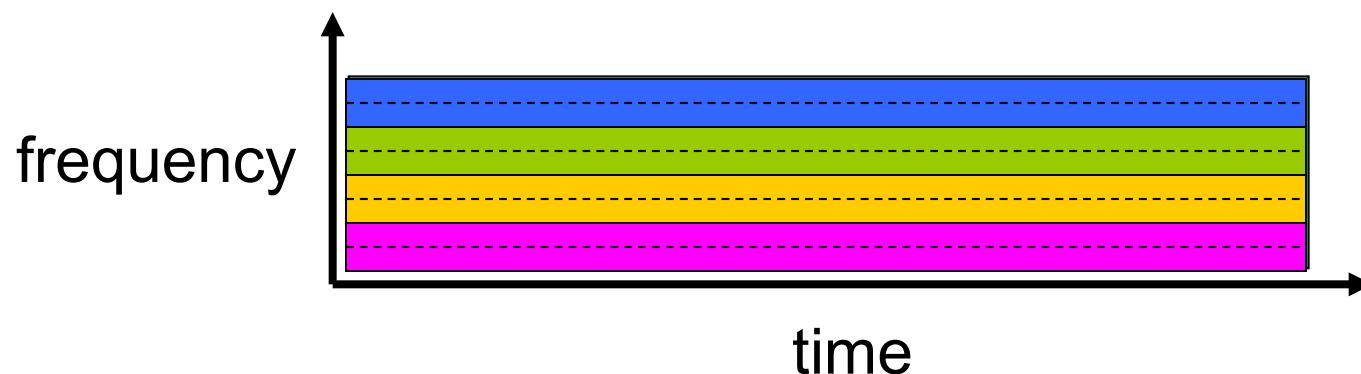
- $L = 7.5 \text{ Mbits}$
- $R = 1.5 \text{ Mbps}$
- transmission delay = 15 sec

Circuit Switching: FDM and TDM

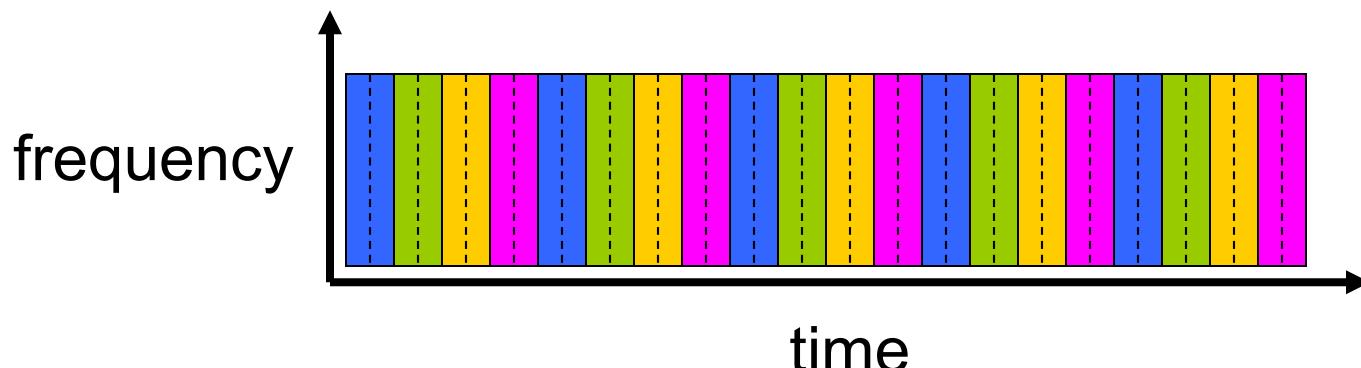
Example:

FDM

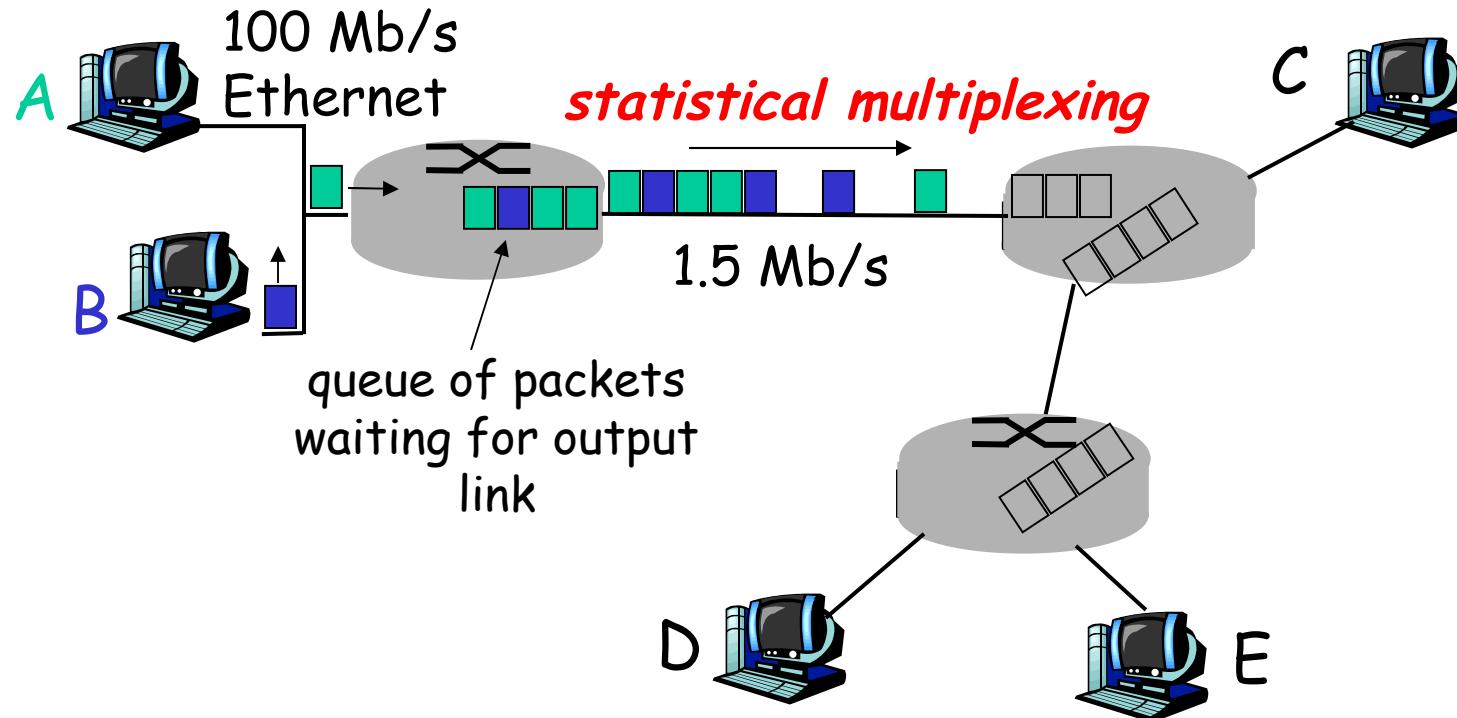
4 users



TDM



Packet Switching: Statistical Multiplexing



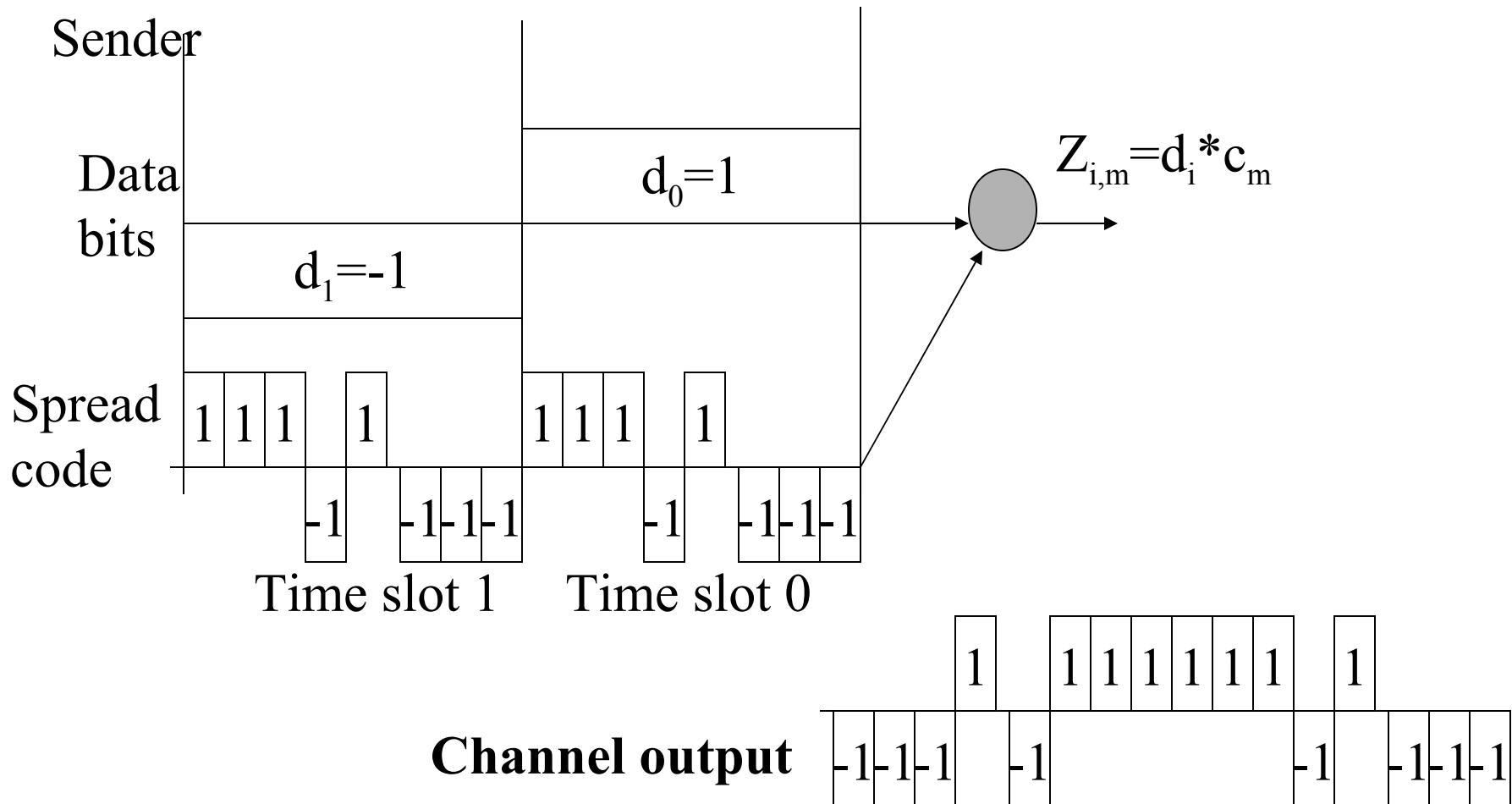
Sequence of A & B packets does not have fixed pattern,
bandwidth shared on demand ***statistical multiplexing***.

TDM: each host gets same slot in revolving TDM frame.

Πολλαπλή Πρόσβαση Διαίρεσης Κύδικα (CDMA)

- Στο CDMA ορίζεται σε κάθε κόμβο ένας διαφορετικός κώδικας
- Οι κώδικες είναι **ορθογώνιοι μεταξύ τους** (δηλ. το εσωτερικό γινόμενο μεταξύ οποιωνδήποτε δύο κωδίκων είναι 0)
- Κάθε κόμβος χρησιμοποιεί το δικό του μοναδικό κώδικα για να κωδικοποιήσει τα bits των δεδομένων που στέλνει
- Οι κόμβοι μπορούν να εκπέμπουν ταυτόχρονα
- **Πολλαπλοί κόμβοι σε κάθε κανάλι**
- Οι αντίστοιχοι προς αυτούς δέκτες
 - ❖ Λαμβάνουν σωστά τα κωδικοποιημένα bits δεδομένων ενός πομπού
 - Θεωρώντας ότι ο δέκτης γνωρίζει τον κώδικα του πομπού, παρά τις παρεμβαλλόμενες μεταδόσεις άλλων κόμβων

Παράδειγμα CDMA



Παράδειγμα CDMA (συνέχεια)

- Όταν δεν υπάρχουν παρεμβάλλοντες πομποί

- ❖ Ο δέκτης

- Λαμβάνει τα κωδικοποιημένα bits
 - Ανακτά τα αρχικά bit δεδομένων, d_i , υπολογίζοντας το

$$d_i = \frac{1}{M} \sum_{m=1}^M Z_{i,m} * c_m$$

- Τα παρεμβάλλοντα εκπεμπόμενα δυαδικά σήματα είναι προσθετικά