



# CS240 — DATA STRUCTURES

## 1st Series of Exercises

### Submission Instructions

Exercises can be submitted to the course assistants on **Monday, October 16, 2023, from 11:00 AM to 12:00 PM** at the TAs' office (**B.208 / B.210**). Exercises submitted after 12:00 PM on Monday, 16/10/2023, will have a penalty. **Late submissions are accepted in electronic format**, and they must be submitted using the `turnin` program. For more information visit the course website.

<b>SEMESTER:</b>	Spring (2023-24)
<b>UNIVERSITY:</b>	University of Crete
<b>DEPARTMENT:</b>	Computer Science
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<b>LAST MODIFICATION:</b>	10 / 10 / 2023

## Exercise 1

[20 points]

Let  $T_1(n)$  be the total number of times the  $x = x + 3$  command is executed by procedure `Puzzle1()`, and  $T_2(n)$  be the total number of times the  $x = x + 3$  command is executed by procedure `Puzzle2()`. Calculate  $T_1(n)$  and  $T_2(n)$ . Analyze the order (denoted as  $O$  notation) of  $T_1(n)$  and  $T_2(n)$ .

```
procedure Puzzle1(integer n) {  
    for i = 1 to sqrt(n) do  
        for j = i to n do  
            for k = 1 to 2n do  
                x = x+3;  
            }  
        }  
    }  
}
```

```
procedure Puzzle2(integer n)  
    for i = 0 to 2n do  
        if i is even then {  
            for j = i to n do x = x+3;  
        }  
    }  
}
```

## Exercise 2

[25 points]

a. Which of the following is true and why?

•  $n^2 - 2n - 7 = \Theta(n^2)$  [5P]

•  $n \times \sum_{i=1}^n (3 \times (i + 1)) = O(n^4)$  [5P]

•  $n^2 \log n^2 - n \log n - n = \Theta(n^2 \log n)$  [5P]

b. Study the asymptotic complexity (using the  $\Theta$  notation) of the function:

$$f(n) = 10n^3 \log n^6 + n^3 \sqrt{n} - n \log n - 5 \quad [10P]$$

## Exercise 3

[20 points]

Solve the following recursive relations using iterative substitution.

a.  $T(n) := \begin{cases} 1, & \text{for } n \in \{0, 1\} \\ 3T(\lfloor \frac{n}{3} \rfloor) + 3n, & \text{otherwise} \end{cases}$  [5M]

b.  $T(n) := \begin{cases} 1, & \text{for } n \leq 0 \\ T(n - c) + c \cdot n, & \text{otherwise, where } c \text{ is an integer constant} \end{cases}$  [5M]

For the recursive relation of question a, the following are requested:

c. Draw the recursion tree. [5P]

d. Verify, using mathematical induction, that the solution you found with iterative substitution is correct. [5P]

## Exercise 4

[15 points]

Consider the following function, which performs Egyptian multiplication:

$$m(x, y) := \begin{cases} 0, & \text{if } y = 0 \\ m(x + x, \frac{y}{2}), & \text{if } y \text{ even, and } y \neq 0 \\ x + m(x, y - 1), & \text{if } y \text{ odd, and } y \neq 0 \end{cases}$$

- a. Present pseudocode for a recursive function [5P]

```
int EgyptianMultiplication(int x, int y)
```

which will perform Egyptian multiplication according to the above recursive formula.

- b. Trace the execution of `EgyptianMultiplication(5, 20)` (i.e., provide a detailed explanation of its execution, as well as the activation records in memory, when called with  $x = 5$  and  $y = 20$ ). [10P]

## Exercise 5

[20 points]

`InsertionSort()` can be expressed as a recursive procedure in the following way. To sort the array  $A[1..n]$ , we recursively sort the subarray  $A[1..n-1]$  and then we insert the element  $A[n]$  into the sorted subarray  $A[1..n-1]$ .

- Present pseudocode for this recursive version of `InsertionSort()`.
- Formulate a recursive relation for the runtime of this recursive version of `InsertionSort()`.
- Solve the recursive relation you proposed in question b. and study (based on the  $O$  and  $\Omega$  notations) the order of complexity of the recursive version of `InsertionSort()`.
- What is the space complexity of the algorithm (i.e., how much memory is required to run the algorithm)?